

COMPLETE CHARACTERIZATION OF PERFECTLY SECURE STEGO-SYSTEMS WITH MUTUALLY INDEPENDENT EMBEDDING OPERATION

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ABSTRACT

Without any assumption on the cover model, this paper presents a complete characterization of all perfectly secure stego-systems that employ mutually independent embedding operation. We show that for a fixed embedding operation, the only perfectly secure stego-systems are those whose cover distribution is an element of a linear vector space with basis vectors determined by the embedding operation. Finally, we show that for mutually independent embedding operation, perfect security as defined by Cachin is equivalent to positivity of Fisher information with respect to the embedding change rate. This result is important for deriving steganographic capacity of covers modeled as Markov chains [1].

Index Terms— steganography, perfect security, mutually independent embedding

1. INTRODUCTION

In steganography, the sender and receiver communicate by hiding their messages in generally trusted media, such as digital images, so that no warden can distinguish between the original (cover) object and the object carrying the message—the stego object. Formally, the security of a stego-system is evaluated using the Kullback-Leibler divergence between the distributions of cover and stego objects [2]. Systems with zero KL divergence are called perfectly secure.

Formally, a stego-system is a combination of an embedding algorithm and a cover source. The vast majority of practical stego-systems hide messages by modifying individual elements of the cover using mutually independent embedding operations. This is the case, for example, for stego-systems that use LSB and ± 1 embedding, the F5 algorithm [3] and its variations [4], perturbed quantization [4], MMx [4], as well as for algorithms making larger modifications, such as variants of 2LSB embedding [5] or stochastic modulation [6].

In this paper, we provide a complete characterization of perfectly secure stego-systems for the class of embedding algorithms with mutually independent (MI) embedding operations. The cover distributions of all perfectly secure systems form a linear vector space spanned by distributions determined by the embedding operation. Moreover, we show that perfect security (zero KL divergence) is equivalent to satisfying a simple condition related to Fisher information.

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The paper is structured as follows. In Section 2, we introduce the notation and definitions and review some preliminary facts that will be put to use later. Section 3 contains the main results of this paper as well as examples of specific stego-systems to illustrate the theoretical contributions. Section 5 concludes this work and further elaborates on its importance.

2. NOTATION, PRELIMINARIES, AND ASSUMPTIONS

We will represent an n -element cover object with $x_1^n \triangleq (x_1, \dots, x_n) \in \mathcal{X}^n$, $\mathcal{X} = \{1, \dots, N\}$, where x_1^n is a realization of a random variable X_1^n distributed according to some general distribution P over \mathcal{X}^n . The stego object $y_1^n \triangleq (y_1, \dots, y_n) \in \mathcal{X}^n$ is assumed to be a realization of random variable Y_1^n distributed according to stego distribution Q_β , where β is a scalar parameter capturing the extent of embedding changes (e.g., it will be helpful to think of β as the change rate).

The definition of steganographic security was given by Cachin [2].

Definition 1 Let P, Q_β be probability distributions of cover, stego objects with n elements embedded with parameter β , respectively. Steganography is said to be perfectly secure iff

$$d(\beta) \triangleq D_{KL}(P||Q_\beta) = \sum_{y_1^n \in \mathcal{X}^n} P(y_1^n) \log \frac{P(y_1^n)}{Q_\beta(y_1^n)} = 0,$$

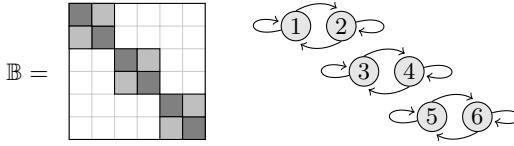
or ϵ -secure if $d(\beta) \leq \epsilon$.

We assume the impact of embedding with parameter $\beta \in [0, \beta_0]$ on the k -th element can be captured using the matrix $b_{i,j}(\beta) \triangleq Pr(Y_k = j | X_k = i) = \delta_{i,j} + \beta c_{i,j}$, for some constants $c_{i,j} \geq 0$ for $i \neq j$, $c_{i,i} = -\sum_j c_{i,j}$, where $\delta_{i,j}$ is the Kronecker delta. In a matrix form, $\mathbb{B}_\beta = \mathbb{I} + \beta \mathbb{C}$, where $\mathbb{B}_\beta \triangleq (b_{i,j}(\beta))$, \mathbb{I} is the identity matrix, and $\mathbb{C} \triangleq (c_{i,j})$. We further assume that the embedding operations are mutually independent, $Pr(Y_1^n | X_1^n) = \prod_{k=1}^n Pr(Y_k | X_k)$. By the definition of $b_{i,j}$, the matrix \mathbb{B}_β is stochastic, $\sum_j b_{i,j} = 1$. Finally, we assume that $b_{i,i}(\beta) > 0$ for all $\beta \in [0, \beta_0]$. Many embedding methods can be represented in this framework (see examples in Figure 1). We use matrix \mathbb{B}_β as a representation of an embedding algorithm with MI embedding operation (simply MI embedding).

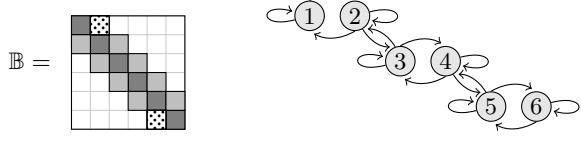
To simplify the language in this paper, we will speak of security of a cover source w.r.t. given MI embedding meaning that the *cover model is perfectly secure w.r.t. \mathbb{B}* , if the resulting stego-system is perfectly secure. It does then make sense to inquire about all possible perfectly secure cover sources w.r.t. MI embedding with matrix \mathbb{B} .

In the rest of this section we review some results from the theory of ergodic classes [7] that will be later applied to the stochastic matrix \mathbb{B}_β . For states $i, j \in \mathcal{X}$, we call j a *consequent* of i

LSB embedding: $\blacksquare = 1 - \beta$ $\square = \beta$



± 1 embedding: $\blacksquare = 1 - \beta$ $\square = \frac{1}{2}\beta$ $\boxtimes = \beta$



F5: $\blacksquare = 1 - \beta$ $\square = \beta$ $\boxtimes = 1$

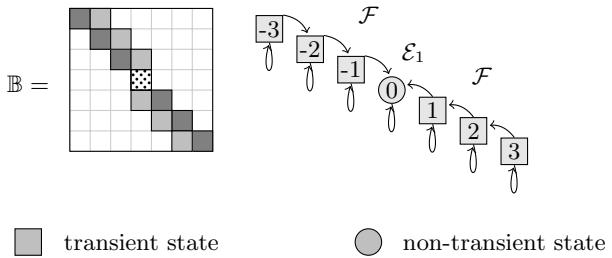


Fig. 1. Examples of several embedding methods and their ergodic classes.

$(i \rightarrow j)$ iff $\exists k, (\mathbb{B}_\beta^k)_{i,j} \neq 0$. We classify each state $i \in \mathcal{X}$ as *transient* if it has a consequent of which it is not itself a consequent, i.e., $\exists j \in \mathcal{X}$ such that $(i \rightarrow j) \Rightarrow (j \not\rightarrow i)$. We say $i \in \mathcal{X}$ is *non-transient* if it is a consequent of every one of its consequents, $\forall j \in \mathcal{X}, (i \rightarrow j) \Rightarrow (j \rightarrow i)$. The set \mathcal{X} is decomposed as $\mathcal{X} = \mathcal{F} \cup \mathcal{E}_1 \cup \dots \cup \mathcal{E}_k$, where \mathcal{F} is the set of all transient states and \mathcal{E}_a , $a \in \{1, \dots, k\}$, are called ergodic classes. We put two non-transient states into one ergodic class if they are consequents of each other.

Let matrix \mathbb{B}_β have k ergodic classes. Then, there exist k linearly independent left eigenvectors, denoted as $\pi^{(1)}, \dots, \pi^{(k)}$, of matrix \mathbb{B}_β corresponding to eigenvalue 1, called *invariant distributions*. If $\pi^{(a)}\mathbb{B}_\beta = \pi^{(a)}$, for some $a \in \{1, \dots, k\}$, then $\pi_i^{(a)} > 0$ for all $i \in \mathcal{E}_a$, and $\pi_i^{(a)} = 0$ otherwise. Every other π satisfying $\pi\mathbb{B}_\beta = \pi$ is obtained by a convex linear combination of $\{\pi^{(a)} | a \in \{1, \dots, k\}\}$. For a complete reference, see [7, Chapter V, §2]. The set of ergodic classes for matrix \mathbb{B}_β , depends only on set $\{(i, j) | b_{i,j}(\beta) = 0\}$. Since $b_{i,j}(\beta) = 0$ iff $c_{i,j} = 0$ for $i \neq j$ and $b_{i,i}(\beta) > 0$ for $\beta \in (0, \beta_0]$, the structure of ergodic classes does not depend on β (we are not interested in particular values of $\pi_i^{(a)}$ if it is positive). Moreover, if $\pi\mathbb{B}_\beta = \pi$ for some $\beta > 0$, then $\pi\mathbb{C} = 0$ and thus all invariant distributions are independent of β , because $\pi\mathbb{B}_{\beta'} = \pi\mathbb{I} + \beta'\pi\mathbb{C} = \pi\mathbb{I} = \pi$. By this reason, we frequently omit the index β .

3. PERFECTLY SECURE COVER SOURCES UNDER MUTUALLY INDEPENDENT EMBEDDING OPERATION

In this section, we let matrix \mathbb{B} represent an arbitrary MI embedding with k ergodic classes \mathcal{E}_a and invariant distributions $\pi^{(a)}$, $a \in \{1, \dots, k\}$. The following example describes a construction of perfectly secure cover models w.r.t. \mathbb{B} .

Example 2 [Perfectly secure cover models] Let $P^{(2)}$ be a probability distribution on 2-element cover objects defined as $P^{(2)}(X_1^2 = (i, j)) = \pi_i^{(a)}\pi_j^{(b)}$ for some $a, b \in \{1, \dots, k\}$. Then $P^{(2)}$ is a perfectly secure cover model w.r.t. \mathbb{B} . This is because

$$\begin{aligned} Q_\beta^{(2)}(Y_1^2 = (i, j)) &= \left(\sum_i b_{i,i} P(X_1 = i) \right) \left(\sum_j b_{j,j} P(X_2 = j) \right) \\ &= (\pi^{(a)}\mathbb{B})_i (\pi^{(b)}\mathbb{B})_j = \pi_i^{(a)}\pi_j^{(b)} = P^{(2)}(X_1^2 = (i, j)), \end{aligned}$$

and thus both distributions $P^{(2)}$, and $Q_\beta^{(2)}$ are identical, which implies perfect security. Since this construction does not depend on the particular choice of $a, b \in \{1, \dots, k\}$, we can create k^2 perfectly secure cover models w.r.t. \mathbb{B} . The probability distributions $P^{(2)}$ obtained from this construction are linearly independent and form a k^2 -dimensional linear vector space. By a similar construction, we can construct k^n n -element linearly independent perfectly secure cover models w.r.t. \mathbb{B} .

We next show that there are no other linearly independent perfectly secure cover models w.r.t. \mathbb{B} .

Theorem 3 [Mutually independent embedding] There are exactly k^n linearly independent perfectly secure probability distributions P on n -element covers. Every perfectly secure probability distribution P w.r.t. \mathbb{B} can be obtained by a convex linear combination of k^n linearly independent perfectly secure distributions described in Example 2.

Proof It is sufficient to prove that there cannot be more than k^n linearly independent perfectly secure probability distributions P on n -element covers. We show the proof for $n = 2$ and later present its generalization.

We define the following matrices $\mathbb{P} \triangleq (p_{i,j})$, $p_{i,j} = P(X_1^2 = (i, j))$, and $\mathbb{Q} \triangleq (q_{i,j})$, $q_{i,j} = Q_\beta(Y_1^2 = (i, j))$. By definition of MI embedding, we have

$$\begin{aligned} q_{ij} &= \sum_{(v,w) \in \mathcal{X}^2} Q_\beta(Y_1^2 = (i, j) | X_1^2 = (v, w)) P(X_1^2 = (v, w)) \\ &= \sum_{v,w \in \mathcal{X}} b_{vi} b_{wj} p_{vw}. \end{aligned}$$

Define matrix $\mathbb{D} \triangleq (d_{u_1^2, v_1^2})$ of size $N^2 \times N^2$, where $d_{u_1^2, v_1^2} = b_{u_1, v_1} b_{u_2, v_2}$. If \vec{p} is defined as one big row vector of elements $p_{i,j}$ and similarly \vec{q} , then assuming perfect security of cover model w.r.t. \mathbb{B} ($\mathbb{P} = \mathbb{Q}$), we have $\vec{q} = \vec{p}\mathbb{D} = \vec{p}$ and thus \vec{p} is left eigenvector of \mathbb{D} corresponding to 1. Matrix \mathbb{D} is stochastic and thus it is sufficient to show that it has k^2 ergodic classes.

We first show that

$$u_1^2 \xrightarrow{(m)} v_1^2 \Leftrightarrow (u_1 \xrightarrow{(m)} v_1) \text{ and } (u_2 \xrightarrow{(m)} v_2), \quad u_1^2, v_1^2 \in \mathcal{X}^2. \quad (1)$$

By $u_1^2 \xrightarrow{(m)} v_1^2$ we mean that v_1^2 is a consequent of u_1^2 of order m in terms of matrix \mathbb{D} . If $u_1^2 \xrightarrow{(m)} v_1^2$, then there exist $m - 1$ intermediate

states ${}_1 w_1^2, \dots, {}_{m-1} w_1^2$, such that $d_{u_1,1} d_{1,w_2} \cdots d_{m-1,w_1} > 0$. Since $d_{u_1^2,v_1^2} = b_{u_1,v_1} b_{u_2,v_2}$, this implies the existence of both paths $u_i \xrightarrow{(m)} v_i$ of order m , $i = 1, 2$. The converse is true by the same reason.

We show that $\mathcal{E}_a \times \mathcal{E}_b$, $a, b \in \{1, \dots, k\}$ are the only ergodic classes. If $u_1 \xrightarrow{(m_1)} v_1$ and $u_2 \xrightarrow{(m_2)} v_2$, then $u_1^2 \xrightarrow{(m_1+m_2)} v_1^2$ for all $u_1, v_1 \in \mathcal{E}_a$ and $u_2, v_2 \in \mathcal{E}_b$, because the path from u_i to v_i can be arbitrarily extended by adding self loops of type $j \rightarrow j$ since all diagonal terms $b_{j,j}$ are positive and thus by (1) we have $u_1^2 \xrightarrow{(m_1+m_2)} v_1^2$. Finally by $u_1, v_1 \in \mathcal{E}_a$ and $u_2, v_2 \in \mathcal{E}_b$, $v_i \rightarrow u_i$ and by the same argument $v_1^2 \rightarrow u_1^2$, and therefore $\mathcal{E}_a \times \mathcal{E}_b$ are ergodic classes. Any other state $u_1^2 \in \mathcal{E}_a \times \mathcal{F} \cup \mathcal{F} \times \mathcal{E}_a \cup \mathcal{F} \times \mathcal{F}$ must be transient w.r.t. \mathbb{D} , otherwise by (1) we obtain contradiction with $u_i \in \mathcal{F}$ for some i .

This proof can be generalized for $n \geq 3$ by proper definition of matrices \mathbb{P} , \mathbb{Q} , and \mathbb{D} . In general, matrix \mathbb{D} has size $N^n \times N^n$. By similar construction we obtain k^n ergodic classes of generalized matrix \mathbb{D} , however we know k^n linearly independent distributions. ■

4. PERFECT SECURITY AND FISHER INFORMATION

It is well known ([2, Sec. 2]) that the KL divergence imposes a bound on the performance of the best possible detector. For small β , the leading term in the Taylor expansion of the KL divergence is the quadratic term with a constant equal to one half of the Fisher information w.r.t. β , $f(0) = 2\partial^2 d(\beta)/\partial\beta^2|_{\beta=0}$. If for some stego-system $d(\beta) = 0$ for $\beta \in [0, \beta_0]$, then $f(0) = 0$ from the Taylor expansion. Even though the opposite does not hold in general, in this section we prove that for MI embedding zero Fisher information implies perfect security. In other words, a stego-system with MI embedding is perfectly secure for $\beta \in [0, \beta_0]$ if and only if the Fisher information w.r.t. the parameter β is zero, $f(0) = 0$. This provides us with a simpler condition for verifying perfect security than the KL divergence. Moreover, the Fisher information is a fundamental quantity that bounds the variance of minimum variance estimators of β (quantitative steganalyzers).

We start by reformulating the condition $f(0) = 0$.

Proposition 4 Let P, Q_β be probability distributions of cover, stego objects with n elements embedded with parameter β , respectively, then the Fisher information is zero if and only if the so called FI-condition is satisfied

$$\forall y_1^n \in \mathcal{X}^n \quad \left(P(X_1^n = y_1^n) > 0 \right) \Rightarrow \left(\frac{d}{d\beta} Q_\beta(y_1^n) \Big|_{\beta=0} = 0 \right). \quad (2)$$

Proof The second derivative of $d(\beta)$ at β , $d''(\beta)$, can be written as

$$f(\beta) = - \sum_{y_1^n \in \mathcal{X}^n} P(y_1^n) \left(\frac{Q_\beta''(y_1^n)}{Q_\beta(y_1^n)} - \left(\frac{Q_\beta'(y_1^n)}{Q_\beta(y_1^n)} \right)^2 \right), \quad (3)$$

where $Q_\beta'(y_1^n) = \frac{\partial}{\partial\beta} Q_\beta(y_1^n)$. By $P(y_1^n) = Q_{\beta=0}(y_1^n)$, the first term in the bracket in (3) sums to zero at $\beta = 0$, and thus $f(0)$ is zero iff $Q_\beta'(y_1^n)|_{\beta=0} = 0$ is zero for all $y_1^n \in \mathcal{X}^n$ for which $P^{(n)}(y_1^n) > 0$ as was to be proved. Here, we assume the KL divergence $d(\beta)$ to be continuous w.r.t. β which is valid by the construction of the matrix \mathbb{B} . ■

In the next theorem, we show that the FI condition (2) is equivalent with perfect security for stego-systems with MI embedding. Besides providing a simpler condition of perfect security this result plays a key role in proving the square root law of steganographic capacity for covers modeled as Markov chains [1].

Theorem 5 [Fisher information condition] There are exactly k^n linearly independent probability distributions P on n -element covers satisfying the FI condition (2). These distributions are perfectly secure w.r.t. \mathbb{B} . Every other probability distribution P satisfying (2) can be obtained by a convex linear combination of k^n linearly independent perfectly secure distributions.

Proof From Example 2, we know k^n linearly independent perfectly secure distributions. By Taylor expansion of $d(\beta)$, these distributions satisfy the FI condition, because $d(\beta) = 0 \Rightarrow f(0) = 0$. It is sufficient to show that there cannot be more linearly independent distributions satisfying the FI condition.

Similarly as in the previous proof, we reformulate the theorem as eigenvector problem and use ergodic class theory to give the exact number of left eigenvectors corresponding to 1. Again, we present the proof for the case $n = 2$ and then show how to generalize it.

If P satisfies (2), then the linear term in the Taylor expansion of $Q_\beta(y_1^2)$ w.r.t. β is zero. By the independence property, $(Q(y_1^n|x_i^n)) = \prod_{i=1}^n Q(y_i|x_i)$, and the form of matrix \mathbb{B} ($\mathbb{B}_\beta = \mathbb{I} + \beta\mathbb{C}$), condition (2) has the following form

$$\begin{aligned} \frac{dQ_\beta(y_1^2)}{d\beta} \Big|_{\beta=0} &= \lim_{\beta \rightarrow 0} \sum_{x_1^2 \in \mathcal{X}^2} P(x_1^2) \frac{d}{d\beta} \prod_{i=1}^2 Q_\beta(y_i|x_i) \\ &= \sum_{x_1 \in \mathcal{X}} c_{x_1,y_1} P(x_1, y_2) + \sum_{x_2 \in \mathcal{X}} c_{x_2,y_2} P(y_1, x_2) = 0. \end{aligned} \quad (4)$$

We define matrix $\mathbb{P} \triangleq (p_{i,j})$ as $p_{i,j} = P(X_1^2 = (i, j))$ and represent it as a row vector \vec{p} . If we define matrix $\mathbb{D} \triangleq (q_{u_1^2,v_1^2})$ of size $N^2 \times N^2$ as

$$d_{u_1^2,v_1^2} = \begin{cases} c_{u_1,v_1} & \text{if } u_1 \neq v_1 \text{ and } u_2 = v_2 \\ c_{u_2,v_2} & \text{if } u_1 = v_1 \text{ and } u_2 \neq v_2 \\ 0 & \text{otherwise,} \end{cases} \quad (5)$$

and diagonal matrix $\mathbb{G} \triangleq (g_{u_1^2,v_1^2})$ of size $N^2 \times N^2$ as $g_{u_1^2,u_1^2} = -c_{u_1,u_1} - c_{u_2,u_2}$, then equation (4) can be written in a compact form as $\vec{p}\mathbb{D} = \vec{p}\mathbb{G}$. Both matrices \mathbb{D} and \mathbb{G} are non-negative by their definitions.

Let $\mathbb{H} = \mathbb{I} + \gamma(\mathbb{D} - \mathbb{G})$. If we put $\gamma = (\max_{u_1^2 \in \mathcal{X}^2} g_{u_1^2,u_1^2})^{-1}$, then matrix \mathbb{H} is stochastic and $\vec{p}\mathbb{H} = \vec{p}$ iff $\vec{p}\mathbb{D} = \vec{p}\mathbb{G}$ and thus (2) is equivalent with an eigenvalue problem for matrix \mathbb{H} .

First, we observe that for $i \neq j$ $c_{ij} > 0$ iff $h_{(i,a),(j,a)} > 0$ for all $a \in \mathcal{X}$, because by (5) $h_{(i,a),(j,a)} = \gamma d_{(i,a),(j,a)} = \gamma c_{ij}$ (the first case when $u_2 = v_2$). Similarly, for $i \neq j$ $c_{ij} > 0$ iff $h_{(a,i),(a,j)} > 0$ for all $a \in \mathcal{X}$ (the second case when $u_1 = v_1$). This means that $i \rightarrow j$ iff $(i, a) \rightarrow (j, a)$ w.r.t. \mathbb{H} for all $a \in \mathcal{X}$ and similarly $i \rightarrow j$ iff $(a, i) \rightarrow (a, j)$ w.r.t. \mathbb{H} for all $a \in \mathcal{X}$. This can be proved by using the previous statement. By this rule used for a given $u_1^2 \in \mathcal{E}_a \times \mathcal{E}_b$, we obtain $u_1^2 \rightarrow v_1^2$ and $v_1^2 \rightarrow u_1^2$ for all $v_1^2 \in \mathcal{E}_a \times \mathcal{E}_b$ and thus $\mathcal{E}_a \times \mathcal{E}_b$ is an ergodic class w.r.t. \mathbb{H} . We show that there can not be more ergodic classes and thus we have all k^2 of them. If $u_1^2 \in \mathcal{F} \times \mathcal{E}$, then u_1^2 has to be transient w.r.t. \mathbb{H} , otherwise we will obtain contradiction with $u_1 \in \mathcal{F}$. This is because the only consequents

of order 1 are of type $(i, a) \rightarrow (j, a)$ or $(a, i) \rightarrow (a, j)$, therefore if $u_1^2 \in \mathcal{F} \times \mathcal{E}$, we choose $v_1^2 \in \mathcal{X} \times \mathcal{E}$, such that $v_1 \not\rightarrow u_1$ (u_1 is transient and thus such v_1 must exist). State u_1^2 must be transient otherwise $u_1^2 \leftrightarrow v_1^2$ implies $u_1 \leftrightarrow v_1$ which results in contradiction with $v_1 \not\rightarrow u_1$. Similarly for $u_1^2 \in \mathcal{E} \times \mathcal{F} \cup \mathcal{F} \times \mathcal{F}$.

This proof can be generalized for $n \geq 3$ by assuming larger matrices $\mathbb{P}, \mathbb{D}, \mathbb{G}$, and \mathbb{H} , obtaining exactly k^n linearly independent perfectly secure distributions satisfying the FI condition. ■

By the proof of both theorems, the set of all possible perfectly secure cover models w.r.t. \mathbb{B} is a linear vector space. By Theorem 2.1 from [7, Chapter V, page 175] and by the construction from Example 2, this space is generated by k^n basis vectors, where each basis vector is obtained from some invariant distribution of \mathbb{B} . By Theorem 2.1 from [7, Chapter V, page 175], ergodic classes \mathcal{E}_a , $a \in \{1, \dots, k\}$, of the stochastic matrix \mathbb{B} can be obtained from positive elements of the following matrix limit $\mathbb{M} = \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{B}^i$. In other words, if we reorder the set \mathcal{X} so that the states from \mathcal{E}_1 are first, then the states from \mathcal{E}_2 etc., and as the last the set of transient states, \mathcal{F} , then the matrix \mathbb{M} will be block-diagonal, where the a -th block is a positive matrix of size $|\mathcal{E}_a| \times |\mathcal{E}_a|$ for $a \in \{1, \dots, k\}$. By this fact, the invariant distribution belonging to ergodic class \mathcal{E}_a , $\pi^{(a)}$, is zero except for indices \mathcal{E}_a , i.e., $\pi_i^{(a)} > 0$ if $i \in \mathcal{E}_a$ and zero otherwise.

For the F5 embedding algorithm [3], the set of states $\mathcal{X} = \{-1024, \dots, 1024\}$. By the nature of the embedding changes (flip towards 0), there is only one ergodic set $\mathcal{E}_1 = \{0\}$ and $\mathcal{F} = \mathcal{X} \setminus \{0\}$. Thus, there is only one invariant distribution, which is singular, $\pi_0 = 1$ and zero otherwise. By the form of the invariant distribution π , no message can be embedded in such covers.

For the case of LSB embedding over $\mathcal{X} = \{0, \dots, 255\}$, we have $\mathcal{E}_a = \{2a, 2a+1\}$ for $a \in \{0, \dots, 127\}$, $\mathcal{F} = \emptyset$ and $\pi_{2a}^{(a)} = \pi_{2a+1}^{(a)} = \frac{1}{2}$ and zero otherwise. This leads to the well known fact that LSB embedding evens out the histogram bins. Thus, sources realized as a sequence of mutually independent random variables with such a distribution are the only possible perfectly secure sources w.r.t. LSB embedding. Figure 1 shows examples of matrices \mathbb{B} and ergodic classes of several known algorithms with MI embedding operation.

5. CONCLUSION

The theory of ergodic classes, originally developed for Markov chains, allows us to gain insight into perfectly secure stego-systems with mutually independent embedding operations. We knew that some i.i.d. cover sources are perfectly secure w.r.t. some embedding algorithm, however it was not immediately clear, if this set of perfectly secure sources cannot be larger. An important corollary can be obtained from both theorems if we constrain ourselves to stationary cover sources. In this case, for a given embedding algorithm we have exactly k (instead of k^n) perfectly secure stationary cover sources, which are i.i.d. sources with some invariant distribution.

Perfect security of stego-systems with MI embedding is completely captured using Fisher information formulated in Section 4 as the FI condition. This result not only provides a simpler and equivalent condition for perfect security, but it finds applications in theoretical steganalysis. For example, in [8], Ker et al. introduced the problem of steganographic capacity w.r.t. the number of cover elements. While the capacity of noisy channels is proportional to n , the authors proposed and practically verified that the capacity of steganographic channels is only proportional to \sqrt{n} if the sender doesn't know the

source completely. This conjecture was proved in [1] for the case of Markov cover sources under the MI embedding operation. There, the FI condition was used to eliminate all perfectly secure stego-systems from analysis, because it is known that capacity of such systems is linear in n .

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