DATA COMPRESSION FOR INFERENCE TASKS IN WIRELESS SENSOR NETWORKS

BY

MO CHEN

BS, Shandong University, China, 1998 MS, Shandong University, China, 2001

DISSERTATION

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Mark L. Fowler,	Department of Electrical and Computer Engineering,
	Binghamton University
Xiaohua (Edward) Li	Department of Electrical and Computer Engineering,
	Binghamton University
N. Eva Wu	Department of Electrical and Computer Engineering,
	Binghamton University
Harold W. Lewis III,	Department of Industry Engineering and System Science,
	Binghamton University

Abstract

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Mo Chen

Chair: Mark L. Fowler

In order for wireless sensor networks to exploit signal, signal data must be collected at a multitude of sensors and must be shared among the sensors. The vast sharing of signals among the sensors contradicts the requirements (energy efficiency, low latency and high accuracy) of wireless networked sensor. Although many approaches have been proposed in the past (routing, sleep modes, low-power electronics, etc.), a new aspect is proposed here: using data compression methods as a tool for accomplishing the optimal trade-off between rate, energy, and accuracy in a sensor network. The ability of data compression to provide energy efficiency rests on the favorable trade between computational energy and transmission energy recently recognized in the literature.

Because a primary task of multi-sensor systems is to make statistical inferences based on the data collected and shared throughout the sensor system, in the viewpoint of rate, energy, and accuracy, it is important to design data compression methods that enable rapid and low-energy consumption sharing while causing only minimal degradation of the quantity of these inferences. By recognizing that MSE-based compression algorithms are not appropriate for such tasks, this dissertation stresses the development of distortion measures that effectively capture the impact of compression on the accuracy of the inferences. Furthermore, sensor systems generally have

multiple inference tasks to accomplish simultaneously, and these multiple inferences generally have conflicting requirements on compression and finding the right way to balance these conflicts is crucial, this dissertation develops theory and algorithms that will allow the optimal trade-off between these conflicting goals. On the other hand, these multiple tasks may occur sequentially (and a step in the sequence could require simultaneous inferences). A task-embedded compression method is developed to compress and send data in a sequential manner that allows optimal attainment of the sequential tasks.

The contributions of this dissertation include: (i) strengthening data compression as a powerful tool for achieving optimal tradeoff among energy, latency and accuracy in sensor networks; (ii) developing new fundamental framework for compression in sensor networks that recognize the inferential characteristics of sensor networks; (iii) developing a significant framework for the "compression for multiple inferences" area that addresses multiple simultaneous and sequential inferences (no results in the literature currently address this important issue); The results of this dissertation will provide the engineer with a systematic means for addressing Rate-Energy-Accuracy issues across the spectrum of sensor network types – from networks of myriad microsensors to networks of several macrosensors.

To my parents and Dr. Mark. L. Fowler

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TABLE OF CONTENTS

COMMITTEE PAGE	iii
BSTRACT	. iv
DEDICATION	vi
CHNOWLEGEMENTS	vii
IST OF TABLES	. X
IST OF FIGURES	. xi
IST OF APPENDICES	xv

CHAPTERS

1	Introduction 1.1 Wireless Sensor Networks 1.2 Importance of Data Compression for Energy-Efficiency in Wireless Sensor Networks 1.3 Compression for Inference Tasks in Wireless Sensor Networks	1 1 . 2 . 9
2	Overview of Compression System for Inference Tasks 2.1 Estimation 2.2 Detection 2 Detection 2 Data Compression and Decompression	13 14 18 20
	2.3 Data Compression and Decompression 2.3.1 Coding 2.3.2 Transform 2.3.3 Quantization 2.3.4 Bit Allocation 2.4 Compression for inference tasks	20 22 23 25 29 30
3	Data Compression for Parameter Estimation and Detection 3.1 Fisher-Information-Based Data Compression for Parameter Estimation 3.1.1 Overview of Previous Work 3.1.2 Fisher Information for Compression 3.1.3 Compressing to Maximize Fisher Information 3.1.4 Example Applications 3.1.4.1 Compression for TDOA Estimation 3.1.4.2 Compression for FDOA Estimation	. 35 35 36 38 . 48 50 55

3.1.4.3 Compression for DOA Estimation623.1.4.4 Compression for Signal Estimation653.2 Chernoff-Distance-Based Data Compression for Detection683.2.1 Overview of Previous Work683.2.2 Compressing to Maximize Chernoff Distance72
 4 Data Compression for Simultaneous Tasks of Multiple Inference Quantities . 76 4.1 Data Compression for Simultaneous Multiple Parameter Estimation
without Detection
4.1.1 Application to Simultaneous Estimation
4.1.1.1 Joint TDOA and DOA Estimation
4.1.1.2 Simulations for Joint TDUA and DUA Estimation
4.1.1.5 JOINT IDUA and FDUA Estimation
4.1.2 Mounted Distortion Measures to include On-Diagonals 100
4.1.5 Simulations of compression using paramet auxiliary 5111
4 2 Data Compression for Parameter Estimation with Detection 110
4.2 Data compression for rarameter Estimation with Detection
5 Data Compression for Sequential Tasks of Multiple Inference Quantities 114
5.1 Simple Detection-Then-TDOA Problem
5.1.1 An Algorithm for Sequential Detect-Then-TDOA
5.1.2 Simulation Scenario
5.2 General Sequential data compression algorithm
6 Conclusion
APPENDICES
BIBLIOGRAHPY

LIST OF TABLES

Table

5.1 Sequential Example, Rough Multiple Estimate - Then- Refine Estimation .	126
5.2 Sequential Example, Detection-Then - Estimation	126
B.1 Fitness of the approximation model with the true mode	141

LIST OF FIGURES

Figure

1.1 Non-distributed compression improves network lifespan when using direct transmission
1.2 Result of showing improvement using compression with a simple routing scheme
2.1 System of Compression for inference tasks, \mathbf{x}_1 and \mathbf{x}_2 are digitized sensed measurements at sensor S_1 and S_2 , and $\hat{\mathbf{x}}_1$ is the decompressed data at sensor S_2 14
2.2 The elements of compression-decompression process
2.3 Graphical representation of scalar quantization
2.4 Embedded scalar quantizers Q_0 , Q_1 , and Q_2 , of rates $R = 1, 2$, and 3 bits/samples 28
2.5 Uniform scalar quantizer with deadzone
2.6 Framework of the compression method
3.1 Classical set-up for compression in a distributed sensor system
3.2 Compression processing, the data vectors, and their corresponding Fisher information
3.3 The spectrum of a typical FM signal used in the simulations
3.4 TDOA accuracy vs SNR of pre-compressed sensor S ₁ signal for a CR of 4:1; the SNR of the sensor S ₂ signal was 40 dB
3.5 TDOA accuracy vs SNR of pre-compressed sensor S1 signal for a CR of 8:1; the SNR of the sensor S2 signal was 40 dB
3.6 TDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 4:1; the SNR of the sensor S_2 signal was set equal to SNR ₁
3.7 TDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 8:1; the SNR of the sensor S_2 signal was set equal to SNR ₁

3.8 Reconstruction accuracy vs SNR of pre-compressed sensor S_1 signal for a compression ratio of 4:1
3.9 Reconstruction accuracy vs SNR of pre-compressed sensor S_1 signal for a compression ratio of 8:1
3.10 FDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 4:1; the SNR of the sensor S_2 signal was 40 dB
3.11 FDOA accuracy vs SNR of pre-compressed sensor <i>S</i> ₁ signal for a CR of 8:1; the SNR of the sensor <i>S</i> ₂ signal was 40 dB
3.12 DOA accuracy vs SNR_1 pf pre-compressed sensor S_1 signal for a CR of 4:1; the SNR of the sensor S_2 signal was equal to SNR_2
3.13 DOA accuracy vs SNR_1 pf pre-compressed sensor S_1 signal for a CR of 8:1; the SNR of the sensor S_2 signal was equal to SNR_2
4.1 The spectrum of a typical FM signal used in the simulations
4.2 Trade-off between TDOA and DOA accuracies as α is varied for $SNR_1 = 15$ dB & $SNR_2 = 15$ dB and compression ratio (a) 3.5:1 and (b) 6:1
4.3 Effect of compression ratio on (a) TDOA and (b) DOA performance
4.4 Trade-off between TDOA and FDOA accuracies as α is varied for compression ratio 3:1 and $SNR_1 = 15 \text{ dB} \& SNR_2 = 15 \text{ dB} \dots 98$
4.5 Effect of compression ratio on (a) TDOA and (b) FDOA performance. A comparison is also made to the case of simply sending less data ("Length Reduced") rather than compressing the data
4.6 Comparison between the determinant optimization method ('Area') and weighted trace method ('Premeter') and MSE for compression ratio 3:1 and $SNR_1 = 15 \text{ dB } \& SNR_2 = 15 \text{ dB} \dots 100$
4.7 Geometry Adaptive TDOA/FDOA System Scheme
4.8 Effect of compression ratio on CEP 105
4.9 Our compression framework using parallel auxiliary STFT processing to evaluate the FIM elements
4.10 (a) Geometry among the two pair of sensors and emitter; (b) The comparison of difference compression algorithms on the trace of error covariance of emitter location

4.11 (a) Geometry among the two pair of sensors and emitter; (b) The comparison of difference compression algorithms on the trace of error covariance of emitter location
5.1 Sequential Tasks with the Initial Trade-Off
5.2 Sequential Tasks with a New Trade-Off 118
5.3 Trade-off for fixed task resources
5.4 Conceptual illustration of the trade-off accomplished via choice of the β parameter. where $SNR_1 = SNR_2 = 15$ dB
5.5 Simulation results illustrating the achieved trade-offs
B.1 Comparison of characteristic function between true model and approximation model
B.2 Comparison of value $I_t(\xi_n(\theta), \Delta_n, \sigma)$ and $I_a(\Delta, \sigma)$ under different values
of $\xi_n(\theta)$ and Δ_n^2 when (a) PSNR =10 dB, (b) PSNR=15 dB
C.1 Three-level wavelet decomposition
C.2 The input and reference images 157
C.3 Reconstructed image from (a) the standard JPEG or EZW compression, (b) the modified JPEG compression with the weight matrix $\sqrt{i^2 + j^2}$ and all-ones mask, (c) the modified JPEG compression with the zigzag weight matrix and the example mask, and (d) the modified EZW compression
C.4 Normalized correlation outputs of the reference and the input image processed (a) without compression, (b) with JPEG or EZW compression, (c) with the modified JPEG, and (d) with the modified EZW compression
C.5 Compression ratio versus quality factor
C.6 Reconstructed image and correlation output when $q=12$
C.7 Reconstructed image and correlation output when $q=35$
C.8 (a) the noisy input image; (b) normalized correlation outputs of the reference and (a) with the modified JPEG; (c) normalized correlation outputs of the reference and (a) with the modified EZW
E.1 Quality of numerical evaluation of FIM via complex filter bank
E.2 Quality of numerical evaluation of FIM via STFT
F.1 Example of a four-channel, full-tree cascaded analysis filter bank

F.2 Example of a four-channel, full-tree cascaded synthesis filter bank
F.3 Equivalent filter bank with decimators moved
F.4 Equivalent filter bank with decimators moved and blocks combined
F.5 Ideal filters for the various stages (w/ decimation moved to end)
F.6 Shows pairs of filters that appear in the four different cascades that form the four channels
F.7 Four Channels' frequency response (w/ decimation moved to end)
F.8 The four channels and how they split up the input DTFT 172
F.9 Channel 1 after decimation - shows all the streched spectral replicas that arise from decimation
F.10 Channel 2 after decimation - shows all the stretched spectral replicas that arise from decimation
F.11 Channel 4 after decimation - shows all the stretched spectral replicas that arise from decimation
F.12 Channel 3 after decimation - shows all the stretched spectral replicas that arise from decimation
F.13 Channel output DTFT's after decimation - only shows standard $-\pi$ to π range 174
F.14 Relationships between the filters for real PR filter banks
F.15 Movement of real filter zeros under $z \rightarrow z^{-1}$
F.15 Movement of real filter zeros under $z \rightarrow z^{-1}$
 F.15 Movement of real filter zeros under z → z⁻¹
 F.15 Movement of real filter zeros under z → z⁻¹
 F.15 Movement of real filter zeros under z → z⁻¹

LIST OF APPENDICES

APPENDIX

A	Separation Theorems of the Fisher Information and the Chernoff Distance	30
В	Fisher Information and Chernoff Distance of Scalar Quantized Data 12	34
С	Modification of Standard Image Compression Methods for Pattern Recognition 14	44
D	Symmetric Indexing	51
E	Correlated TDOA/FDOA Estimates and FIM for TDOA/FDOA of a Chirp Signal	52
F	Complex PR Filter Banks	57

CHAPTER 1 Introduction

1.1 Wireless Sensor Network

Advances in sensor and communication technology have focused interest on using wireless sensor networks, which are formed by a set of small unterthered sensor devices that are deployed in an ad hoc fashion to cooperate on sensing a physical phenomenon, making the inferences, and transmitting the data [12]-[20]. Typically, each individual sensor can sense in multiple modalities but has limited communication and computation capabilities. Wireless sensor networks hold the promise of revolutionizing sensing in a wide range of application domains because of their reliability, accuracy, flexibility, cost-effectiveness, and ease of deployment.

Wireless sensor networks share many of the challenges of traditional wireless networks, including limited energy available to each node and bandwidth-limited, error-prone channels. Among these challenges, energy is typically more limited in wireless sensor networks than in other wireless networks because of the nature of the sensing devices and the difficulties in recharging their batteries. Usually, the following three metrics are used to evaluate the design of any wireless sensor networks [18]:

- Energy efficiency/system lifetime: As sensor nodes are battery-operated, the design of the wireless sensor network must be energy-efficient to maximize system lifetime.
- Latency: The phenomena of interests or inference results must be reported within a given delay.

3. Accuracy: Obtaining accurate information is the primary objective.

There are also constraints on fault tolerance and scalability [18] but we don't address those here. Any good design of wireless sensor networks must be adaptive to obtain the optimal tradeoff over metrics assessing energy efficiency, communication latency, and accuracy of the conveyed information; for example, a well designed network achieves the desired accuracy and delay while optimizing the sensor energy usage, or maximizes the inference accuracies given the desired energy expenditure and small latency.

1.2 Importance of Data Compression for Energy-Efficiency in Wireless Sensor Networks

Energy efficiency in wireless sensor networks has principally been addressed through routing protocols, sleeping strategies, low-power architectures, and energy-efficient modulation schemes. Accuracy is generally controlled through optimal processing strategies as well as using accurate sensors deployed in optimal ways. Latency and channel capacity issues in sensor networks can be addressed through routing strategies and data compression [16],[17]. It is very important to understand the interplay between the compression method and routing. In the following, by investigating a well recognized routing scheme, we demonstrate that data compression can bring more energy efficiency to a network than does recently proposed combinations of routing and data aggregation.

There can be many different scenarios for sensor networks; here we focus on the so-called "reach-back" issue: communicating the data collected within the network back to a single information sink (e.g., base station, central command, etc.) with minimal latency and energy use. Energy efficiency in reach back has been previously addressed by many researchers including [20], where energy efficiency was measured in terms of network lifespan. A related study has

been carried out in [17] to show the need for compression to address the latency issue. The usefulness of data compression for energy efficiency is less clear.

Before discussing our results, we put our study into perspective with recent related results in [17] and [20]. The results in [20] address energy efficiency for reach-back through use of a combination of routing and data fusion/aggregation (called "LEACH"). By using data fusion/aggregation to combine two or more collected data sets that become correlated during transmission of the data through the network towards the sink, LEACH significantly reduces the overall data needed to be transferred and increases network lifespan. However, it is not clear in these papers how data fusion/aggregation can be relied on in *general* in a sensor network. In particular, in [20] the data from sensors grouped into clusters get fused through processing, where a stated assumption is that this fusion specifically implements beamforming; therefore, data aggregation is possible in a general sensor network setting. However, it is not clear that fusion/aggregation is possible in a general framework using data compression rather than fusion/aggregation that would give similar gains in energy efficiency.

The results in [17] don't consider reach-back but rather the task of conveying the network's total collected information to each and every sensor node. The results in [17] establish fundamental information theoretic limits on the rate of information transferal through the network and show that data compression is needed to transfer the data without latency under a channel rate constraint. But for us, the key idea established in [17] is the effectiveness of combining classical source codes with routing algorithms and that this is competitive with distributed compression methods such as in [16], which remove common information between two nodes without sharing any data between them.

For ease of comparing results, we use the same radio model used in [20] the radio dissipates 50 nJ/bit in the transmitter circuitry, 50 nJ/bit in the receiver circuitry, and 100 pJ/bit/m2 for the transmitter amplifier. Because we aren't using a *specific* compression algorithm in this part of our study it is hard to specify how much energy is spent compressing the data, so we use the same energy cost that is used for data fusion/ aggregation via beamforming in LEACH, namely 5 nJ/bit/message. To compare our methods with LEACH we ran tests using the following scenario: 100 sensors, randomly placed uniformly inside a 50m×50m square of real estate. LEACH randomly selects 5% of its nodes as cluster heads; data from all the nodes in a cluster are beamformed together (data aggregation). This gives LEACH an inherent "compression ratio" of 20:1 since at each cluster head, 20 signals get beamformed into one. One of the key published conclusions for LEACH is that sending *directly* to the sink is inferior to LEACH; however, this really is an unfair comparison since the direct transmission method did not use any form of compression in [20]. Therefore we performed a simple simulation to show that using general compression without any routing provides better network lifetime- thus, it is LEACH's beamforming-achieved compression, not the routing protocol, that achieves the energy efficiency. Of course, it is granted that the routing does have the advantage that it uniformly spreads sensor deaths over the network.

A. Direct Transmission with Non-Distributed Compression

We simulated LEACH as well as direct transmission with compression, with the later using a compression of 6:1 and 10:1. As in [20], each "round" of network transmission consisted of each node receiving 2000 bits of data and the network transferring the data to the sink. In LEACH, cluster heads are randomly selected on each round and the remaining nodes are assigned to clusters. Each cluster head receives 2000 bits from each of its cluster nodes and beamforms them

into a single 2000 bit signal, which is then transmitted from the cluster head to the sink. Alternatively, direct-with-compression compresses the 2000 bits received at each node and then transmits the resulting bits to the source. The lifetime of the network is assessed by noting the number of nodes still alive at each round, where a live node is taken to be a node that has energy remaining. As mentioned above, comparable amounts of computational energy are assumed to be spent for beamforming and compression.

Our study shows that if a compression ratio of 6:1 is achievable at each sensor prior to sending the data directly to the sink, then the time it takes for 50% of the nodes to die is comparable to LEACH, as seen in Figure 1.1. Obviously, higher compression ratios further improve the direct w/compression curve, as is shown in Figure 1.1 for the case of 10:1, where the direct-with-compression now clearly outperforms LEACH.

This result is important because previous results [20] indicated that direct transmission was to be avoided and that special routing schemes were the answer to the energy efficiency problem. In our results we see that it is not the routing in LEACH that makes the difference, it is the beamforming-achieved compression. In understanding our result it is important to keep in mind that the direct method has no energy cost for reception since sensor nodes don't receive any transmissions; that compensates for the excess compression ratio that LEACH is assumed to have here (20:1 vs. 6:1 or 10:1).

However, it should also be mentioned that [20] points out that a problem with the direct method is that the death of nodes begins with the nodes farthest from the sink and sweeps through the network towards the sink– this is generally undesirable and one nice feature of LEACH is that node deaths are uniformly distributed. Clearly direct-with-compression isn't directly applicable but it does point out the importance of data compression.



Figure 1.1: Non-distributed compression improves network lifespan when using direct transmission.

B. Direct Transmission with Compression Exploiting Spatial Redundancy

The direct-transmission method discussed above does not exploit any redundancy between signals received at closely located sensor nodes. As a *simple* demonstration of the effectiveness of exploiting spatial redundancy we ran simulations to characterize the effect of exploiting this spatial redundancy. The results are shown in Figure 1.2 where we have simulated the effect of signals within a radius of 10m of a randomly selected set of primary nodes (making up 10% of the total number of nodes) as having virtually the same information content. This is more like LEACH but with beamforming-based aggregation replaced by general data compression. The spatially similar data is compressed and then the remaining data sets are compressed using non-spatial methods having a compression ratio of 6:1, after which all compressed data is sent to the information sink using direct transmission. The results in Figure 1.2 show the potential of exploiting this spatial redundancy through routing and local compression (rather than distributed compression). By randomly rotating which sensors are used as the central compression sites the "sweep-of-death" for direct transmission is eliminated as it is in LEACH. The reason that this scheme far outperforms LEACH even though its CR is only 6:1 compared to LEACH's 20:1 is



Figure 1.2: Result of showing improvement using compression with a simple routing scheme.

that we have also reaped the benefit of compression as each sensor sends its data to the cluster head.

It is important to realize that we are not proposing that either of these two direct-transmission approaches should be considered as viable real-world methods; rather we are simply using them to illustrate that data compression plays an important role not only in controlling latency in sensor networks (as in [20]) but plays an important role in addressing network energy efficiency.

From above observations, we thereby assert that data compression – normally viewed as a primary tool for the capacity/latency vs. accuracy trade-off – is a natural tool for achieving the desired trade-off between energy, accuracy, and latency in sensor networks. This expanded role of compression arises as follows: Transmission of data is one of the most energy-expensive tasks a node undertakes – thus, using data compression to reduce the number of bits sent reduces energy expended for transmission because the transmission energy is linearly proportional to the number of data to transmit [20][85]. However, compression requires computation, which also expends energy. Fortunately, trading computation for transmission has potential to save energy: a recent paper [15] asserts that typically on the order of 3000 instructions can be executed for the

energy cost required to transmit one bit over a distance of 100 m by radio – thus, the trade favors spending computation to reduce the amount of transmission needed.

C. Rate-Energy-Accuracy function for Sensor Network Compression

Classical data compression theory relies on tradeoffs between rate (R) and distortion (D) in terms of a *R-D* function. Rate is usually measured in terms of bits/sample and distortion is often measured as a mean square error between the original and reconstructed signal. In the classical view, rate impacts latency and distortion impacts the accuracy of the signal reconstruction. As we explored above, in sensor networks the rate can also impact energy efficiency. Thus, for sensor networks we propose the use of a 3-D extended version of the *R-D* function: the Rate-Energy-Accuracy (R-E-A) function. The rate axis equals the length of compressed data and the accuracy axis is proportional to distortion caused by compression and reflects the effect of the compression energy and transmission energy) needed to move the compressed data to the desired destination. Clearly, decreasing the rate decreases the amount of transmission energy spent and the duration of transmission time but a decreased rate comes at the expense of computational energy and time from compression algorithms. A simple characterization of this is

$$\Delta E(R) = E_C(R) + E_T(R) , \qquad (1.1a)$$

$$\Delta T(R) = T_C(R) + T_T(R) , \qquad (1.1b)$$

where $\Delta E(R)$ and $\Delta T(R)$ are the changes in energy and time respectively due to compression to a rate of R, $E_C(R)$ and $T_c(R)$ are the computational energy and time used to compress to R, similarly, $E_T(R)$ and $T_T(R)$ are the energy and time needed to transmit at the rate R, respectively. To maintain certain accuracy (quality of information), if R decreases, both $E_C(R)$ and $T_c(R)$ increase (better compression requires more computation), in the meantime, $E_T(R)$ and $T_T(R)$ decreases (more compression requires less transmission). Clearly, these measures depend on the computational efficiency of the compression algorithm, the energy efficiency of the computational architecture, and the energy efficiency of the transmission hardware. If we can specify a desired operating point in R-E-A space, then we develop compression algorithms (as well as low-power computing & transmitting architectures) that achieve it. However, optimizing of such multi-objective functions (1.1) should be based on the trinity of compression algorithm, hardware architecture and transmission, which is a horrible task even if it seems possible in theory. An alternative way to balance the trade-off between $E_T(R)$ and $E_C(R)$ (or $T_T(R)$ and $T_C(R)$) is to apply the idea of separate optimality as in communication [1]. More specifically, for a given rate, we can choose to maximize the accuracy using the compression algorithm which is a efficient as possible (minimizing the computation).

1.3 Compression for Inference Tasks in Wireless Sensor Networks

The primary task of sensor networks is to make statistical inferences based on the data collected throughout the network, it is important to design compression methods that cause minimal degradation of the accuracy of these inferences. Traditionally, the mean-square-error (MSE) is the primary distortion measure to guide the compression algorithms in literature. MSE is a natural choice of the distortion measures for the application whose goal is to reconstruct the data as near to the original data as possible. However, MSE is not able to capture the effect of the compression error on the final use of the data – namely, the making of statistical inferences. For example, if the inference task is estimation, then the accuracy measure should capture the impact of the compression on the estimation accuracy instead of the reconstruction accuracy. Compression with the goal to maximize the inference task accuracy must be found instead. The key to addressing this question is to use distortion measures that accurately reflect the ultimate

performance on the tasks. For estimation tasks the ultimate performance is the variance of the estimation error (at least in the unbiased estimate case) and MSE is only part of what determines the variance. Similarly, for decision tasks the ultimate performance is probability of detection. To design compression algorithms with respect to these performance goals it is essential to have appropriate, tractable metrics that measure the impact of reducing the rate on the inference performance. For estimation we propose using the Fisher Information Matrix (FIM) to provide a guide to how to reduce the rate while minimizing the impact on estimation performance. For decision we propose exploring the use of the Chernoff distance to derive the distortion measures for the compression algorithms [21]-[28], [60], [75].

Data compression for distributed multi-sensor systems previously considered in the literature (See the details in Chapter 3) all hold the view that multiple sensors encode received signals and transfer the coded data directly to a central inference base station. Thus, these prior results are not made in a true sensor network context. For example, they don't consider the energy constraint as mentioned above. Nor do they consider the impact and needs of inter-sensor communication and cooperation. An even more important aspect not considered in these previous results is that sensor networks may have multiple inference tasks to accomplish (either simultaneously or sequentially). Multiple inferences may have conflicting compression requirements and finding the right way to balance these conflicts is crucial. Thus, another of our assertions is that compression for sensor networks must consider the case of multiple inferences.

In a sensor network there are a lot of cases where inference tasks are naturally done sequentially and therefore the sharing of data to complete those tasks can also occur sequentially. As an example of multiple inferences, consider the case where a sensor network is deployed to detect and then locate vehicles. This is a case of multiple sequential inferences where the

compression can be done sequentially as well. For example, at first sensors might share their collected data for the purpose of improved detection. After detecting the presence of a vehicle, data would be shared among sensors to estimate the vehicle's position, direction and velocity. Our proposed viewpoint for a novel approach to such a compression scenario is what we call task-embedded compression: at each task stage, send only that data needed to supplement the already-delivered data for optimal processing for the current task. For example, (i) the data stream that is shared during the detection phase is optimally compressed for detection, then (ii) the data stream that is shared during the estimation phase is the additional data "layer" needed to perform optimal estimate. It should be pointed out that in the last task of this example there might be simultaneous multiple inferences (estimate position and velocity) that may very well have conflicting data compression requirements – thus, we must find ways to compress data that allow the proper trade-off in this conflict. The key tools we bring to bear on this area are: (i) the use of multiple distortion measures that are designed to assess the quality of data subsets for use in the multiple inferences, and (ii) the use of numerical optimization methods to achieve desired trade-offs based on these quality assessments.

This dissertation is organized as follows. In Chapter 2, we review the basic elements of inference tasks (estimation and detection) and the framework of compression (transform coding). In Chapter 3, we first limit our attention to the single inference task case and propose using Fisher information and Chernoff distance to derive the effective distortion measures for estimation and detection task respectively. In Chapter 4, we attack optimization of compression for the *general* simultaneous multi-parameter estimation problem and simultaneous multi-inference tasks (joint detection/estimation). In Chapter 5 we consider a sequence of inference

tasks. Chapter 6 presents conclusions from this work. Some mathematical deduction needed for the body of the dissertation and some supplemental materials are in Appendices.

CHAPTER 2

Overview of Compression System for Inference Tasks

The overall nature of the transportation of sensed data can be vastly different for different scenarios; however all share one common thread: sending data collected in one place to some other place where it is processed. As illustrated in Figure 2.1, for the work we report in this dissertation, we limited our focus on the compression of data received at one sensor node (S_1) , which is then transmitted to a second node (S_2) where it is combined with data received locally at S_2 to obtain some estimates of some desired environmental parameters that is reflected in the data or to make some decision of the presence of anomalies from the data. This scenario is likely to occur as data is routed through a micro-sensor network; it also is common in macro-sensor networks. Our approach can also be applied to the case where the data at multiple sensors must be compressed and transmitted to a central node where estimation is done; We are particularly (although not exclusively) interested in one aspect of cooperative signal processing where it is not possible to perform an inference (detection or estimation) based on a single sensor's data (e.g., centralized detection, source location), because in scenarios where it is possible for each sensor to make an inference it is likely to be more efficient to transmit sensor-local inference results that are then fused (decentralized system). However, it is generally the case that better inferences can be made when the fusion center has access to the raw data rather than sensor-local estimates; thus, the



Figure 2.1: System of Compression for inference tasks, \mathbf{x}_1 and \mathbf{x}_2 are digitized sensed measurements at sensor S_1 and S_2 , and $\hat{\mathbf{x}}_1$ is the decompressed data at sensor S_2 .

proper design of a sensor network should consider the trade-offs between fusing sensorlocal estimates and making centralized estimates based on compressed sensor data.

In the following, each component of the system in the system of Figure 2.1 will be further explained:

2.1 Estimation

Given the received signal as

$$x[n] = \eta[n; \theta] + w[n], \qquad n = -N/2, \dots, N/2 - 1$$
 (2.1)

where x[n] is the result of sensing measurement at the time instant n, $\eta[n; \theta]$ is a deterministic function whose dependency on θ is assumed known; $\theta^T = [\theta_1, \theta_2, ..., \theta_p]$ is the unknown, $p \times 1$ deterministic parameters sought to be estimated, and w[n] is the measurement noise. The measurement x[n] is random due to that w[n] is modeled as random noise and its probability distribution ($p(x[1],...,x[n]|\theta)$) is functionally dependent on the parameter vector θ . The equation (2.1) can be written in vector form as

$$\mathbf{x} = \mathbf{\eta}_{\mathbf{\theta}} + \mathbf{w} \tag{2.2}$$

In judging an estimation problem, it is important to know the maximum estimation accuracy that can be obtained with the measurements. As is well known, the Cramer Rao (CR) bound provides the powerful lower bounds on the estimation accuracy [4][73]. Further, since they are lower bounds for any estimator, their parameter dependences reveal characteristic features of the estimation problem.

If we let $\hat{\theta}(\mathbf{x})$ denote some unbiased estimate of unknown parameter vector $\boldsymbol{\theta}$ based on the measurement of \mathbf{x} and it is assumed that the PDF $p(\mathbf{x} | \boldsymbol{\theta})$ satisfies the "regularity" conditions

$$E\left[\frac{\partial \ln p(\mathbf{x};\boldsymbol{\theta})}{\partial \boldsymbol{\theta}}\right] = 0 \quad \text{for all } \boldsymbol{\theta} , \qquad (2.3)$$

where the expectation is taken with respect to $p(\mathbf{x}; \boldsymbol{\theta})$, then the covariance matrix of estimation error of any unbiased estimator $\hat{\boldsymbol{\theta}}$ satisfies

$$E\left[(\hat{\boldsymbol{\theta}} - \boldsymbol{\theta})(\hat{\boldsymbol{\theta}} - \boldsymbol{\theta})^{T}\right] = \mathbf{C}_{\hat{\boldsymbol{\theta}}} \ge \mathbf{J}^{-1}(\boldsymbol{\theta})$$
(2.4)

where ≥ 0 is interpreted as meaning that the matrix is positive semidefinite. The Fisher information matrix $\mathbf{J}(\mathbf{\theta})$ is given as

$$\left[\mathbf{J}(\mathbf{\theta})\right]_{ij} = -E\left[\frac{\partial^2 \ln p(\mathbf{x};\mathbf{\theta})}{\partial \theta_i \partial \theta_j}\right], \qquad 1 \le i, j \le n,$$
(2.5)

where the derivatives are evaluated at the true value of θ and the expectation is taken with respect to $p(\mathbf{x}; \theta)$. Moreover, for a positive semi-definite matrix the diagonal elements are nonnegative. Hence,

$$\left[\mathbf{C}_{\hat{\boldsymbol{\theta}}} - \mathbf{J}^{-1}(\boldsymbol{\theta})\right]_{ii} \ge 0$$
(2.6)

and therefore

$$\operatorname{var}(\hat{\boldsymbol{\theta}}_{i}) = \left[\mathbf{C}_{\hat{\boldsymbol{\theta}}}\right]_{ii} \ge \left[\mathbf{J}^{-1}(\boldsymbol{\theta})\right]_{ii}$$
(2.7)

Whenever it is desired to estimate $\alpha = g(\theta)$ for g, an r-dimensional function, then,

$$\mathbf{C}_{\hat{\alpha}} - \frac{\partial g(\boldsymbol{\theta})}{\partial \boldsymbol{\theta}} \mathbf{J}^{-1}(\boldsymbol{\theta}) \frac{\partial g(\boldsymbol{\theta})^{T}}{\partial \boldsymbol{\theta}} \ge 0, \qquad (2.8)$$

where

$$\frac{\partial \mathbf{g}(\mathbf{\theta})}{\partial \mathbf{\theta}} = \begin{bmatrix} \frac{\partial g_1(\theta)}{\partial \theta_1} & \frac{\partial g_1(\theta)}{\partial \theta_2} & \cdots & \frac{\partial g_1(\theta)}{\partial \theta_p} \\ \frac{\partial g_2(\theta)}{\partial \theta_1} & \frac{\partial g_2(\theta)}{\partial \theta_2} & \cdots & \frac{\partial g_2(\theta)}{\partial \theta_p} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial g_r(\theta)}{\partial \theta_1} & \frac{\partial g_r(\theta)}{\partial \theta_2} & \cdots & \frac{\partial g_r(\theta)}{\partial \theta_p} \end{bmatrix}.$$

Geometrically, the covariance matrix $C_{\hat{\theta}}$ can be visualized in the space of the estimation error by the concentration ellipsoid

$$(\hat{\boldsymbol{\theta}} - \boldsymbol{\theta})^T \mathbf{C}_{\hat{\boldsymbol{\theta}}}^{-1} (\hat{\boldsymbol{\theta}} - \boldsymbol{\theta}) = \boldsymbol{\kappa}$$
(2.9)

where κ is a constant that determines the size of the *p*-dimensional region enclosed by the surface. In these terms an equivalent formulation of the CR inequality reads: For any unbiased estimate of θ the concentration ellipsoid (2.9) lies inside or on the bound ellipsoid defined by

$$(\hat{\boldsymbol{\theta}} - \boldsymbol{\theta})^T \mathbf{J} (\hat{\boldsymbol{\theta}} - \boldsymbol{\theta}) = \kappa .$$
(2.10)

The size and orientation of the ellipsoid (2.10) can be best described in terms of the eigenvalues and eigenvectors of the symmetric $p \times p$ matrix **J**. To this end the eigenvalue problem

$$\mathbf{J}\boldsymbol{\varsigma}_i = \lambda_i \boldsymbol{\varsigma}_i , \quad i = 1, \dots, p , \qquad (2.11)$$

has to be solved, where $\lambda_1, ..., \lambda_p$ are eigenvalues of **J** and $\varsigma_1, ..., \varsigma_p$ the corresponding eigenvectors. The mutually orthogonal eigenvectors coincide with the principal axes of the bound ellipsoid and form an orthogonal matrix $\mathbf{A} = (\varsigma_1, ..., \varsigma_p)$ that diagonalize **J**

$$\mathbf{A}^{T}\mathbf{J}\mathbf{A} = \begin{pmatrix} \lambda_{1} & 0 \\ & \ddots & \\ 0 & & \lambda_{n} \end{pmatrix}.$$
 (2.12)

Thus, rotating the coordinate axes by means of the transformation \mathbf{A}^{T} in the new variables defined by

$$\boldsymbol{\xi} = \mathbf{A}^T (\hat{\boldsymbol{\theta}} - \boldsymbol{\theta}) \,. \tag{2.13}$$

then the ellipsoid (2.13) takes the form

$$(\hat{\boldsymbol{\theta}} - \boldsymbol{\theta})^T \mathbf{J}(\hat{\boldsymbol{\theta}} - \boldsymbol{\theta}) = \sum_{i=1}^n \lambda_i \varsigma_i^2 = \kappa.$$
(2.14)

This is in the coordinates ς_i the equation of an ellipsoid in its standard form with semiaxes of length $\sqrt{\kappa/\lambda_i}$. The bigger λ_i , the more accurate on this semiaxis ς_i , on the contrary, the smaller λ_i , the wider estimation error range will be, if $\lambda_i = 0$, then the value of the corresponding coordinate ς_i can be chosen arbitrarily, i.e. (2.14) describes a degenerate ellipsoid extending to infinity in this coordinate direction. The Fisher information matrix exhibits the information content of the estimation problem. Thus, physically $\lambda_i = 0$ means that there is no information at all about the corresponding coordinate ς_i , i.e., ς_i is unobservable.

CR can not only provide us an intuitive understanding and a deeper insight into the compression for estimation problem by investigating the influence of compression

(selection or quantization) on the estimation accuracy, but also CR is asymptotically achievable by taking the maximum likelihood estimation (MLE) as estimation procedure.

The MLE for a vector parameter $\boldsymbol{\theta}$ is defined to be the value that maximizes the likelihood function $p(\mathbf{x}; \boldsymbol{\theta})$ over the allowable domain for $\boldsymbol{\theta}$. Assuming a differentiable likelihood function, the MLE is found from

$$\frac{\partial \ln p(\mathbf{x}; \mathbf{\theta})}{\partial \theta}\Big|_{\mathbf{\theta}=\hat{\mathbf{\theta}}, w} = 0, \qquad (2.15)$$

If we let $\hat{\boldsymbol{\theta}}_m$ designate the maximum-likelihood estimate of $\boldsymbol{\theta}$ based on N i.i.d. random variables and $\boldsymbol{\theta}_0$ be the true value of the parameter, $\hat{\boldsymbol{\theta}}_m$ converges in law (also called convergence in distribution) to a normal random variable: That is

$$\sqrt{N}(\hat{\boldsymbol{\theta}}_m - \boldsymbol{\theta}_0) \rightarrow Y.$$
(2.16)

where

$$Y \sim N(\mathbf{0}, \mathbf{J}^{-1}(\mathbf{\theta}_0)). \tag{2.17}$$

The ML estimate $\hat{\theta}_m$ is asymptotically efficient in the sense that asymptotically it attains the Cramer-Rao lower bound as $N \to \infty$.

2.2 Detection

Detection can be formulated as a binary statistical hypothesis test. If H_0 and H_1 refer to the hypotheses that the target is absent or present, respectively, we have

$$\begin{cases} H_0 & \mathbf{x} \sim p_0(\mathbf{x}) \\ H_1 & \mathbf{x} \sim p_1(\mathbf{x}, \mathbf{\theta}) \end{cases}$$
(2.18)

where H_0 and H_1 have nonzero priori probabilities P_0 and P_1 , respectively. Under hypothesis H_0 and H_1 , **x** is distributed according to the pdf's p_0 and p_1 . Besides a function of **x**, p_1 is also dependent on the parameter vector **\theta**. The likelihood ratio $L(\mathbf{x}) = p_1(\mathbf{x}, \mathbf{\theta}) / p_0(\mathbf{x})$ is a sufficient statistic for detection, i.e., all we need to know is the likelihood ratio for deciding between the hypotheses H_0 and H_1 . When the parameter vector **\theta** is unknown, the generalized likelihood ratio is usually applied by replacing **\theta** with its MLE (2.15). In the subsequent equations in this section, **\theta** is omitted for the simplicity. The likelihood ratio is invariant to invertible operations such as the transform in the transform coder discussed in the next section.

Under a variety of optimality criteria, the detection algorithms take the form of an LRT

$$L(\mathbf{x}) = \frac{p_1(\mathbf{x})}{p_0(\mathbf{x})} \mathop{\geq}_{H_0}^{H_1} \tau.$$
 (2.19)

where τ is an appropriate threshold. The value of τ depends on the optimality criterion. In a Neyman-Pearson test, the threshold τ is chosen such that, for a given probability of false alarm ($P_f = p_0(L(\mathbf{x}) > \tau)$, the probability of miss ($P_{miss} = p_1(L(\mathbf{x}) < \tau)$) is minimized, or for a given P_{miss} , P_f is minimized. Under the minimum-probability-oferror rule (Bayes rule), the optimal decision is $\hat{i} = \arg \max_i P[H_i | \mathbf{x}] = \arg \max_i [P_i p_i(\mathbf{x})]$, where P_i is priori probability for hypothesis H_i . The LRT in (2.19) is then optimal when τ is equal to the ratio P_0 / P_1 of the priori probabilities. The probability of error in this case is

$$P_{e} = \int \min(P_{0}p_{0}(\mathbf{x}), P_{1}p_{1}(\mathbf{x}))d\mathbf{x}.$$
(2.20)

The problem of choosing an optimal decision rule is treated in a broad literature [5], relatively small attention is paid to how to quantify the discriminating ability of data on the detection performance. The discriminating ability of data (or the performance of the detection) can be evaluated through the Chernoff distance. Chernoff distance gives an upper bound on both P_f and P_{miss} :

$$P_{f} \leq \tau^{-s} e^{-\mu_{s}(p_{0}(\mathbf{x}), p_{1}(\mathbf{x}))} , \qquad (2.21)$$

$$P_{miss} \le \tau^{1-s} e^{-\mu_s(p_0(\mathbf{x}), p_1(\mathbf{x}))}, \qquad (2.22)$$

where μ_s is the Chernoff distance defined by:

$$\mu_{s}(p_{0}(\mathbf{x}), p_{1}(\mathbf{x})) = -\ln \int p_{0}(\mathbf{x}) \left(\frac{p_{1}(\mathbf{x})}{p_{0}(\mathbf{x})}\right)^{s} d\mathbf{x}, \quad 0 < s < 1.$$
(2.23)

and τ is the threshold in the LRT of (2.19). When the Bayes rule is applied, $\tau = P_0 / P_1$ and (2.20) together with the fact that $\min(a,b) \le a^s b^{1-s}$, $0 \le s \le 1$, give an upper bound on P_e

$$P_{e} \leq P_{0}^{1-s} P_{1}^{s} e^{-\mu_{s}(p_{0}(\mathbf{x}), p_{1}(\mathbf{x}))}.$$
(2.24)

The bound is very tight within a scale factor [49][73]. The Chernoff bounds (2.21), (2.22), and (2.24) on P_f , P_{miss} , and P_e hold for any distribution of the data and any sample size N.

2.3 Data Compression and Decompression

Data compression methods are commonly developed either under a classical ratedistortion viewpoint [1] or an operational rate-distortion viewpoint [11],[29]. The classical viewpoint strives to develop methods that are optimal *on average*, over an ensemble of realizations of a random process model; this necessarily demands a random model for the signal and knowledge of a probability structure. The operational viewpoint specifies a compression framework (whose design is often based on insights from the classical viewpoint) and then optimizes the operating point of that framework for the particular signal at hand; this has the advantage of relaxing the assumptions made on the signal (e.g., can assume it is deterministic) but has the disadvantage that side information describing the operating point must be included as overhead in the compressed bit stream. Because a sensor network would likely be required to operate in an abundance of differing signal environments, in this dissertation, we focus on the operational viewpoint to remove the necessity of assuming (limiting) statistical models for the signal. Typically, the operational framework uses numerically computable allocations of rates resources (see [29], [33]) rather than classical closed forms such as reverse water-filling [1].

Moreover, we discuss compression under the umbrella of transform-based coders, which are ubiquitous in practice. The choice of transform coding is due to (i) it provides us the rigid theoretical analysis and design of optimal compression for detection and estimation tasks; (ii) it provides us good compression performance and low computational complexity [31][32] because the performance of compression is proportional to the time consumed and hardware complexity of compression algorithm and transform coding provides us the best tradeoff in the energy-rate-accuracy space, i.e., transform coding can achieve a given certain compression performance in a more efficient way than other compression schemes.

21
samples ×	Transform $\mathbf{y} = T(\mathbf{x})$	samples y	Quantizer $\mathbf{q} = Q(\mathbf{y})$	indices ► q	Encoder $\mathbf{c} = C(\mathbf{q})$	
Reconstruction	Transform $\hat{\mathbf{x}} = T^{-1}(\hat{\mathbf{y}})$	samples ◀ ŷ	Dequantizer $\hat{\mathbf{y}} = \widetilde{Q}^{-1}(\mathbf{q})$	indices q	Decoder $\mathbf{q} = C^{-1}(\mathbf{c})$	

Figure 2.2: The elements of compression-decompression process.

The structure of a transform coder is illustrated in Figure 2.2, the first step is to transform the original digitized sensed data $\mathbf{x} = [x[0], \dots, x[N-1]]^T$ (x can be a vector for one-dimensional signal or a matrix for two-dimensional image) into a new set of samples, which are more amenable to compression according to a criteria best for inference tasks at hand (see Chapter 3 for details). For this step, we write y = T(x), where y is another finite sequence, having the same number of elements as \mathbf{x} . The operator T is invertible (not a singular matrix); For its properties will be discussed shortly. The decompressor employs the inverse transform, T^{-1} , and no distortion is introduced by this step. The second step is a many-to-one mapping to represent the transform samples approximately using a sequence of quantization indices. For this step we write $\mathbf{q} = Q(\mathbf{y})$, where \mathbf{q} denotes the finite sequence of quantization indices. The set of possible outcomes for each quantization index is generally much smaller than that for the transform samples; also, the number of such quantization indices is no larger and may be smaller than the number of transform samples. The decompressor uses an approximate inverse, \tilde{Q}^{-1} . Thus, the quantization mapping function, Q, introduces distortion to the data and makes the reconstructed $\hat{\mathbf{x}}$ not identical to \mathbf{x} . The quantization indices are coded to form the final

bit-stream. We write $\mathbf{c} = C(\mathbf{q})$. This step is invertible and introduces no distortion so that the decompressor may recover the quantization indices as $\mathbf{q} = C^{-1}(\mathbf{c})$.

2.3.1 Coding

The purpose of coding is to exploit statistical redundancy among the quantization indices. The quantization and transform elements are designed in such a way as to ensure that the redundancy is localized. Ideally, the underlying random variables are all statistically independent. In that case, the indices may be coded independently and the only form of statistical redundancy which need be considered is that associated with any non-uniformity in their probability distribution.

2.3.2 Transforms

The transform is responsible for mapping the original samples into a form which enables comparatively simple quantization and coding operations. On the one hand, the transform should capture the essence of statistical dependencies among the original samples so that the group of adjacent transform samples and the quantization indices possess common characteristics and exhibit at most only very local dependencies, ideally, independent; On the other hand, the transform should separate irrelevant information from relevant information according to certain criteria so that the irrelevant samples can be identified and quantized more heavily or even discarded whereas relevant samples are quantized lightly.

We can consider an $N \times N$ linear transform matrix **A**, which maps an *N*-dimensional input vector, **x**, into an *N*-dimensional output vector, **y**, according to

$$\mathbf{y} = \mathbf{A}\mathbf{x} \,. \tag{2.25}$$

Here, we restrict our attention to invertible transforms, writing the inverse as

$$\mathbf{x}_1 = \mathbf{S}\mathbf{y} , \qquad (2.26)$$

where **S** is the inverse of **A**; i.e., SA = I, the $N \times N$ identity matrix. In this case, **S** is the unique inverse of **A**, which we may write as

$$S^{-1} = A$$
, or $S = A^{-1}$. (2.27)

Observe that the transform coefficients may be expressed as

$$y_q = \mathbf{a}_q \mathbf{x}$$
, $q = 0, 1, ..., N - 1$. (2.28)

where \mathbf{a}_q is the *q* th column of the $n \times n$ matrix, **A**. We refer to \mathbf{a}_q as the q^{th} "analysis vector," since it "analyzes" the original vector **x**, to determine its q^{th} transform coefficient. Accordingly, we refer to **A** as the analysis matrix. Also, the inverse transform may be expressed as

$$\mathbf{x}_1 = \sum_{q=0}^{n-1} \mathcal{Y}_q \mathbf{s}_q , \qquad (2.29)$$

where \mathbf{s}_q is the q^{th} column of the $n \times n$ matrix, \mathbf{S} . We refer to \mathbf{s}_q as the q^{th} "synthesis vector," since \mathbf{x}_1 is "synthesized" from a linear combination of the \mathbf{s}_q , with the transform coefficients serving as the weights. Accordingly, the matrix, \mathbf{S} , is known as the "synthesis matrix".

A transform is said to be orthonormal if the analysis vectors are all mutually orthogonal and have unit norm (length); i.e.,

$$\mathbf{a}_i^H \mathbf{a}_j = 0, \ \forall i \neq j, \tag{2.30a}$$

$$\mathbf{a}_i^H \mathbf{a}_i = \left\| \mathbf{a}_i \right\|^2 = 1, \ \forall i$$
(2.30b)

This means that $AA^{H} = A^{H}A = I$, so that $S = A^{H}$ is a unitary matrix. Equivalently, the analysis and synthesis vectors for orthonormal transforms are identical. An orthonormal

transform performs an orthonormal expansion of the input signal as the sum of its projections onto each of the basis vectors; i.e.,

$$\mathbf{x} = \sum_{q} y_{q} \mathbf{s}_{q} = \sum_{q} \left(\mathbf{x}^{H} \mathbf{s}_{q} \right) \cdot \mathbf{s}_{q} .$$
 (2.31)

An important property of orthonormal transforms/expansions is that they are "energy preserving," meaning that

$$\sum_{p} \left| \boldsymbol{x}_{p} \right|^{2} = \left\| \boldsymbol{x} \right\|^{2} = \boldsymbol{x}^{H} \boldsymbol{x} = \left(\sum_{i} y_{i} \boldsymbol{s}_{i}^{H} \right) \left(\sum_{j} y_{j} \boldsymbol{s}_{j} \right)$$
$$= \sum_{i,j} y_{i}^{*} y_{j} \boldsymbol{s}_{i}^{H} \boldsymbol{s}_{j} = \sum_{q} \left| y_{q} \right|^{2} = \left\| \boldsymbol{y} \right\|^{2}$$
(2.32)

In words, the sum of the squares of the input samples (energy of the input), is identical to the sum of the squares of the transform coefficients (energy of the output). This property is often known as Parseval's relation.

To appreciate the significance of this property for compression, consider the compression system of Transform Coding Structure. Let $\mathbf{e}_y = \mathbf{y} - \hat{\mathbf{y}}$ denote the error introduced into the transform coefficients by quantization. Similarly, let $\mathbf{e}_x = \mathbf{x}_1 - \hat{\mathbf{x}}_1$ denote the error introduced into the reconstructed samples by the entire compression system. By linearity of the transform, $\mathbf{e}_x = \mathbf{S}\mathbf{e}_y$ and if the transform is orthonormal, $\|\mathbf{e}_x\|^2 = \|\mathbf{e}_y\|^2$. In words, the error energy in the samples is identical to the error energy in the transform coefficients. Minimizing the mean square error (MSE) of the quantized transform coefficients is then identical to minimizing the MSE of the reconstructed samples. The use of an orthonormal transform has no impact on the rate-distortion function and the selection of an "ideal" transform ensures that the transform bands may be quantized and coded independently without penalty.



rigure 2.5. Oraphical representation of scalar quantiz

2.3.3 Quantization

Quantization is the element of lossy compression systems responsible for reducing the precision of data (reduce the wordlength of samples) in order to make them more compressible. In most lossy compression systems, quantization is the sole source of distortion. Next, the most widely used quantization (scalar quantization) will be reviewed.

Scalar quantization (SQ) is the simplest of all lossy compression schemes. It can be described as a function that maps each element in a subset of the real line to a single value in that subset. Consider partitioning the real line into M disjoint intervals

$$I_q = [t_q, t_{q+1}], \qquad q = 0, 1, \dots, M - 1.$$
 (2.33)

with

$$-\infty = t_0 < t_1 < \cdots < t_M = +\infty.$$

Within each interval, a point \hat{x}_q is selected as the output value (or codeword) of I_q . The process of scalar quantization is illustrated in Figure 2.3. A scalar quantizer is basically a mapping from continuous \Re to discrete $\{0,1,\ldots,M-1\}$. Specifically, for a given x, Q(x) is the index q of the interval I_q which contains x. The dequantizer is given by

$$\widetilde{Q}^{-1}(q) = \hat{x}_q$$

when $x \in I_q = [t_q, t_{q+1})$, that is $\tilde{Q}^{-1}(Q(x)) = \hat{x}_q$. Clearly, the t_q can be thought of as thresholds, or decision boundaries for the \hat{x}_q . The size of the quantization index set is constant, M, therefore only $\log_2 M$ bits will be needed to signal the codeword chosen by the quantizer, which is why a SQ that has $M = 2^b$ boundaries is often called *b*-bit quantizer. There are a number of scalar quantizers (such as uniform quantizer, Lloydmax scalar quantizer, entropy coded scalar quantization) which differ in the choice of the boundaries t_q and the reconstruction value \hat{x}_q . For example, for uniform quantization [10] the decision boundaries are spaced evenly, except for the two outer intervals, the reconstruction values are also spaced evenly, with the same spacing as the decision boundaries: in the inner intervals, they are the midpoints of the intervals. One of the most used uniform quantizer is the midrise quantizer that can be represented as

$$Q(x) = \Delta \left\lfloor \frac{x}{\Delta} \right\rfloor + \frac{\Delta}{2},$$

Here, Δ denotes the step size and $\lfloor \cdot \rfloor$ denotes an operator that rounds downwards to the nearest integer. Different from the uniform quantizer, the Lloyd-Max quantizer [10][11] sets the decision boundaries and reconstruction values to minimize MSE between the input samples and reconstructed values subject to the size of the code (*M*) according to the probability density function (PDF) of the samples. The entropy coded scalar quantizer [11] minimizes MSE subject to a constraint on the entropy of the quantization indices. No matter what type of quantizer, for the high rates (*M* is large or the number of bits of quantizer is large¹) and the data is individually independent

¹ The large is a loose concept, one or two bits could be called large, see [31][32]

distributed (IID), the distortion-rate function, or quantization noise power, can be modeled in the form of

$$d(b) \cong C\sigma^2 2^{-2b}$$
. (2.34)

where σ^2 is the variance of the samples and *C* is a constant which is different for different quantizers [11][30], and is usually determined heuristically in most cases.

A very desirable feature of compression schemes is the ability to successively refine the reconstructed data as the bit-stream is decoded. In this situation, a (perhaps crude) representation of the data to be compressed becomes available after decoding only a beginning small subset of the compressed bit-stream for one purpose. As more of the compressed bit-stream is decoded, the representation of data can be improved incrementally for another purpose. Compression systems possessing this property are facilitated by embedded quantization and this is the basis for the compression for sequential task embedded tasks introduced in Chapter 1.

In embedded quantization, the intervals of higher rate quantizers are embedded within the intervals of lower rate quantizers. Equivalently, the intervals of lower rate quantizers are partitioned to yield the intervals of higher rate quantizers. Consider a sequence of K embedded scalar quantizers Q_0 , Q_1 , Q_2 , ..., Q_{K-1} . The intervals of Q_{K-1} are then embedded within the intervals of Q_{K-2} , which in turn are embedded within those of Q_{K-3} , and so on. Equivalently, the intervals of Q_0 are partitioned to get the intervals of Q_1 , which in turn are partitioned to get the intervals of Q_2 , and so on.

Specifically, each interval of $Q_0 (I_{q_0} \quad q_0 = 0, 1, ..., M_0 - 1)$ is partitioned into M_1 intervals $I_{q_0,q_1} \quad q_0 = 0, 1, ..., M_1 - 1$. The total number of intervals of Q_1 is then M_0M_1 .



Figure 2.4: Embedded scalar quantizers Q_0 , Q_1 , and Q_2 , of rates R = 1, 2, and 3 bits/samples.



Similarly, the intervals of Q_1 are partitioned to obtain the intervals of Q_2 as I_{q_0,q_1,q_2} $q_2 = 0,1,...,M_2 - 1$. In general then, Q_k (k = 0,1,...,K - 1) has $\prod_{j=-}^k M_j$ intervals, given by $I_{q_0,q_1,...,q_k}$. With this partitioning, it is natural to take the comma separated list $q_0,q_1,...,q_k$ as the "quantization index" of $I_{q_0,q_1,...,q_k}$. This situation is illustrated in Figure 2.4 for K = 3. A particular elegant (and important) example is the uniform dead zone quantizer, which is showed in Figure 2.5

2.3.4 Bit Allocation

The compression gain is achieved by limiting the number of bits assigned to the quantizers of transform coefficients. In many cases, since different parts of transform coefficients have different statistics and importance for the tasks at hand, each quantizer has to be optimized for its own transform coefficients and the quantizers are not identical. For example, it is well known that the statistics of high-frequency components of audio

are significantly different from those of the lower frequency components, therefore a subband coder (audio MP3 coder [74]) assigned different quantizer to different spectral bands of audio. The optimal bit allocation (or operational rate distortion optimization for a particular signal at hand instead of classical operational rate distortion optimization for an ensemble of realizations of a random process models) can be formulated as: First, transform coefficients are grouped into M coding units where each coding unit can be quantized by K different available quantizers. For each coding unit i quantized by a quantizer j, we have an average number bits per coefficient r_{ij} for coding unit i and the corresponding distortion d_{ij} caused by the quantizer j. We make no assumptions of any particular structure for the r_{ij} and d_{ij} ; we simply use the convention that quantization indices are listed in order of increasing "coarseness"; i.e. j = 1 is the final quantizer (highest r_{i1} and lowest d_{i1}) and j = K is coarsest. We consider here that the distortion d_{i1} due to r_{i1} is known, or it is possible to replace measured r_{ij} and d_{ij} by values that are estimated on models, but this would not affect the bit allocation algorithm. We then define an objective distortion function $f(\cdot)$ that is a function of d_{ij} and optimize it under the constraint that the total number of bits is upper bounded by the budget R_T . Mathematically speaking, the goal of bit allocation is to find the optimal quantizer b_i^* for each coding unit *i* such that

$$\sum_{i=1}^{N} r_{ib_{i}^{*}} \le R_{T}$$
(2.35)

and the specified form of distortion measure $f(d_{1b_1^*}, d_{2b_2^*}, \dots, d_{Nb_N^*})$ is optimized.

2.4 Compression for Inference Tasks

The last but the most important part of the compression system is for the sensor S_2 to use the decompressed data $\hat{\mathbf{x}}_1$ and its local data \mathbf{x}_2 to do the specific detection or estimation. If we let the output of the estimation component at the sensor S_2 be $h(\hat{\mathbf{x}}_1, \mathbf{x}_2)$, then it is obvious that the error ellipse determined by the estimation error covariance matrix $\mathbf{C}_{\hat{\mathbf{0}}}(\hat{\mathbf{x}}_1, \mathbf{x}_2) = E\{(\mathbf{0} - h(\hat{\mathbf{x}}_1, \mathbf{x}_2))(\mathbf{0} - h(\hat{\mathbf{x}}_1, \mathbf{x}_2))^T\}$ will be greater than that without compression i.e., $\mathbf{C}_{\hat{\mathbf{0}}}(\hat{\mathbf{x}}_1, \mathbf{x}_2) \ge E\{(\mathbf{0} - h(\mathbf{x}_1, \mathbf{x}_2))(\mathbf{0} - h(\mathbf{x}_1, \mathbf{x}_2))\}^T$ due to the reason that the compression algorithm causes certain distortion on \mathbf{x}_1 . Similarly, the resulting probability of detection error is $P_e(\hat{\mathbf{x}}_1, \mathbf{x}_2)$ from the hypothesis test $L(\hat{\mathbf{x}}_1, \mathbf{x}_2)$, it is expected that $P_e(\hat{\mathbf{x}}_1, \mathbf{x}_2) \ge P_e(\mathbf{x}_1, \mathbf{x}_2)$. Therefore, the ultimate goal of the compression algorithm is to minimize its generated distortion on \mathbf{x}_1 such that $h(\hat{\mathbf{x}}_1, \mathbf{x}_2)$ is as close to the true parameter $\mathbf{0}$ as possible, and $P_e(\hat{\mathbf{x}}_1, \mathbf{x}_2)$ is as small as possible.

Although the problem of designing optimal compression algorithms for various applications has received considerable attention and many important results have been established [10], until now, most of the efforts are concentrated primarily on designing compression that performs well in terms of MSE criteria; in other words, the goal of most compression algorithms is to minimize the differences between the original data \mathbf{x}_1 and the reconstructed data $\hat{\mathbf{x}}_1$, i.e., minimize $\|\hat{\mathbf{x}}_1 - \mathbf{x}_1\| = 1/n \sum_{i=1}^n (x_1[i] - \hat{x}_1[i])$. This choice of distortion measure is natural when the sample-to-sample reconstruction is needed. In this spirit, for orthonormal transform and scalar quantization, the bit allocation formulation (2.35) together with quantization distortion (2.34) is [11]

$$f(d_{1b_{1}^{*})}, d_{2b_{2}^{*}}, \dots, d_{Nb_{N}^{*}}) = \sum_{i=1}^{M} C_{i} \eta_{i} \sigma_{i}^{2} 2^{-2r_{ib_{i}^{*}}}$$
(2.36)

where C_i is a constant determined by the ON transform and SQ quantizer used and η_i denotes the ratio between the number of coefficients in the coding unit *i* and the total number of coefficients, and σ_i^2 is the variance the coding unit *i*. The value of σ_i is usually estimated from the coefficients in unit *i*. According to [11], (2.36) can be extended to the more complex quantizer, such as vector quantizers. The classical compression methods for non audio/video signals are typically based on (2.36); even for audio/video, (2.36) is also the primary choice of distortion function [29]. However, when compression algorithms are studied as an element of an inference system, for estimation tasks the ultimate performance measure should be the variance of the estimation error, i.,e., $C_{\hat{x}}(\hat{x}_1, x_2)$, and similarly, for detection tasks the ultimate performance measure should be the probability of detection error, i.e., $P_e(\hat{\mathbf{x}}_1, \mathbf{x}_2)$. MSE is not only inappropriate since it is weakly related to $C_{\hat{x}}$ and P_e , but is not required because inference accuracy, not reconstruction of signal, is the primary objective. Therefore appropriate, useable metrics that measure the impact of rate reduction on the inference performance must be found. Besides, we assume that the inference (detection/ estimation) processing and compression processing are not jointly designed, i.e., the detection/estimation processing uses methods that are optimal in the absence of compression and the compression process is optimal for optimal detector or efficient estimators. Although joint design is preferred from an ultimate optimality perspective, we believe it is of limited value in the inter-operability environment of practice because sensor networks are more likely to be called on to provide data to other systems that are



Figure 2.6 : Framework of the compression method.

independently designed (e.g., "legacy" systems and other "system of systems" scenarios); Since there are numerous approaches to estimate a parameter or detect a signal, the compression algorithms must be universal.

The synopsis of our compression framework is summarized as follows: We model the collected sensor data as a deterministic signal plus additive noise. As illustrated in Figure 2.6, for simplicity and with its own importance, we seek to compress a block of data **x** collected at sensor S_1 so that it can be transmitted to S_2 using no more than a budgeted *R* bits while making the decision about the presence of the signal with the lowest possible probability of detection error and/or making an estimate θ at S_2 (using the compressed data and S_2 -local data) with the smallest mean-square estimation error. The processing box label "**T**" is a linear invertible transform (preferably orthogonal) whose output coefficients are quantized by the set of quantizers $\{Q_i\}$; each coefficient is evaluated to assess its importance or quality for the desired inference tasks and bits are allocated to the various quantizers on that basis. We use Fisher information (2.5) to assess the impact on estimation accuracy and use Chernoff distances (2.23) to assess the impact on decision accuracy. As will be seen, different inference tasks rely on different sets of coefficients. If multiple inference tasks are required, it is also crucial to find methods to properly determine the optimal allocation of the bits to enact the desired trade-offs. The key part of all cases is finding a *method* that can be used to assess the coefficients' quality based on Fisher information and Chernoff distances. We will explore this in detail in the next Chapter.

CHAPTER 3

Data Compression for Parameter Estimation and Detection

3.1 Fisher-Information-Based Data Compression for Parameter Estimation

3.1.1 Overview of Previous Work

Data compression for distributed multi-sensor estimation problems has been previously considered to some degree by many others [34]-[41]. To put our results in context we briefly review the main results of these researchers. Generally, these results are cast in the form shown in Figure 3.1. The source emits a signal $s(\theta)$ depending on an unknown source parameter θ , the *i*th sensor receives signal $X_i(\theta)$ that is related to $s(\theta)$ and encodes it as $\hat{X}_i(\theta)$, which is transmitted to a central processing center where the encoded signals are decoded and then used to make an estimate; Figure. 3.1 shows the estimation of either (i) the parameter θ , or (ii) the signal $s(\theta)$ itself.

These previous results can be grouped in several ways, based on what is being estimated, what compression framework is used, etc. The focus is either on (i) optimizing compression for the purpose of estimation of the source *signal* [16],[17],[34], [35],[58] or on (ii) optimizing compression for the purpose of estimating the source *parameter* [36]-[39],[40],[41]. Many focus on restricting the encoder to being a scalar quantizer that is then designed according to some optimization criteria [34], [35], [36],



Figure 3.1: Classical set-up for compression in a distributed sensor system.

[38], [39], [41], [58]. Others focus on more general compression structures [16],[17],[37],[40], although, most of these ([17],[37],[40]) focus on establishing information theoretic results and bounds. Of these, the closest to our approach is the approach by Lam and Reibman [39], who used the idea of optimizing the compression with respect to an estimation-centric distortion function (the Cramer-Rao lower bound), but they only considered scalar quantization of signal samples rather than a more powerful compression framework like transform coding, as considered here. The results in [41] take an inference-centric viewpoint, although it chooses a quite different metric driven by the specific application they considered; their approach seems considerably less general than that considered here.

3.1.2 Fisher Information for Compression

Fisher information is a well-known concept in estimation theory [4]. It provides a numerical assessment of how useful a set of data is for estimating a parameter. The amount of Fisher information that a set of data possesses is proportional to the parameter estimation accuracy. Its properties include: (i) the Fisher information is invariant under application of invertible maps to the data [78] and is decreased under application of quantization (Appendix B), (ii) the reciprocal of the Fisher information yields the Cramer-Rao Lower Bound on the variance of any unbiased parameter estimator, and (iii) it can be viewed as an "information" measure because it is nonnegative and it is additive for independent observations (Appendix A). However, it is fundamentally different from Shannon information, which measures how many bits must be used to represent a signal. A Shannon-Fisher connection has been explored in [1] in a classical rate-distortion information-theoretic viewpoint.

Let the real¹ data vector **x** be drawn from a probability density function (PDF) $p(\mathbf{x}; \theta)$ that is parameterized by θ , which is to be estimated. The Fisher information (FI) of this estimation problem is defined to be

$$J(\theta; \mathbf{x}) = E\left\{ \left[\frac{\partial \ln p(\mathbf{x}; \theta)}{\partial \theta} \right]^2 \right\}, \qquad (3.1)$$

where the expected value is taken with respect to $p(\mathbf{x}; \theta)$ and therefore the Fisher information is not a function of the data vector \mathbf{x} . However, the *notational* dependence on \mathbf{x} shown on the left-hand side of (3.1) is included merely to keep track of the data set or data subset for which the Fisher information is computed. As indicated on the lefthand side of (3.1), the Fisher information can be a function of the parameter to be estimated, although in many cases it is not.

Clearly, compression of the data vector \mathbf{x} using a lossy algorithm changes the underlying PDF to that of the post-compression data $\hat{\mathbf{x}}$ available for estimation, and therefore alters the Fisher information. More specifically, compression \mathbf{x} decreases the

¹ Modifications to handle the complex case are straightforward [4].

Fisher information (Appendix B). Then, our goal here is to seek *operational* ratedistortion methods to maximize the amount of Fisher information remaining in the data set $\hat{\mathbf{x}}$ while satisfying a budget *R* on the rate used to represent the data set. However, when the Fisher information depends on the parameter θ to be estimated a question arises as to how to precisely characterize our desire to maximize the Fisher information without knowing at what parameter value we should be concentrating our efforts. Thus, in general, we propose a minimax approach. Let $J(\theta; \mathbf{x})$ and $J(\theta; \hat{\mathbf{x}})$ be the Fisher information of the original data \mathbf{x} and the compressed data $\hat{\mathbf{x}}$, respectively. The minimax approach is to compress to *R* bits such that we satisfy

$$\min_{\hat{\mathbf{x}}} \max_{\theta} \left[J(\theta; \mathbf{x}) - J(\theta; \hat{\mathbf{x}}) \right], \tag{3.2}$$

where the maximization is over all the possible parameter values and the minimization is over all compressed $\hat{\mathbf{x}}$ that meet the bit budget R. Clearly, when the Fisher information does not depend on the parameter θ or when the Fisher information can be represented by the product of two independent factor, such as $J(\theta; \hat{\mathbf{x}}) = J_1(\theta)J_2(\hat{\mathbf{x}})$, the minimax approach (3.2) is equivalent to maximizing $J(\hat{\mathbf{x}})$ or $J_2(\hat{\mathbf{x}})$, which are the cases we will focus on. There are surely tasks for which the minimax would be needed, but we found no typical important estimation tasks that required it, and therefore, this general minimax approach is left for future work.

3.1.3 Compressing to Maximize Fisher Information

Let noisy measurement \mathbf{x}_k be the set of data collected at sensor node S_k (k = 1, 2) to support estimation of the parameter θ . We will model the received signal vectors \mathbf{x}_k as

$$\mathbf{x}_k = \mathbf{s}_k(\theta) + \mathbf{w}_k, \quad k = 1, 2, \tag{3.3}$$

where $\mathbf{s}_k(\theta)$ is an unknown deterministic signal vector dependent on the unknown scalar parameter θ to be estimated, and \mathbf{w}_k is a additive noise vector with \mathbf{w}_1 and \mathbf{w}_2 independent of each other. Without loss of generality, we assume that the noise vector \mathbf{w}_1 is white (if not, prewhitening filter [4][5] can be used before compression) and the variance σ_1^2 of \mathbf{w}_1 is known from design considerations and/or sensor measurements in the absence of a signal; the characteristics of \mathbf{w}_2 is not needed to be known for developing the compression.

As shown in Figure 2.1 and Figure 2.6, measurement vector \mathbf{x}_1 will be compressed at S_1 and sent to S_2 ; the decompressed vector $\hat{\mathbf{x}}_1$ is then used with \mathbf{x}_2 to estimate the parameter θ . So the general goal of our compression method is as follows, expressed as transform coding with a Fisher-information-based distortion. Given some orthonormal (ON) decomposition basis set $\{\varphi_n\}_{n=1}^N$ (e.g., wavelet transform, DFT, DCT, etc.) we can write the signal vector to be compressed as

$$\mathbf{x}_{1} = \sum_{n=1}^{N} \chi_{n} \boldsymbol{\varphi}_{n}$$

= $\sum_{n=1}^{N} \xi_{n}(\theta) \boldsymbol{\varphi}_{n} + \sum_{n=1}^{N} \omega_{n} \boldsymbol{\varphi}_{n}$, (3.4)

where χ_n are the data coefficients, $\xi_n(\theta)$ are the parameterized signal coefficients, and ω_n are the noise coefficients.

We wish to select which coefficients χ_n should be quantized and transmitted to achieve a desired rate-Fisher information goal. Only those coefficients with significant contribution to the Fisher information should be selected. The quantization should be done such that only minimal degradation in Fisher information occurs. Let $\Omega \subset \{1, 2, ..., N\}$ be the set of indices of a selected set of coefficients; let $B = \{b_n \mid n \in \Omega\}$ be a set of bit allocations to the selected coefficients; let $\{\hat{\chi}_n \mid n \in \Omega\}$ be the selected coefficients after quantization using the allocation set $B = \{b_n \mid n \in \Omega\}$. The resulting compressed signal is then given by

$$\hat{\mathbf{x}}_1 = \sum_{n \in \Omega} \hat{\boldsymbol{\chi}}_n \boldsymbol{\varphi}_n \quad . \tag{3.5}$$

Our goal, then, is to find a selection set Ω and a bit allocation $B = \{b_n | n \in \Omega\}$ that solves

$$\max_{\Omega,B} J(\hat{\mathbf{x}}_1, \mathbf{x}_2) \quad \text{subject to} \quad \sum_{n \in \Omega} b_n \le R \ . \tag{3.6}$$

Note that the maximization in (3.6) is of the Fisher information of the combination of the two data sets; however, only one of those data sets is available at S_1 , where compression processing is done. However, because the two data sets are statistically independent (due to the assumed independence of the noise vectors), the result in Appendix A shows that $J(\hat{\mathbf{x}}_1, \mathbf{x}_2) = J(\hat{\mathbf{x}}_1) + J(\mathbf{x}_2)$ which means that we can replace (3.6) with

$$\max_{\Omega,B} J(\hat{\mathbf{x}}_1) \quad \text{subject to} \quad \sum_{n \in \Omega} b_n \le R \ . \tag{3.7}$$

(Note that if both sensor signals are to be compressed and transmitted to a fusion center, then the 2nd sensor also would (independently) compress its data to maximize its own $J(\hat{\mathbf{x}}_2)$; we focus here only on compression of a single sensor's data - the later case will be considered in Section 3.1.4.)

We proceed with our development by expressing the Fisher information $J(\mathbf{x}_1)$ in terms of the ON decomposition coefficients. Using (3.4), we can recast our signal model in terms of the transform coefficients as



Figure 3.2: Compression processing, the data vectors, and their corresponding Fisher information.

$$\boldsymbol{\chi}_1 = \boldsymbol{\xi}_1(\boldsymbol{\theta}) + \boldsymbol{\omega}_1, \tag{3.8}$$

where χ_1 is the *N*×1 vector of the coefficients χ_n in (3.4), $\xi_1(\theta)$ is the *N*×1 vector of the signal coefficients $\xi_n(\theta)$ in (3.4), and ω_1 is the *N*×1 vector of the noise coefficients ω_n in (3.4). As outlined above, Fisher information is invariant under any invertible maps and an ON transform is indeed an example of an invertible transform, which gives

$$J(\boldsymbol{\chi}_1) = J(\mathbf{x}_1) , \qquad (3.9)$$

which means that the ON expansion conserves Fisher information; thus, only selection and quantization processing affects the Fisher information, as shown in Figure 3.2.

To develop the compression processing we need to quantify the impact of the selection/quantization processing on the Fisher information. Selection processing is easily modeled because the Fisher information of any discarded coefficient χ_n is zero. The hard part is modeling the effect of quantization of coefficients on the decrease of Fisher information.

In our transform coding framework, the quantizer will be applied to each coefficient χ_n such as

$$r_n = \mathcal{Q}_{b_n}(\chi_n) = \mathcal{Q}_{b_n}(\xi_n(\theta) + \omega_n), \qquad (3.10)$$

where $Q_{b_n}(\cdot)$ denotes $M = 2^{b_n}$ level scalar quantization on χ_n . As showed in Figure 2.3, any b_n -bit scalar quantization can be described as a mapping from \Re to some discrete values $\{q_0, q_1, \dots, q_{M-1}\}$. More specifically, consider partitioning a real line into $M = 2^b$ disjoint intervals

$$I_{j} = [t_{j}, t_{j+1}], \qquad j = 0, 1, \dots, M-1.$$

with

$$-\infty = t_0 < t_1 \cdots < t_{M-1} < t_M = +\infty$$
.

Within each interval, a point q_j is selected as the output value of I_j . Due to quantization (multiple-to-one mapping), r_n takes the discrete form and can only take one value from $\{q_0, q_1, \dots, q_{M-1}\}$, and therefore, the PMF of r_n can be expressed as

$$\operatorname{Prob}(r_n = q_j) = f_{r_n}^{q_j}(b;\theta)$$
$$= \int_{I_j} p(\chi_n) d\chi_n = \int_{I_j}^{I_{j+1}} p(\chi_n) d\chi_n \quad (3.11)$$

where $p(\chi_n)$ is probability density function (PDF) of χ_n . Since r_n are independent discrete values, the Fisher information of r_n can be changed from (3.1) to

$$\hat{J}_{r_n}(b_n,\theta) = \frac{1}{f_{r_n}^{q_j}(b_n;\theta)} \left(\frac{\partial f_{r_n}^{q_j}(b_n;\theta)}{\partial \theta}\right)^2.$$
(3.12)

If we know the probability density function (PDF) of ω_n , and denote it as $p_{\omega_n}(x)$, the closed form of $\widetilde{J}_{r_n}(b_n, \theta)$ can be analytically calculated as

$$\begin{split} \hat{J}_{r_{n}}(b_{n},\theta) &= \left(\frac{\partial(\xi_{n}(\theta))}{\partial\theta}\right)^{2} \left[\frac{p_{\omega_{n}}^{2}(t_{1}-\xi_{n}(\theta))}{\eta_{\omega_{n}}(t_{1}-\xi_{n}(\theta))} + \sum_{j=1}^{M-2} \frac{\left(p_{\omega_{n}}^{2}(t_{j+1}-\xi_{n}(\theta))-p_{\omega_{n}}^{2}(t_{j}-\xi_{n}(\theta))\right)^{2}}{\eta_{\omega_{n}}(t_{j+1}-\xi_{n}(\theta))-\eta_{\omega_{n}}(t_{j}-\xi_{n}(\theta))} \\ &+ \frac{p_{\omega_{n}}^{2}(t_{M-1}-\xi_{n}(\theta))}{1-\eta_{\omega_{n}}(t_{M-1}-\xi_{n}(\theta))}\right] \\ &= \Gamma(\xi_{n})^{2} \left[\frac{p_{\omega_{n}}^{2}(t_{1}-\xi_{n}(\theta))}{\eta_{\omega_{n}}(t_{1}-\xi_{n}(\theta))} + \sum_{j=1}^{M-2} \frac{\left(p_{\omega_{n}}^{2}(t_{j+1}-\xi_{n}(\theta))-p_{\omega_{n}}^{2}(t_{j}-\xi_{n}(\theta))\right)^{2}}{\eta_{\omega_{n}}(t_{j+1}-\xi_{n}(\theta))-\eta_{\omega_{n}}(t_{j}-\xi_{n}(\theta))}\right] \\ &+ \frac{p_{\omega_{n}}^{2}(t_{M-1}-\xi_{n}(\theta))}{1-\eta_{\omega_{n}}(t_{M-1}-\xi_{n}(\theta))}\right] \quad, \end{split}$$
(3.13)

where $\{t_1, ..., t_j, t_{j+1}, t_{M-1}\}$ are thresholds of the chosen quantizer and we have defined $\Gamma(\xi_n) = \partial \xi_n(\theta) / \partial \theta$, which captures the signal's sensitivity to a change in the parameter and put the factors due to quantization and noise inside bracket of (3.13); $\eta_{\sigma_n}(x) = \int_{-\infty}^x p_{\omega_n}(t) dt$ is the cumulative distribution function for PDF of ω_n . The derivation of (3.13) is given in Appendix B. The total Fisher information possessed by quantized ON transform coefficients $\hat{\chi}_1$ is

$$\widetilde{J}_{\hat{\chi}_1} = \sum_n \widetilde{J}_{\chi_n}(b_n, \theta) \,. \tag{3.14}$$

Technically, to implement (3.13), one must derive the form of $\Gamma(\xi_n)$ for the desired estimation task as a function of the *signal* coefficients ξ_n . However, in the operational rate-distortion setting we must instead evaluate the results using the noisy coefficients χ_n . In such case, we have to use χ_n to replace ξ_n in (3.13).

Given the type of quantizer ($\{t_j\}$ are set), we can surely maximize (3.14) under the bit constraint (3.6), however, there is an issue preventing direct optimization of (3.14) from being applied: the computation complexity of (3.13) is severe when b_n is large. Especially, when the size of data N is large, the operations needed to calculate (3.13) will be formidable and reduce the leverage of compression to achieve energy-time efficiency for wireless sensor network. We have to find ways to keep the computation of the Fisher information of quantized data as small as possible. In this dissertation, we propose to use an approximate noise model to complement (3.13) to solve this problem.

The idea behind the approximate model is to model the quantization noise as an additive noise to the measurement; mathematically speaking, after selection/quantization the signal model becomes

$$\hat{\boldsymbol{\chi}}_1 = \widetilde{\boldsymbol{\xi}}_1(\boldsymbol{\theta}) + \widetilde{\boldsymbol{\omega}}_1 + \boldsymbol{\varepsilon},$$
 (3.15)

where a tilde indicates a "reduced-by-selection" version of a vector and $\boldsymbol{\epsilon}$ is the noise vector due to quantization of the selected coefficients. Without loss of generality, we assume that $\tilde{\boldsymbol{\omega}}_1$ is the additive, white Gaussian noise, and uniform quantizer is used such that $\boldsymbol{\epsilon}$ is modeled as white, uniformly distributed, zero mean, and independent of the sensor-noise vector $\tilde{\boldsymbol{\omega}}_1$, where the independence assumption is reasonable because of the dithering effect of the sensor noise [61],[62]. Because of the independence, the variance of the sum of these two noises is the sum of their two variances.

There is a difficulty here because the PDF of $\hat{\chi}_1$ is now not strictly Gaussian, which makes exact evaluation of the Fisher information difficult. Therefore, we choose here to use an approximation that still yields good results: as long as the variance of the uniformly distributed quantization noise ε is small, the sum of the two noises in (3.15) is approximately distributed Gaussian. To see this we note that the PDF of the total noise $\mathbf{v} = \widetilde{\boldsymbol{\omega}}_1 + \varepsilon$ is the convolution of the individual PDFs, therefore the Fourier transform of this convolution is

$$P_{\mathbf{v}}(f) = P_{\tilde{\mathbf{\omega}}}(f)P_{\varepsilon}(f), \qquad (3.16)$$

where $P_{\mathbf{v}}(f)$ is the Fourier transform of the PDF of \mathbf{v} , etc. Because $\widetilde{\mathbf{\omega}}_1$ has a Gaussian PDF, $P_{\tilde{\omega}_1}(f)$ is also a Gaussian function, but with the reciprocal variance; the larger you make the variance of $\widetilde{\boldsymbol{\omega}}_1$ the narrower you make $P_{\widetilde{\boldsymbol{\omega}}_1}(f)$. Because $\boldsymbol{\varepsilon}$ is uniformly distributed, its PDF is a rectangular function so $P_{\varepsilon}(f)$ is a *sinc* function. If the variance of ε is small enough so that the *sinc* function $P_{\varepsilon}(f)$ is broad enough to be approximately flat over the region where $P_{\tilde{\omega}_1}(f)$ is significant, then the product will have approximately the shape of a Gaussian and its inverse Fourier transform (the PDF of the summed noises) will also be approximately Gaussian. See Appendix B for more details of this approximation along with numerical simulation results and comparison between this approximation model and the exact model from (3.13). Through the simulations, we can see that the fitness of approximation to the exact result depends on the number of bits of quantization and the variance σ^2 of $\widetilde{\omega}_1$. In terms of Gaussian noise and uniform quantization, the goodness of approximate model is shown in the Table B.1. For different quantizers, we can have an empirical function called $\gamma(\sigma)$, which defines a threshold above which the approximation model perfectly matches the exact model. Thus, whenever the approximation model is used, we can rewrite (3.13) as

$$\hat{\boldsymbol{\chi}}_1 = \widetilde{\boldsymbol{\xi}}_1(\boldsymbol{\theta}) + \boldsymbol{v}, \qquad (3.17)$$

where **v** is zero-mean Gaussian with variance equal to the sum of the variances of $\boldsymbol{\omega}_1$ and $\boldsymbol{\varepsilon}$. Since the quantizers won't all have the same number of bits allocated, the variance of $\boldsymbol{\varepsilon}$ changes from element to element; let q_n^2 be the variance of the n^{th} element of $\boldsymbol{\varepsilon}$.

Applying the Gaussian noise assumption to the signal model in (3.3), without compression, the Fisher information of the n^{th} coefficient quantized with respect to unknown parameter θ is [4]

$$J(\chi_n) = \frac{1}{\sigma_1^2} \left[\frac{\partial \xi_n(\theta)}{\partial \theta} \right]^2 = \frac{1}{\sigma_1^2} \Gamma^2(\xi_n), \qquad (3.18)$$

If the approximation is used, i.e., when the bits allocated for $\xi_n(\theta)$ is greater than $\gamma(\sigma)$, instead of (3.13), the Fisher information of the n^{th} coefficient quantized to b_n bits can be numerically computed by

$$J_a(\hat{\chi}_n) = \left[\frac{\partial \xi_n(\theta)}{\partial \theta}\right]^2 \frac{1}{\sigma_1^2 + q_n^2} = \frac{\Gamma^2(\xi_n)}{\sigma_1^2 + q_n^2}.$$
(3.19)

Note that the form of (3.19) has a more understandable interpretation than (3.13): we are compressing to maintain a high level of what we will call the "signal-sensitivity-to-noise-ratio" (SSNR). Our goal then is to maximize (3.19) subject to the constraint in (3.6).

Similarly, because we must compute $J(\hat{\chi}_1)$ from our *data* coefficients χ_n , we have to instead use a noisy version given by

$$J(\hat{\chi}_{n}) = \frac{\Gamma^{2}(\chi_{n})}{\sigma_{1}^{2} + q_{n}^{2}},$$
(3.20)

where the hat is used to indicate that the quantity uses the noisy quantities that are available from the data rather than the noise-free values really needed. If we let $\hat{J}_{\iota}(\hat{\chi}_n)$ represents the distortion computation of (3.13), the operational distortion function we use to assess the n^{th} coefficient quantized to b_n bits is

$$\hat{J}_{n} = \begin{cases} 0, & \text{if } b_{n} = 0\\ J_{t}(\hat{\chi}_{n}), & \text{if } 1 \le b_{n} < \gamma(\sigma), \\ J_{a}(\hat{\chi}_{n}), & \text{if } b_{n} \ge \gamma(\sigma) \end{cases}$$
(3.21)

and we seek a bit allocation set $B = \{b_n \ge 0 \mid n \in 1, 2, ..., N\}$ that solves

$$\max_{B} \left\{ \sum_{n=1}^{N} \hat{J}_{n} \right\} \quad \text{subject to} \quad \sum_{n=1}^{N} b_{n} \le R \ . \tag{3.22}$$

Lagrangian and dynamic programming optimization methods are commonly used to determine the bit allocation in operational rate-distortion methods [29]. Lagrangian method is much more effective than dynamic programming, especially when the objective function is a sum of objective functions, and (3.22) is just this case. Although the expensive dynamic programming is said to be superior to Lagrangian method because it contains the exhaustive search, as far as the simulations in this dissertation are concerned, we found out that the differences brought by these two different methods are slight. Therefore, Lagrangian methods are preferred to use for sensor applications. To apply (3.21) we need to have a numerical way to determine the value of J_n as a function of the allocated b_n . Establishing the exact form of quantization noises is generally not possible and it is common to use an approximate model (2.34). We use small blocks of coefficients to estimate the variance needed in (2.34) and then apply the same quantizer to all coefficients in a block for computational efficiency.

There is also the issue of the choice of the ON transform. An advantage of our approach is that once the designer has derived the functional form of the signal sensitivity function $\Gamma(\cdot)$, it generally provides insight into the proper choice of a suitable transform. For example, as will be seen in Section 3.1.4, when the parameter to be estimated is the

relative delay between two sensors' signals, the Fisher information depends on a specific bandwidth measure, thus choosing a frequency domain transform makes sense in that case.

3.1.4. Example Applications

We are particularly interested in problems where data *must* be shared between sensors because neither sensor can estimate the parameter by itself. We call these types of problems "dual-sensor-critical" problems. Such problems often arise in passive systems due to lack of knowledge about a transmitted signal that has been perturbed by some parameter.

It is important to keep in mind that the real essence of FI is that it captures the sensitivity of the data to a *change* in the parameter; this is due to the derivative in the definition of FI in (3.1). Thus, in a passive sensor setting the data at one sensor may be very sensitive to the parameter (i.e., the FI is non-zero) yet you may still be unable to estimate the *value* of the parameter. But with two sets of data (each having a different parameter value) you can estimate the difference in the two parameter values. For example, the data can be *sensitive* to a change in time delay but without a reference it is impossible to estimate the delay. What makes an active sensor problem different is that you have a known reference signal: estimation of the parameter's change from the reference signal leads to estimation of the actual value. In a passive sensor scenario you often need the data at the other sensor to play the role of the reference.

Thus, to make our approach work we only need that the data is sensitive to the parameter, and that is what Fisher information assesses.

48



Figure 3.3 : The spectrum of a typical FM signal used in the simulations.

There are many estimation tasks that are likely to arise in a sensor network and we can't explore them all here. Our goal here is to choose examples that are common distributed estimation problems and (i) illustrate the overall effectiveness of our estimation-centric method relative to MSE-driven methods, and (ii) illustrate how the form of the Fisher information can drive the choice of the transform. We show the application of our method to four cases: (i) time-difference-of-arrival (TDOA) [66], (ii) frequency-difference-of-arrival (FDOA) [66], (iii) direction-of-arrival (DOA) [66], and (iv) signal reconstruction from compressed data. Maximum likelihood estimate is used for each estimate [67][68][70].

The signal we use in our simulations is a complex baseband FM signal with a pseudorandom modulating signal that mimics the spectrum of speech. A realization of the signal spectrum is shown in Figure 3.3.

Compression-estimation processing is simulated as follows: (i) generate a signal at sensor S_1 , (ii) create the signal at sensor S_2 by imparting a TDOA or FDOA or DOA to the S_1 signal, add computer-generated white Gaussian noise to each signal to yield signals

with signal-to-noise ratios of SNR_1 and SNR_2 , compress the noisy S_1 signal, estimate the TDOA or FDOA or DOA or reconstructed signal using maximum likelihood estimators. Under each test condition of interest (i.e., SNR values and compression ratio) we evaluated the RMS error of the estimates by averaging over 400 Monte Carlo runs. In each case we compare the performance of our method to that of standard minimum-MSE transform-based compression (using the same transform as used on our method). We compare the performance on the basis of RMS error of the parameter estimates:

Estimation RMS Error =
$$\sqrt{\frac{1}{M} \sum_{m=1}^{M} (\theta - \hat{\theta}_m)^2}$$

where θ is the true parameter value, $\hat{\theta}_m$ is the estimate made on the m^{th} of *M* Monte Carlo runs. We also compare the MSE performance on the basis of the average vector norm of the signal error between the data vector (\mathbf{x}_1) and the decompressed version of the data vector ($\hat{\mathbf{x}}_1$):

Average Signal Error Norm =
$$\frac{1}{M} \sum_{m=1}^{M} \|\mathbf{x} - \hat{\mathbf{x}}_m\|$$
.

The MSE distortion function is based on (2.36). We should also state that we perform no post-quantization entropy coding (i.e. lossless compression), application of which would likely provide further improvement in the compression ratio with no further degradation in the parameter accuracy or signal MSE; we do this to clearly focus on what is possible using only the FI-based compression.

3.1.4.1 Compression for TDOA Estimation

It is possible to locate a signal source by first estimating the TDOA between several pairs of sensors and then using the TDOA estimates to estimate the source location [66].

The continuous-time signal model for two passively-received complex baseband signals having an unknown TDOA of τ is given by

$$x_{1}(t) = s(t - (t_{0} + \tau/2)) + w_{1}(t)$$

$$x_{2}(t) = s(t - (t_{0} - \tau/2)) + w_{2}(t)$$
(3.24)

where t_0 is also an unknown parameter that can not be estimated, and $w_i(t)$ is complex bandlimited white Gaussian noise. In the frequency domain this model becomes

$$X_{1}(2\pi f) = S(2\pi f)e^{-j2\pi f((t_{0}+\tau/2))} + W_{1}(2\pi f)$$

$$X_{2}(2\pi f) = S(2\pi f)e^{-j2\pi f((t_{0}-\tau/2))} + W_{2}(2\pi f) , \qquad (3.25)$$

where S(f) is the Fourier transform of the source signal. It is well known (see, e.g., [67],[68]) that for TDOA estimation, the Cramer-Rao bound is inversely proportional to the so-called RMS bandwidth B_{rms} of the signal , which is defined by

$$B_{rms} = \sqrt{\frac{\int_{-\infty}^{\infty} f^2 |S(f)|^2 df}{\int_{-\infty}^{\infty} |S(f)|^2 df}}.$$
 (3.26)

Because the Fisher information is the reciprocal of the Cramer-Rao bound, the Fisher information for TDOA estimation is directly proportional to B_{rms} . This analysis of the Fisher information for TDOA estimation explicitly drives the choice of a frequency-domain transform to provide the ability to discard frequency components that contribute little to the Fisher information. Thus we will use the DFT as our transform in our compression processing.

Now assume that we have samples of the signals in (3.25) such that the samples of the bandlimited white noises give discrete-time white noises with variances σ_i^2 , with only σ_1^2 assumed known. We also assume that the sampling of the underlying signal $s(t - (t_0 + \tau/2))$ is such that aliasing is made negligible. Taking the DFT $X_1[k]$ of these samples (taking care to minimize leakage errors) leads to the frequency domain model

$$X_{1}[k] \approx S[k] \exp[-j\frac{2\pi}{N}k(t_{0} + \tau/2)] + W_{1}[k]$$

$$k = -N/2, -N/2 + 1, \dots, N/2 - 1, \qquad (3.27)$$

where the S[k] are the DFT coefficients (for negative and positive frequencies) of the samples of signal s(t) and $W_1[k]$ are the DFT coefficients of the noise.

This is clearly a dual-sensor-critical problem. Each data set is sensitive to changes in the time-of-arrival $t_0 + \tau/2$ but neither sensor by itself can estimate $t_0 + \tau/2$, t_0 alone, or τ alone. In fact, even using both data sets it is impossible to estimate the nuisance parameter t_0 . Although this appears to be a *two*-parameter problem that might require a 2-D FI *matrix*, there really is only *one* thing that matters: sensitivity of the first sensor's data to $\theta \triangleq t_0 + \tau/2$. Of course that is equivalent to sensitivity to τ for a fixed t_0 . Thus, we only need to consider the FI with respect to τ .

Note that because the DFT is an orthogonal, but not ortho*normal*, transform we have to account for the impact of this on the noise variance; thus, $W_1[k]$ is a complex Gaussian random variable with variance $N\sigma_1^2$. Using (3.20), the FI of k^{th} quantized DFT coefficient applying the approximation model becomes

$$\hat{J}_{a}(X_{1}[k]) = \frac{\left(\frac{2\pi}{N}\right)^{2} k^{2} |X_{1}[k]|^{2}}{N\sigma_{1}^{2} + q_{n}^{2}}, \quad k = -N/2, -N/2 + 1, \dots, N/2 - 1.$$
(3.28)

The numerator of (3.28) shows that the DFT coefficients get quadratically weighted based on their frequency's distance from DC. For the condition where the

approximation model fails, we have to turn to the true model. However, because the signal is complex, the scalar quantizer has to be applied to the real and image part of $X_1[k]$ respectively. Let

$$R[k] = \mathcal{Q}_M(\mathfrak{R}(X_1[k]))$$

and

$$\widetilde{R}[k] = \mathcal{Q}_M(\mathfrak{I}(X_1[k]))$$

where $\Re(\cdot)$ and $\Im(\cdot)$ denotes the real parts and imaginary parts of the quantity between the parenthesis, respectively. Since r[k] and $\widetilde{r}[k]$ are independent, the total Fisher information of DFT coefficient $X_1[k]$ is

$$\hat{J}_{t}(X_{1}[k]) = \hat{J}_{t}(R[k]) + \hat{J}_{t}(\widetilde{R}[k]).$$
(3.29)

Applying (3.13), we have

$$\hat{J}_{t}(\tilde{R}[k]) = \frac{\left(\frac{2\pi}{N}\right)^{2} k^{2} (\Im(X_{1}[k]))^{2}}{\pi \sigma_{1}^{2}} \left[\frac{\exp\left(-2\left(\frac{t_{1}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)^{2}\right)}{\Phi\left(\sqrt{2} \frac{t_{1}-\Im(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)} + \frac{\exp\left(-2\left(\frac{t_{d-1}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)^{2}\right)}{\Phi\left(\sqrt{2} \frac{t_{d-1}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right) - \exp\left(-\left(\frac{t_{d}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)^{2}\right)\right)^{2}} + \frac{\exp\left(-2\left(\frac{t_{d-1}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)^{2}\right)}{1-\Phi\left(\sqrt{2} \frac{t_{d-1}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)}\right]$$

$$\hat{J}_{t}(\tilde{R}[k]) = \frac{\left(\frac{2\pi}{N}\right)^{2} k^{2} (\Re(X_{1}[k]))^{2}}{\pi \sigma_{1}^{2}} \left[\frac{\exp\left(-2\left(\frac{t_{1}-\Im(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)^{2}\right)}{\Phi\left(\sqrt{2} \frac{t_{1}-\Im(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)} + \frac{\exp\left(-2\left(\frac{t_{d-1}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)^{2}\right)}{1-\Phi\left(\sqrt{2} \frac{t_{d-1}-\Re(X_{1}[k])}{\sqrt{N\sigma_{1}}}\right)}\right]$$
(3.31)

$$\sum_{j=1}^{M-2} \frac{\left(\exp\left(-\left(\frac{t_{j+1}}{\sqrt{N}\sigma_1}, \Im(X_1[k])\right)^2\right) - \exp\left(-\left(\frac{t_j}{\sqrt{N}\sigma_1}, \Im(X_1[k])\right)^2\right)\right)^2}{\Phi\left(\sqrt{2}\frac{t_{j+1}}{\sqrt{N}\sigma_1}, \Im(X_1[k])\right) - \Phi\left(\sqrt{2}\frac{t_j}{\sqrt{N}\sigma_1}, \Im(X_1[k])\right)} + \frac{\exp\left(-2\left(\frac{t_{M-1}}{\sqrt{N}\sigma_1}, \Im(X_1[k])\right)^2\right)}{1 - \Phi\left(\sqrt{2}\frac{t_{M-1}}{\sqrt{N}\sigma_1}, \Im(X_1[k])\right)}\right)$$

where $\Phi(\cdot)$ is the cumulative distribution function for a standard normal PDF and is defined as

$$\Phi(x) = \int_{-\infty}^{x} \frac{1}{\sqrt{2\pi}} \exp(-\frac{1}{2}t^2) dt,$$

and $\{t_j\}$ are thresholds that are decided by the chosen scalar quantizer, and range of both $\Re(X_1[k])$ and $\Im(X_1[k])$.

Using (3.28) to (3.31) in (3.21) with the Lagrange optimization method produces the results shown in Figure 3.4 to Figure 3.9, which show the performance of our compression method for a few scenarios. Figure 3.4 and Figure 3.5 show (for compression ratios (CR) of 4:1 and 8:1, respectively) the TDOA accuracy performance of our method (labeled "Fisher") vs. the performance of standard MSE-optimum DFT-based transform compression (labeled "MSE") as the SNR_1 at sensor S_1 is varied; the value of SNR_2 at sensor S_2 is fixed at 40 dB. The legend on the plot indicates the cases where just the signal at S_1 is compressed (labeled "S1") and the cases where the signals at S_1 and S_2 were both compressed (labeled "S1&S2" and with dashed curves on the plots for easier identification). For comparison, the performance with no compression is given (labeled "w/o comp"). The effect of varying SNR₂ is shown in Figure 3.6 and Figure 3.7 (for CRs of 4:1 and 8:1, respectively), where results are shown for the case of both SNR's changing but set equal to each other. Figure 3.8 and Figure 3.9 show, for further insight, the performance of our method and the standard MSE-optimal DFT-based transform compression relative to the reconstruction error for compression ratios of 4:1 and 8:1, respectively; the vertical axis shows the norm of the error vector between the original

(noisy) signal and its compressed version, which assesses the MSE performance of the two methods.

In all cases shown, our method provides better TDOA accuracy than the MSEoptimized method; at moderately high SNR our method is nearly the same as when no compression is used – even when the CR is 8:1. In some cases the TDOA accuracy of the MSE-based method is quite poor (especially at higher CRs and when both signals are compressed) – indicating that optimizing MSE provides no guarantee of good performance relative to parameter estimation accuracy. Also note that when both sensor's signals have been compressed at 8:1, our method's performance is degraded a small amount where as the performance of the MSE-based method is severely degraded. However, notice in Figure 3.8 and Figure 3.9 that the TDOA-optimized compression performs very poorly from the reconstruction point of view – this shows that compression for reconstruction and compression for TDOA accuracy are in conflict with each other; our method favors TDOA accuracy at the expense of MSE while the reference method favors MSE at the expense of TDOA accuracy.

The above algorithm can be also extended to apply to image recognition applications. The task in image recognition is finding a specific match in an image. However this image may be distorted and compressed from the original. The Fisher-information-based algorithms that aim to maximize the ability of images to recognize are detailed in Appendix C.

3.1.4.2 Compression for FDOA Estimation

It is possible to locate a signal source by first estimating the FDOA between several pairs of sensors and then using the FDOA estimates to estimate the source location [66].



Figure 3.4: TDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 4:1; the SNR of the sensor S_2 signal was 40 dB.



Figure 3.5: TDOA accuracy vs SNR of pre-compressed sensor S1 signal for a CR of 8:1; the SNR of the sensor S2 signal was 40 dB.



Figure 3.6: TDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 4:1; the SNR of the sensor S_2 signal was set equal to SNR₁.



Figure 3.7: TDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 8:1; the SNR of the sensor S_2 signal was set equal to SNR₁.


Figure 3.8: Reconstruction accuracy vs SNR of pre-compressed sensor S_1 signal for a compression ratio of 4:1.



Figure 3.9: Reconstruction accuracy vs SNR of pre-compressed sensor S_1 signal for a compression ratio of 8:1.

It is well known (see, e.g.,[67],[68]) that for FDOA estimation, the Cramer-Rao bound is inversely proportional to the so-called RMS duration D_{rms} of the signal , which is defined by

$$D_{rms} = \sqrt{\frac{\int_{-\infty}^{\infty} t^2 |s(t)|^2 dt}{\int_{-\infty}^{\infty} |s(t)|^2 dt}} .$$
(3.32)

Because the Fisher information is the reciprocal of the Cramer-Rao bound, the Fisher information for FDOA estimation is directly proportional to D_{rms} . This analysis of the Fisher information for FDOA estimation explicitly drives the choice of the identity transform (which is an ON transform) to provide the ability to discard time components that contribute little to the Fisher information. Thus we will directly quantize the complex-valued signal samples, using individual quantizers for the real and imaginary parts.

Given this insight, we now derive the form for the signal sensitivity function $\Gamma(\cdot)$ in terms of the samples of the received signal. The signal model for two passively-received signals having an unknown FDOA of v is given by

$$x_{1}[n] = s[n]e^{j(v_{0}+v/2)n} + w_{1}[n] \quad n = -N/2, -N/2+1, \dots, N/2$$

$$x_{2}[n] = s[n]e^{j(v_{0}-v/2)n} + w_{2}[n] \quad n = -N/2, -N/2+1, \dots, N/2 , \qquad (3.33)$$

where v_0 is an unknown nuisance parameter that can not be estimated, and $w_i[n]$ is complex Gaussian noise with variance of σ_i^2 , with σ_1^2 assumed known. The discussion in Appendix D addresses why symmetric indexing is chosen in (3.33). The model in (3.33) is mathematically identical to the TDOA model for the DFT transform in (3.27) and therefore we can use the previous results to immediately state that the FI of quantized per-sample in terms of approximation model becomes

$$\hat{J}_{a}(x_{1}[n]) = \frac{1}{2} \frac{n^{2} |x_{1}[n]|^{2}}{\sigma_{1}^{2} + q_{n}^{2}}, \quad n = -N/2, -N/2 + 1, \dots, N/2.$$
(3.34)

On the other hand, in terms of calculation of $\hat{J}_t(x_1[n])$ by applying (3.13), we let $r[n] = Q_M(\Re(x_1[n]))$ and $\tilde{r}[n] = Q_M(\Im(x_1[n]))$, due to the reason that r[n] and $\tilde{r}[n]$ are independent, the total Fisher information of quantized samples $\hat{x}_1[n]$ is

$$\hat{J}_{t}(\hat{x}_{1}[n]) = \hat{J}_{t}(r[n]) + \hat{J}_{t}(\tilde{r}[n]).$$
(3.35)

where

$$\hat{J}(r[n]) = \frac{n^{2} (\mathfrak{I}(x_{l}[n]))^{2}}{4\pi\sigma_{l}^{2}} \left[\frac{\exp\left(-2\left(\frac{t_{l}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)^{2}\right)}{\Phi\left(\sqrt{2} \frac{t_{l}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)} + \frac{\exp\left(-\left(\frac{t_{l+1}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)^{2}\right) - \exp\left(-\left(\frac{t_{l}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)^{2}\right)\right)^{2}}{\Phi\left(\sqrt{2} \frac{t_{l+1}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right) - \Phi\left(\sqrt{2} \frac{t_{l}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)} + \frac{\exp\left(-2\left(\frac{t_{M-1}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)^{2}\right)\right)}{1-\Phi\left(\sqrt{2} \frac{t_{M-1}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)} + \frac{\hat{J}(\tilde{r}[n]) = \frac{n^{2} (\mathfrak{R}(x_{l}[n]))^{2}}{4\pi\sigma_{l}^{2}} \left[\frac{\exp\left(-2\left(\frac{t_{l}-\mathfrak{I}(x_{l}[n])}{\sigma_{l}}\right)^{2}\right)}{\Phi\left(\sqrt{2} \frac{t_{l}-\mathfrak{I}(x_{l}[n])}{\sigma_{l}}\right)} + \frac{\exp\left(-2\left(\frac{t_{M-1}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)\right)}{1-\Phi\left(\sqrt{2} \frac{t_{M-1}-\mathfrak{R}(x_{l}[n])}{\sigma_{l}}\right)} \right] \tag{3.37}$$

where the definitions of $\Phi(\cdot)$ and $\{t_j\}$ are same as in (3.30) and (3.31).

Using these results (3.34)-(3.37) in (3.21) with the Lagrange optimization method produces the results shown in Figure 3.10 and Figure 3.11, which show the performance of our method (labeled "Fisher") and the standard MSE-optimum time-domain



Figure 3.10: FDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 4:1; the SNR of the sensor S_2 signal was 40 dB.



Figure 3.11: FDOA accuracy vs SNR of pre-compressed sensor S_1 signal for a CR of 8:1; the SNR of the sensor S_2 signal was 40 dB.

compression (labeled "MSE") relative to the FDOA RMS error value for compression ratios of 4:1 and 8:1 as SNR₁ at sensor S_1 is varied; the value of SNR₂ at sensor S_2 is fixed at 40 dB, respectively; We have included the case where just the signal at S_1 is compressed (labeled "S1") as well as the case where the signals at S_1 and S_2 were both compressed (labeled "S1&S2"); these figures also show the FDOA RMS error that is achieved when no compression is used (labeled "w/o compression"). In all cases shown in Figure 3.10 and Figure 3.11, our method provides better FDOA accuracy than the MSE-optimized method.

3.1.4.3 Compression for DOA estimation

Direction of arrival (DOA) estimation is a classical problem that has been extensively studied from numerous aspects. A wide variety of techniques have been proposed for the DOA estimation of narrowband sources in the far-field [4],[70], [71]. Let us consider an array of M uniformly spaced omni-directional sensors arranged in a uniform linear array (ULA). Let d denote the spacing between two consecutive sensors. Let λ denote the signal's wavelength, and let φ (-90° < φ < 90°) denote DOA of a narrowband signal that is illuminating the array.

The DOA is measured clockwise with respect to the normal to the line of sensors, then the matrix of observations (snapshots) is given by

$$\mathbf{x}(t_i) = \mathbf{\alpha}(\varphi)s(t_i) + \mathbf{w}(t_i), \quad i = 1, 2, \dots, N,$$
(3.38)

where received signal vector $\mathbf{x}(t_i) = [x_1(t_i), \dots, x_M(t_i)]^T$ and receive noise vector $\mathbf{w}(t_i) = [w_1(t_i), \dots, w_M(t_i)]^T$ and the steering vector is given by

$$\boldsymbol{\alpha}(\boldsymbol{\varphi}) = \begin{bmatrix} 1, e^{-jw_s}, \cdots, e^{-j(M-1)w_s} \end{bmatrix}^T$$
(3.39)

where the spatial frequency w_s is given by

$$w_s = 2\pi f_s = \frac{2\pi d\sin\varphi}{\lambda}$$

Let us look at a two-sensor array $(S_1 \text{ and } S_2)$ to be consistent with the context in this Chapter. Despite being of interest in itself, the two-sensor array case also gives us a glimpse at the effects quantization has on multi-sensor array. Let the snapshot at the two sensor $(S_1 \text{ and } S_2)$ be

$$x_{1}[m] = s[m] \exp[-j(w_{0} - w_{s}/2)] + w_{1}[m] \quad m = 0, 1, ..., N-1$$

$$x_{2}[m] = s[m] \exp[-j(w_{0} + w_{s}/2)] + w_{2}[m] \quad m = 0, 1, ..., N-1$$
(3.40)

After we stack all the received samples $x_1[m]$ into a vector $\mathbf{x}_1(w_s)$ and then we perform any ON transform on $\mathbf{x}_1(w_s)$ to get $\boldsymbol{\chi}_1(w_s)$ whose elements are $\boldsymbol{\chi}_1[m]$.

$$\chi_1[m] = \xi_1[m] \exp[-j(w_0 - w_s/2)] + \omega_1[m] \quad m = 0, 1, \dots, N-1$$
(3.41)

It should be noted that the best ON transform is the phase-invariant transform, otherwise, (3.40) does not hold. The Fisher information with respect to the unknown parameter w_s according to (3.18) is computed by

$$J_{w_s w_s} = \frac{1}{2\sigma_1^2} \sum_{m=1}^{N} \left| \xi_1(m) \right|^2, \qquad (3.42)$$

where $w_s = \frac{2\pi d \sin \varphi}{\lambda}$, $\varphi = \sin^{-1} \left(\frac{\lambda w_s}{2\pi d} \right)$. Then the Fisher information with respect to

DOA φ is computed by [4]

$$J_{\varphi} = \left(\frac{\partial \varphi(w_s)}{\partial w_s}\right)^2 J_{w_s w_s} = \left(\frac{\lambda}{2\pi d \cos \varphi}\right)^2 \left[\frac{1}{2\sigma_1^2} \sum_{m=1}^N \left|\xi_1(m)\right|^2\right].$$
 (3.43)

Observing (3.43), it is clear to see that the Fisher information of DOA φ can be decomposed into the product of two independent terms, one is a function of DOA φ and distance between the sensors *d* and the signal's wavelength λ and the other completely depends on the data $\xi_1(m)$ and variance of noise σ_1^2 . Clearly this problem can be handled by applying our "maximize Fisher information" approach only to the part of the Fisher information that is within the square brackets in (3.43).

Proceeding as before, in terms of the partial Fisher information we intend to maximize, according to (3.20),

$$\hat{J}_{a}(\chi_{m}) = \frac{1}{2} \frac{|\chi(m)|^{2}}{\sigma_{1}^{2} + q_{m}^{2}},$$
(3.44)

and

Using these results (3.44)-(3.46) in (3.21) with the Lagrange optimization method produces the results shown in Figure 3.12 and Figure 3.13, which show the performance of our method (labeled "Fisher") and the standard MSE-optimum time-domain compression (labeled "MSE") relative to the DOA RMS error value for compression ratios of 4:1 and 8:1 as SNR_1 at sensor S_1 is equal to SNR_2 at sensor S_2 and is varied from 10 dB to 40 dB; The wavelet ON transform (Coif5 wavelet, see [72]) is used. Regarding the parameter in (3,42), we choose the distance between the two sensors are equal to the wavelength, $d = \lambda$, and the true DOA $\varphi = 0$ which corresponds to the direction that has the better worse than other angles (See our discussion on the minimax problem). In all cases shown in Figure 3.12 and Figure 3.13, our method provides better DOA accuracy than the MSE-optimized method. The main reason why MSE is inferior to maximizing SNR is in that MSE pays equal weight to all the coefficients, even the coefficients of small amplitude are severely corrupted by the noise whereas maximizing SNR give more bits to those coefficients whose magnitude is much higher than the noise lever. In this case, maximizing SNR has partial denoising functionality.

3.1.4.4 Compression for Signal Estimation

As a further illustrative (though perhaps not entirely practical) example of the flexibility of our Fisher-information-driven viewpoint we consider the scenario

$$x_{1}[n] = s[n - n_{0}] + w_{1}[n] \quad n = 0, 1, ..., N - 1$$

$$x_{2}[n] = s[n] + w_{2}[n] \quad n = 0, 1, ..., N - 1$$
(3.47)

where the goal is to compress the data $x_1[n]$ at sensor S_1 and send it to sensor S_2 where it would be used together with the S_2 -local data $x_2[n]$ to estimate the underlying signal s[n]. For simplicity here we assume that the relative delay n_0 is known either from previous



Figure 3.12: DOA accuracy vs SNR_1 pf pre-compressed sensor S_1 signal for a CR of 4:1; the SNR of the sensor S_2 signal was equal to SNR_2 .



Figure 3.13: DOA accuracy vs SNR_1 pf pre-compressed sensor S_1 signal for a CR of 8:1; the SNR of the sensor S_2 signal was equal to SNR_2 .

estimation processing or from some other *a priori* knowledge (and this is where we veer from a practical example into an illustrative one).

We can estimate the DFT of s[n] instead of estimating s[n] directly, so after taking the DFT of $x_1[n]$ and $x_2[n]$, we have

$$X_{1}[m] = S[m] \exp[-j2\pi mn_{0}] + W_{1}[m] \quad m = 0, 1, ..., N - 1$$

$$X_{2}[m] = S[m] + W_{2}[m] \qquad m = 0, 1, ..., N - 1$$
(3.48)

From this we see that $X_1[m]$ and $X_1[m]$ satisfy the complex classical linear model [4] so (3.48) is equivalent to

$$\mathbf{X}_{1} = \begin{bmatrix} \exp[-j2\pi n_{0}0] & 0 & \cdots & 0 \\ 0 & \exp[-j2\pi n_{0}] & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & \exp[-j2\pi n_{0}(N-1)] \end{bmatrix} \mathbf{S} + \mathbf{W}_{1}$$
(3.49)

and

$$\mathbf{X}_{2} = \begin{bmatrix} 1 & 0 & \cdots & 0 \\ 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 1 \end{bmatrix} \mathbf{S} + \mathbf{W}_{2}.$$
 (3.50)

Applying results for the Fisher information for the classical linear model from [4] to (3.48) and (3.49) gives

$$J(\mathbf{X}_{1}, \mathbf{X}_{2}) = \frac{1}{N\sigma_{1}^{2}} + \frac{1}{N\sigma_{2}^{2}}$$
(3.51)

From this, let us assume fine quantization on \mathbf{X}_1 , that is when the approximation model is valid, we see that after quantization the per-coefficient Fisher information to use in (3.20) for optimizing the compression of sensor S_1 's data becomes

$$\widetilde{J}_n = \frac{1}{N\sigma_1^2 + q_n^2}, \quad n = 0, 1, \dots, N-1.$$
 (3.52)

However, maximizing this is equivalent to minimizing MSE. Thus, this is an illustration that our method is flexible enough to cover the two cases mentioned in Section 3.1.1: parameter estimation and signal estimation. Also, it shows that when MSE is appropriate our method inherently uses it.

In practice, sensor networks may also be tasked with multiple estimation tasks from multiple users, this would make it a multiple parameter problem where the compression requirements are in conflict with each other. This is discussed in Chapter 4.

3.2 Chernoff-Distance-Based Data Compression for Detection

3.2.1 Overview of Previous Work

Using the same set-up as in Section 3.1, we can formulate target detection based on compressed data $\hat{\mathbf{x}}_1$ from S_1 and local data \mathbf{x}_2 at S_2 as a statistical hypothesis testing problem:

$$L(\hat{\mathbf{x}}_{1}, \mathbf{x}_{2}) = \frac{p_{1}(\hat{\mathbf{x}}_{1}, \mathbf{x}_{2})}{p_{0}(\hat{\mathbf{x}}_{1}, \mathbf{x}_{2})} \mathop{\geq}\limits_{H_{0}}^{H_{1}} \tau, \qquad (3.53)$$

where $p_0(\hat{\mathbf{x}}_1, \mathbf{x}_2)$ and $p_0(\hat{\mathbf{x}}_1, \mathbf{x}_2)$ are joint PDF of $\hat{\mathbf{x}}_1$ and \mathbf{x}_2 , and the threshold τ in the LRT (3.53) can be chosen to minimize the probability of error P_e or the probability of miss (P_{miss}). Unfortunately, both P_e and P_{miss} are notoriously intractable functions of the *N*-variate distributions p_0 and p_1 [73], and in general, can only be evaluated experimentally. Hence, it is not feasible to optimize the parameters of complex systems such as lossy coders with respect to P_e or P_{miss} . Some alternative performance measure must be used instead to provide the tractable metric to quantify the degradation of compression on detection accuracy.

In comparison with the few discussions of compression for estimation in literature, various optimality criteria have been analyzed for the problem of optimal quantization for hypothesis testing. Most focus on restricting the encoder to being a scalar quantizer. The composite hypotheses $\theta = 0$ versus $\theta > 0$ was considered in [42] and the efficacy of the sufficient statistic was used as the objective function. Asymptotically optimal companding functions for scalar quantizers were then derived. In [38] [43], a generalized "f-divergence", of the which the Kuallback-leibler distance is a special case, was used as an optimality criterion and its asymptotic quantization effects are studied. A deflection criterion similar to signal-to-noise ratio (SNR) under one of two simple hypotheses is considered in [44], where it is found that maximization of this deflection criterion is by quantizing the likelihood ratio rather then the observation itself. Some properties of these so-called likelihood ratio quantizers are further explored in [45], and their optimality with respect to statistical divergence are also investigated. A method of combining vector quantization and classification was developed in [46][47] by defining an objective function that incorporates MSE and Bayes risk. [48] uses Ali-Silvey distances as optimality criteria and investigate non-asymptotic quantizer effects. The use of alpha entropy as an optimality criterion was proposed in [49] because it gives the exponential decay to zero of the total probability of error of a binary hypothesis test with equal priors, according to a theorem of Chernoff. The distributed hypothesis testing problem with quantized observations is directly addressed in [50]. Optimal scalar quantizers are derived with the Bhattacharyya distance as the objective function and an iterative design algorithm is developed. The Chernoff distance and Kullback-Leibler distance are used to set the information-theoretic bounds on target recognition performance based on degraded image data [51] [52].

The success of designing compression for estimation based on Fisher information motivates us to investigate performance measures that provide tight bounds on P_e and P_{miss} . Since the ability to distinguish between two statistical hypotheses depends on the respective data distributions, using measures of distance or dissimilarity between two distributions sounds natural and is promising.

Ali and Silvey studied a generic category of distances that measure the dissimilarity between two distribution [48]. The Ali-Silvey class of distances is based on an axiomatic definition and takes the general form:

$$d(p_{0}, p_{1}) = f\{E_{0}(C(L))\},$$
(3.54)

where f is any increasing function, C is any convex function on $[0,\infty)$, $L = p_1/p_0$ is the likelihood ratio for the data, and E_0 is the expectation under hypothesis H_0 . It is convenient to also allow pairs (f,C), where f is decreasing and C is concave.

The Ali-Silvey class of distances is decreased under application of many-to-one maps such as quantization. For our problem at hand, it means

$$d(p_0(\hat{\mathbf{x}}), p_1(\hat{\mathbf{x}},)) \le d(p_0(\mathbf{x}), p_1(\mathbf{x})), \qquad (3.55)$$

where $\hat{\mathbf{x}}$ and \mathbf{x} are some $N \times 1$ compressed data vector and unquantized data vector.

Although it has been showed in [43] and [48] how Ali-Silvey distance can be used for optimal quantizer design in detection problems, they limited their design to scalar quantizers. Besides the decrease of distance under quantization, the other attractive property of the Ali-Silvey distance is that they are invariant under application of invertible maps to the data, which directly makes the powerful transform-based compression coding possible.

Chernoff distance (2.23) is in the Ali-Silvey class. More specifically,

$$\mu(s) = -\ln \int p_0(\mathbf{x}) \left(\frac{p_1(\mathbf{x})}{p_0(\mathbf{x})}\right)^s dx , \quad 0 < s < 1$$

$$= -\ln E_0[L^s(\mathbf{x})] \qquad (3.56)$$

where $f(\cdot) = -\ln(\cdot)$ is convex increasing and $C = (\cdot)^s$, 0 < s < 1 is concave. Chernoff distance is the key factor for the Chernoff upper bound of probability of error P_e . [51] provides several bounds derived from Ali-Silvey distance, including Chernoff bound and uses these bounds to guide the design of a quantizer. It is natural to suspect the design that reduces only an upper bound criterion may not always reduce the probability of error. However, Chernoff bound is different from other bounds in that Chernoff's theorem [57] clearly exhibits the exponential tightness of the Chernoff bound given that the samples that make up **x** are I.I.D. Exponential tightness means that for large N.

$$\exp[N(-\mu(s) - \varepsilon)] < P_e < \exp[N(-\mu(s) + \varepsilon)], \qquad (3.57)$$

where $\mu(s)$ is the Chernoff distance for each sample. Clearly any design which increase $\mu(s)$ must (eventually for large *N*) reduce the probability of error.

Clearly, compression of the data vector \mathbf{x} using a lossy algorithm changes the underlying PDF to that of the post-compression data $\hat{\mathbf{x}}$ available for detection, and therefore decreases the magnitude of the Chernoff distance. Similarly, our goal here is to seek *operational* rate-distortion methods to maximize the magnitude of the Chernoff

distance represented by the data set $\hat{\mathbf{x}}$ while satisfying a budget *R* on the rate used to represent the data set.

Let $\mu_s(\mathbf{x})$ and $\mu_s(\hat{\mathbf{x}})$ be the Chernoff distance of the original data \mathbf{x} and the compressed data $\hat{\mathbf{x}}$, respectively. The optimal approach is to compress to *R* bits such that we satisfy

$$\min_{\hat{\mathbf{x}}} [\mu_s(\mathbf{x}) - \mu_s(\hat{\mathbf{x}})], \qquad (3.58)$$

where the minimization is over all compressed $\hat{\mathbf{x}}$ that meet the bit budget R. This is equivalent to maximizing the post-compression $\mu_s(\hat{\mathbf{x}})$.

3.2.2 Compressing to Maximize Chernoff Distance

Let noisy measurement \mathbf{x}_k collected at sensor node S_k (k = 1, 2) be modeled as

$$\begin{cases} \mathbf{x}_k = \mathbf{w}_k \\ \mathbf{x}_k = \mathbf{s} + \mathbf{w}_k \end{cases}, \quad k = 1, 2, \qquad (3.59)$$

where **S** is an unknown deterministic signal vector and \mathbf{w}_k is a noise vector with \mathbf{w}_1 and \mathbf{w}_2 independent of each other. Measurement vector \mathbf{x}_1 will be compressed at S_1 and sent to S_2 ; the decompressed vector $\hat{\mathbf{x}}_1$ is then used with \mathbf{x}_2 to decide the presence of \mathbf{s} . By independence (See Appendix A), $\mu_s(\hat{\mathbf{x}}_1, \mathbf{x}_2) = \mu_s(\hat{\mathbf{x}}_1) + \mu_s(\mathbf{x}_2)$, so we need only maximize $\mu_s(\hat{\mathbf{x}}_1)$.

Similar to the goals in Section 3.1, but expressed here with the Chernoff distance, the general goal of our transform-coding compression method is as follows: we use an orthonormal (ON) basis $\{\varphi_n\}_{n=1}^N$ with χ_n as the coefficients for \mathbf{x}_1 , and ON expansion conserves Chernoff distance, hence $\mu_s(\chi_n) = \mu_s(\mathbf{x}_1)$. Only those χ_n with significant

contribution to the Chernoff distance should be selected and quantized. Let $\Omega \subset \{1, 2, ..., N\}$ be the set of indices of a selected set of coefficients; let $B = \{b_n \mid n \in \Omega\}$ be a set of bit allocations to the selected coefficients; let $\{\hat{\chi}_n \mid n \in \Omega\}$ be the selected coefficients after quantization using the allocation set $B = \{b_n \mid n \in \Omega\}$. The resulting compressed signal is then given by

$$\hat{\mathbf{x}}_1 = \sum_{n \in \Omega} \hat{\boldsymbol{\chi}}_n \boldsymbol{\varphi}_n$$

Our goal, then, is to find a selection set Ω and a bit allocation $B = \{b_n | n \in \Omega\}$ that solves

$$\max_{\Omega,B} \mu_s(\hat{\mathbf{x}}_1) \quad \text{subject to} \quad \sum_{n \in \Omega} b_n \le R \tag{3.60}$$

Let $r_n = Q_b(\chi_n)$ and the PMF of r_n under different hypotheses H_i , i = 0,1 can be

calculated as

$$Prob(r_{n} = \hat{x}_{j} | H_{i}) = f_{r[n]}^{x_{j}}[b_{n}; H_{i}] = \int_{I_{i}} p_{i}(\chi_{n}) d\chi_{n} = \int_{t_{j}}^{t_{j+1}} p_{i}(\chi_{n}) d\chi_{n} , \qquad (3.61)$$

The total Chernoff distance of the independent discrete quantized coefficient r_n can be computed by

$$\boldsymbol{\mu}_{s}(r_{n}) = -\sum_{k=0}^{N-1} \left(\ln \sum_{j} \left(f_{r_{n}}^{\hat{x}_{j}} [b_{n}; H_{0}] \right)^{1-s} \left(f_{r_{n}}^{\hat{x}_{j}} [b_{n}; H_{1}] \right)^{s} \right)$$
(3.62)

See appendix B, the close form of Chernoff distances with respect to each r_n is

$$\mu_{t,s}^{r_n}(\chi_n) = \ln \left[\eta(t_1)^{1-s} \eta(t_1 - \xi_n)^s + \frac{M^{-2}}{2} \left(\eta(t_{j+1}) - \eta(t_j) \right)^{1-s} \times \left(\eta(t_{j+1} - \xi_n) - \eta(t_j - \xi_n) \right)^s + \left(1 - \eta(t_{M-1}) \right)^{1-s} \left(1 - \eta(t_{M-1} - \xi_n) \right)^s \right]$$
(3.63)

where ξ_n is the ON transform of noise-free signal s[n] and $\eta(\cdot)$ is the cumulative distribution function for PDF of ON transformed noise vector \mathbf{w}_1 .

If we do not insist on the optimum selection of the Chernoff parameter s, we may obtain another distance, called the Bhattacharyya distance which also gives us the upper bound

$$\mu_{1/2}(\mathbf{x}) == -\ln E_0[L^{1/2}(\mathbf{x})]. \tag{3.64}$$

We prefer using the Bhattacharyya distance rather than the Chernoff because the Bhattacharyya distance is a popular, simple measure of similarity between two distributions or convenient measure of class separability, if the number of classes is two. Furthermore, it gives an upper bound of the Bayes error [2] if the distributions are normal. Other divergence measures in the literature are not as closely related to Bayes errors as is the Bhattacharyya bound. However, all discussions about the Bhattacharya distance in this dissertation could be extended to the Chernoff.

If the noise \mathbf{w}_1 at S_1 is zero-mean, white Gaussian noise, Bhattacharyya distance takes a simple form like

$$\mu_{1/2}(\mathbf{x}_1) = \frac{1}{8} \frac{\mathbf{s}^H \mathbf{s}}{\sigma_1^2} \quad . \tag{3.65}$$

After ON transformation, and applying the approximation model that we used in Section 3.1 jointly, the post-compression Bhattacharyya distance becomes

$$\mu_{a,1/2}(\chi_n) = \frac{1}{8} \frac{\left|\xi_n\right|^2}{\sigma_1^2 + q_n^2}.$$
(3.66)

It is obvious that maximizing (3.66) in terms of compression is equivalent to maximizing the Signal-to-Noise-and-Quantization ratio. This suits for the compression under the assumption that the signal is deterministic. The goal for compression for detection will be to keep the signal level above the noise and quantization noise floor as far as possible. Notice that (3.66) is only for Gaussian noise, the expression for Bhattacharyya distance under other distribution of noise will be different from (3.66).

If we use the noisy quantities χ_n rather than the noise-free values ξ_n , the operational Bhattacharyya-distance-based distortion function we use to assess the n^{th} coefficient quantized to b_n bits is

$$\hat{\mu}_{1/2}(\hat{\chi}_n) = \begin{cases} 0, & \text{if } b_n = 0\\ \hat{\mu}_{a,1/2}(\hat{\chi}_n), & \text{if } 1 \le b_n < \gamma(\sigma), \\ \hat{\mu}_{t,1/2}(\hat{\chi}_n), & \text{if } b_n \ge \gamma(\sigma) \end{cases}$$
(3.67)

and we seek a bit allocation set $B = \{b_n \ge 0 \mid n \in 1, 2, ..., N\}$ that solves

$$\max_{B} \left\{ \sum_{n=1}^{N} \hat{\mu}_{1/2}(\hat{\chi}_{n}) \right\} \quad \text{subject to} \quad \sum_{n=1}^{N} b_{n} \leq R \,.$$
(3.68)

Simulation results are skipped because maximizing signal SNR for detection is well known and a metric very similar to (3.66) has been applied before in [51][52], although they do not provide the rigid operational function (3.67) as we do. The real interest here is to complete the inference-centric compression framework and explore how multiple measures (for multiple inferences) drive the specification of new multiresolutional characteristic of the data needed to support multiple sequential and simultaneous inferences.

CHAPTER 4

Data Compression for Simultaneous Tasks of Multiple Inference Quantities

In Chapter 3, we have considered data compression for single estimation/detection tasks. However, as pointed out elsewhere [27], sensor networks may also be tasked with multiple inference tasks from multiple users. For example, sensor S_2 may use the shared compressed data stream from sensor S_1 to simultaneously accomplish joint TDOA and FDOA measurements. However, compression for TDOA (3.27)-(3.30) and compression for FDOA (3.33)-(3.36) conflict such that compression for TDOA favors the frequencydomain coefficients of high frequencies whereas compression for FDOA favors the timedomain coefficients of early or late times. For such simultaneous estimation tasks, it is crucial to find the right way to balance the conflicts of compression for different parameters. On the other hand, sensor S_1 may send data for sensor S_2 , where detection is first performed to check if both sensors have the signals intercepted from the same emitter, and then once detected, joint TDOA/FDOA will be performed to locate the emitter. So we also need to find ways to balance the conflicting compression requirements between detection and estimation. To solve this problem, in this Chapter, we will derive multiple distortion measures to assess the information utility of data subsets for the multiples inferences and develop the corresponding numerical optimization methods to achieve desired trade-offs based on the information utility.

4.1 Data Compression for Simultaneous Multiple Parameter Estimation without Detection

Different from single parameter estimation, now we have a parameter vector $\boldsymbol{\theta} = [\theta_1, \dots, \theta_p]$ sought to be estimated, intuitively, we will say that the compression algorithm ε_1 is distinct from the compression ε_2 if at least one of the estimated value θ_i by ε_1 is distinct from the corresponding variable by ε_2 . To design the optimal compression, we need to have criteria for comparing compression algorithms. As suggested in Chapter 2, the natural and ultimate choice of criteria to compare different algorithm relies on their error covariance matrix which is lower bounded by and can be quantified by the Cramer-Rao Lower Bound matrix (2.7). Therefore, we can claim that compression algorithm ε_1 is considered as preferable to compression algorithm $\varepsilon_2(\varepsilon_1 > \varepsilon_2)$ if the difference of $CRB(\mathbf{x}, \varepsilon_2) - CRB(\mathbf{x}, \varepsilon_1)$ is a positive matrix. However, it is not obvious how to use the matrix $CRB(\mathbf{x}, \varepsilon_k)$ as a distortion measure because compression algorithms often require a scalar measure for efficiency. The open question for us is how we map $CRB(\mathbf{x}, \varepsilon_k)$ into a scalar measure of information utility which exploits the information content that is carried by $CRB(\mathbf{x}, \varepsilon_k)$ and is proportional to the "size" of the high probability uncertainty region of the estimate of θ as well.

The determinant det(CRB($\mathbf{x}, \varepsilon_k$)) is proportional to the volume of the rectangular region enclosing the core ellipsoid of covariance ellipsoid of any estimate. Hence, the first information utility function we choose is det(CRB($\mathbf{x}, \varepsilon_k$)). Although the volume of the high probability regions seems to be a useful measure, there are cases in which this measure underestimates the residual uncertainty. In case the smallest principal axis shrinks to zero, the volume of the uncertainty ellipsoid is zero, while the uncertainties along the remaining principal axes might remain large. An alternative measure that avoids this shortcoming of det(CRB($\mathbf{x}, \varepsilon_k$)) would be tr(CRB($\mathbf{x}, \varepsilon_k$)), which is proportional to the circumference of the rectangular region enclosing the covariance ellipsoid. Therefore, the distortion measure for multiple parameter estimation must be able to minimize det(CRB($\mathbf{x}, \varepsilon_k$)) and tr(CRB($\mathbf{x}, \varepsilon_k$)) under the bit constraints.

However, we found that deriving closed-form distortion measures based on either $det(CRB(\mathbf{x}, \varepsilon_k))$ or $tr(CRB(\mathbf{x}, \varepsilon_k))$ is difficult - even for a single parameter case, such as TDOA-only case [24],[60]. Therefore, with the desire to design the most efficient compression algorithms, we resort again to Fisher information.

In Chapter 3, we have used Fisher information to derive transform-coding schemes for the signal parameter cases where the Fisher information (FI) is a scalar. However, when multiple parameter parameters are to be estimated the FI becomes the Fisher information matrix (FIM). The Cramer-Rao bound (2.7) states that the error covariance Σ of any unbiased estimator of x satisfies

$$\Sigma \ge \mathbf{J}^{-1}(\mathbf{x}, \varepsilon_k) = \operatorname{CRB}(\mathbf{x}, \varepsilon_k).$$
(4.1)

It can be shown that the Fisher information matrix (FIM) is related to the surface area of a high probability region, which is a notion of the "size" of the region [1]. It is natural to assume that maximizing the size of the FIM is equivalent to minimization of the size of the CRB. Similar to the measures for CRB, possible forms of the information utility function using the FIM are det($\mathbf{J}(\mathbf{x}, \varepsilon_k)$) and tr($\mathbf{J}(\mathbf{x}, \varepsilon_k)$). Besides easier manipulation of **J** rather then **CRB**, det($\mathbf{J}(\mathbf{x}, \varepsilon_k)$) and det(CRB($\mathbf{x}, \varepsilon_k$)) or tr($\mathbf{J}(\mathbf{x}, \varepsilon_k$)) and tr(CRB($\mathbf{x}, \varepsilon_k$)) can be verified as equivalent to each other due to the close relationships between **J** and **CRB** in (4.1). First of all, it is easy to show from matrix theory that minimizing det(CRB($\mathbf{x}, \varepsilon_k$)) is equivalent to maximizing det($\mathbf{J}(\mathbf{x}, \varepsilon_k)$). Second, the following relationship between CRB($\mathbf{x}, \varepsilon_k$) and $\mathbf{J}^{-1}(\mathbf{x}, \varepsilon_k)$ provides the justification that the maximization of tr($\mathbf{J}(\mathbf{x}, \varepsilon_k)$) is equivalent to the minimization of tr(CRB($\mathbf{x}, \varepsilon_k$)).

Because **J** is positive definite, $CRB = \mathbf{J}^{-1}$ is positive definite, too. Let the eigenvalues of **J** be $\lambda_i > 0$ (i=1,...,n), then $tr(\mathbf{J}) = \sum_{i=1}^n \lambda_i$ and $tr(\mathbf{J}^{-1}) = \sum_{i=1}^n 1/\lambda_i$. Using the Cauchy-Schwarz inequality gives

$$n = \begin{bmatrix} \sqrt{\lambda_i} & \dots & \sqrt{\lambda_n} \end{bmatrix} \begin{bmatrix} 1/\sqrt{\lambda_i} \\ \vdots \\ 1/\sqrt{\lambda_n} \end{bmatrix} \le \sqrt{\sum_{i=1}^k \lambda_i} \sqrt{\sum_{i=1}^k 1/\lambda_i} = \sqrt{\operatorname{tr}(\mathbf{J})\operatorname{tr}(\mathbf{J}^{-1})} , \qquad (4.2)$$

That is, $\operatorname{tr}(\mathbf{J}^{-1}) \ge n^2 / \operatorname{tr}(\mathbf{J})$. Thus the maximization of $\operatorname{tr}(\mathbf{J})$ *tends* to bring the minimization of $\operatorname{tr}(\mathbf{J}^{-1})$, providing further justification for the choice of the trace. Hence, in the sequel, we will use $\operatorname{det}(\mathbf{J}(\mathbf{x},\varepsilon_k))$ or $\operatorname{tr}(\mathbf{J}(\mathbf{x},\varepsilon_k))$ instead of $\operatorname{det}(\operatorname{CRB}(\mathbf{x},\varepsilon_k))$ or $\operatorname{tr}(\operatorname{CRB}(\mathbf{x},\varepsilon_k))$ to derive the theory and algorithm for compression for simultaneous parameter estimation.

Suppose we have two sensor nodes that intercept a common signal and wish to share data to estimate a parameter vector $\boldsymbol{\theta}$. At sensor node S_k we model the received signal vector \mathbf{x}_k as

$$\mathbf{x}_k = \mathbf{s}_k(\mathbf{\theta}) + \mathbf{w}_k, \quad k = 1, 2 , \tag{4.3}$$

where $\mathbf{s}_k(\mathbf{0})$ is an unknown deterministic vector dependent on the unknown deterministic parameter vector $\mathbf{0}$, and \mathbf{w}_k is an additive noise vector with \mathbf{w}_1 and \mathbf{w}_2 independent. To simplify the discussion, but without loss of generality, we will assume \mathbf{w}_1 is a zero-mean, Gaussian noise vector whose covariance matrix Σ is known or estimated. The extension to other noise distributions is obvious because the following procedure for Gaussian noise is quite general (see Chapter 3, the comprehensive study of other noise distribution will be left for further work). Second, we will focus on the two parameter case. Here we use an orthonormal (ON) transform for the linear invertible transform **T** in Figure 2.6. The unitary matrix of this transform will be denoted as $\mathbf{\Phi}$ with columns $\{\mathbf{\varphi}_n\}_{n=1}^N$ that form an ON basis. Expanding data vector \mathbf{x}_1 with respect to this ON basis gives the coefficients $\{\chi_n\}_{n=1}^N$. The coefficients χ_n are then quantized using a set of bit allocations $B = \{b_n \mid n = 1, 2, ..., N\}$. The compressed signal is $\hat{\mathbf{x}}_1 = \sum_n \hat{\chi}_n \mathbf{\varphi}_n$; grouping these coefficients into vector form gives

$$\hat{\boldsymbol{\chi}}_1 = \boldsymbol{\Phi} \mathbf{x}_1 = \boldsymbol{\xi}_1(\boldsymbol{\theta}) + \boldsymbol{\omega}_1 + \boldsymbol{\varepsilon}$$

= $\boldsymbol{\xi}_1(\boldsymbol{\theta}) + \mathbf{v}$, (4.4)

where vector $\xi_1(\theta) = \Phi \mathbf{s}_1(\theta)$ holds the selected signal coefficients $\xi_n(\theta)$, vector $\boldsymbol{\omega}_1$ holds the corresponding noise coefficients $\boldsymbol{\omega}_n$, and $\boldsymbol{\varepsilon}$ is the quantization noise vector that is independent with $\boldsymbol{\omega}_1$. We let **Q** denote the covariance matrix of $\boldsymbol{\varepsilon}$.

As in Chapter 3, we assume that the sensor noises are independent and therefore the total FIM is the sum of the FIMs at the two sensors; thus, we only need to consider the effect of compression on the FIM of the data at sensor S_1 . The FIM for θ based on the data from sensor S_1 is the 2×2 matrix J with elements given by [4]

$$\mathbf{J} = 2 \operatorname{Re} \left\{ \mathbf{G}^{H} \boldsymbol{\Sigma}^{-1} \mathbf{G} \right\}, \tag{4.5}$$

where $\mathbf{G} = \begin{bmatrix} \frac{\partial \mathbf{s}_1}{\partial \theta_1} & \frac{\partial \mathbf{s}_1}{\partial \theta_2} \end{bmatrix}$ is an *N*×2 matrix of the signal's sensitivities to the parameters. The

FIM is invariant under transformation of the data by a unitary matrix; thus, we have that

$$\mathbf{J} = 2 \operatorname{Re} \left\{ \widetilde{\mathbf{G}}^H \boldsymbol{\Sigma}^{-1} \widetilde{\mathbf{G}} \right\} \quad . \tag{4.6}$$

where $\tilde{\mathbf{G}} = \boldsymbol{\Phi}\mathbf{G}$. The FIM specifies an information ellipse via $\boldsymbol{\theta}^T \mathbf{J}^{-1} \boldsymbol{\theta} = \kappa$ – the larger this ellipse the better – with semi-axes along the FIM's eigenvectors and whose lengths are proportional to the square roots of the eigenvalues. Lossy compression of the data vector \mathbf{x} changes all the elements of the FIM, making the data inferior for estimation of $\boldsymbol{\theta}$ and thus shrinking and rotating the information ellipse represented by the FIM. Under the model in (4.4) we have that the FIM after compression is given by

$$\hat{\mathbf{J}} = 2 \operatorname{Re} \left\{ \widetilde{\mathbf{G}}^{H} [\mathbf{\Sigma} + \mathbf{Q}]^{-1} \widetilde{\mathbf{G}} \right\},$$
(4.7)

where the ^ symbol indicates "after compression".

A. Determinant of Fisher Information matrix

The area of the FIM ellipsoid can be represented as $A = C_A \sqrt{\lambda_1 \lambda_2}$, where C_A is a constant and λ_i is the *i*th eigenvalue of the FIM. It is clear from this point that maximizing the area is equivalent to maximizing det(**J**) = $\lambda_1 \lambda_2$. It is also recognized that maximizing det(**J**) is also equivalent to maximizing the entropy of estimation uncertainty when the noise is Gaussian [79],[80] and coincides with other optimal criteria discussed in [78].

The determinant of the FIM has been used as an objective function to be optimized in the solution of engineering problems (see e.g., [76], [77], [78]). However, it is well known that additive distortion yields simpler operational rate-distortion optimization schemes [29] whereas the determinant provides a multiplicative distortion function $(\det(\mathbf{J}) = J_{11}J_{22} - J_{12}^2)$ rather than an additive one, which make it hard to formulate a simple numerical distortion measure to guide compression. However, under the fine quantization condition, we can develop an additive distortion measure aimed at allowing compression that maximizes the determinant.

For convenience of discussion consider that the noise is an i.i.d. process with variance σ^2 , i.e., $\Sigma = \sigma^2 \mathbf{I}$ (a prewhitening filter can be used if the noise is colored), and that the signal data is real valued. After the compression algorithm, the Fisher information matrix of compressed data becomes¹

$$\hat{\mathbf{J}} = \widetilde{\mathbf{G}}^T \operatorname{diag}\left\{\frac{1}{\sigma^2 + q_1^2}, \frac{1}{\sigma^2 + q_2^2}, \dots, \frac{1}{\sigma^2 + q_N^2}\right\} \widetilde{\mathbf{G}}, \qquad (4.8)$$

and its determinant can be manipulated into the form

$$\det(\hat{\mathbf{J}}) = \det\left(\frac{1}{\sigma^{2}}\widetilde{\mathbf{G}}^{T}\widetilde{\mathbf{G}} - \widetilde{\mathbf{G}}^{T}\operatorname{diag}\left\{\left[\frac{q_{1}^{2}}{\sigma^{2}(\sigma^{2}+q_{1}^{2})} \quad \frac{q_{2}^{2}}{\sigma^{2}(\sigma^{2}+q_{2}^{2})} \quad \cdots \quad \frac{q_{N}^{2}}{\sigma^{2}(\sigma^{2}+q_{N}^{2})}\right]\right\}\widetilde{\mathbf{G}}\right)$$

$$= \frac{1}{\sigma^{4}}\det(\widetilde{\mathbf{G}}^{T}\widetilde{\mathbf{G}})\det\left(\mathbf{I} - \underbrace{\left(\widetilde{\mathbf{G}}^{T}\widetilde{\mathbf{G}}\right)^{-1}\widetilde{\mathbf{G}}^{T}\operatorname{diag}\left\{\left[\frac{q_{1}^{2}}{\sigma^{2}+q_{1}^{2}} \quad \frac{q_{2}^{2}}{\sigma^{2}+q_{2}^{2}} \quad \cdots \quad \frac{q_{N}^{2}}{\sigma^{2}+q_{N}^{2}}\right]\right\}\widetilde{\mathbf{G}}\right).$$

$$(4.9)$$

$$= \mathbf{R} \ (defined)$$

Therefore, maximizing $det(\hat{\mathbf{J}})$ is equivalent to $\underset{\{q_0,q_1,\cdots,q_{N-1}\}}{\text{Maximize }}det(\mathbf{I}-\mathbf{R})$. Performing an

eigenvector decomposition on R gives

¹ For the ease of deduction, we use the approximation model $1/(\sigma^2 + q_i^2)$ to quantify the effect of quantization on the reduction of FI. But this does not impair the correctness of (4.8). Referring to Appendix B, the true effect of quantization $I_t(\xi_n(\theta), \Delta, \sigma)$ is always less than that without compression $I(\sigma), 1/(\sigma^2 + q_i^2)$ can be used to represent this reduction of the quantity $1/(\sigma^2)$. q_i^2 is not necessarily

equal to (2.36) any more in this case. Later, we will use this unified approach in our discussion

$$det(\mathbf{I} - \mathbf{R}) = det(\mathbf{Q}^T \mathbf{Q} - \mathbf{Q}^T \operatorname{diag}\{\lambda_1, \lambda_1, \dots, \lambda_N\}\mathbf{Q}) = det(\mathbf{I} - \operatorname{diag}\{\lambda_1, \lambda_1, \dots, \lambda_N\})$$

= 1 - tr(\mathbf{R}) + {sum of product terms}, (4.10)

where "product terms" refers to terms that consist of products of combinations of the eigenvalues λ_i .

<u>Proposition</u>: Under the fine quantization condition, maximizing $det(\hat{J})$ is equivalent to minimizing of $tr(\mathbf{R})$.

<u>Proof</u>: First, if fine quantizers are used $(q_i^2 \text{ is very small})$, the reduction of the determinant of FIM will be small, which makes $\det(\mathbf{I} - \mathbf{R}) \rightarrow 1$, and all λ_i are very small. Thus the product terms in (4.10) are small compared to $\operatorname{tr}(\mathbf{R})$. In this point of view, maximizing $\det(\mathbf{I}_{2\times 2} - \mathbf{R})$ is equivalent to minimizing $\operatorname{tr}(\mathbf{R})$.

This proposition motivates using minimization of $tr(\mathbf{R})$ as an *approximating* alternative to maximization of $det(\hat{\mathbf{J}})$ even though this strictly holds only in the fine quantization case. We can further simplify $tr(\mathbf{R})$ as follows:

$$\operatorname{tr}(\mathbf{R}) = \operatorname{tr}\left(\left(\widetilde{\mathbf{G}}^{T}\widetilde{\mathbf{G}}\right)^{-1}\widetilde{\mathbf{G}}^{T}\operatorname{diag}\left\{\left[\frac{q_{0}^{2}}{\sigma^{2}+q_{0}^{2}} \quad \frac{q_{1}^{2}}{\sigma^{2}+q_{1}^{2}} \quad \cdots \quad \frac{q_{N-1}^{2}}{\sigma^{2}+q_{N-1}^{2}}\right]\right\}\widetilde{\mathbf{G}}\right)$$
$$= \operatorname{tr}\left(\widetilde{\mathbf{G}}\left(\widetilde{\mathbf{G}}^{T}\widetilde{\mathbf{G}}\right)^{-1}\widetilde{\mathbf{G}}^{T}\operatorname{diag}\left\{\left[\frac{q_{0}^{2}}{\sigma^{2}+q_{0}^{2}} \quad \frac{q_{1}^{2}}{\sigma^{2}+q_{1}^{2}} \quad \cdots \quad \frac{q_{N-1}^{2}}{\sigma^{2}+q_{N-1}^{2}}\right]\right\}\right)$$
$$(4.11)$$
$$= \sum_{i=0}^{N-1} v_{i} \frac{q_{1}^{2}}{\sigma^{2}+q_{1}^{2}}$$

where v_i is the *i*th element of the vector **v** given by $\mathbf{v} = \text{diag}\left(\widetilde{\mathbf{G}}(\widetilde{\mathbf{G}}^T\widetilde{\mathbf{G}})^{-1}\widetilde{\mathbf{G}}^T\right)$, which is called effective independence distribution (EID) in [76] and [77]. Thus, a compression algorithm for approximately maximizing the determinant of the FIM can be formulated as

$$\min_{\{b_i\}} \left\{ \sum_{i=0}^{N-1} v_i \frac{c \sigma_{s,i}^2 2^{-2b_i}}{\sigma^2 + c \sigma_{s,i}^2 2^{-2b_i}} \right\} \qquad \text{subject to} \qquad \sum_{i=0}^{N-1} b_i \le R , \tag{4.12}$$

which is an additive distortion measure that can be optimized effectively using the numerical Lagrange multiplier method of [33].

B. Weighted Trace of Fisher Information matrix

Although it involves the effect of the compression on all the elements of $\hat{\mathbf{J}}$, maximizing $det(\hat{\mathbf{J}})$ does not allow the possibility of setting an importance-weighting on the accuracy of the two parameter estimates because it is dependent on the product of the length of semi-axis, $\lambda_1 \lambda_2$ and there are cases where λ_1 is very large while λ_2 is very small. Sometimes, one may wish to favor the accuracy of one parameter at the expense of the others. This allows user-imposed trade-offs between parameters, which is important in multiple parameter estimation problems in that a user may need better accuracy on a subset of the parameters than the rest. Using the determinant, of course, can not fulfill this need, which makes the perimeter approach attractive. The perimeter of the ellipse is an alternative to the area, but it is quite complicated to compute exactly; however, we can use an approximation given by $P \approx C_P \sqrt{\lambda_1 + \lambda_2}$, where C_P is a constant; because the perimeter involves a sum of the eigenvalues it is clear that maximizing the perimeter is equivalent to maximizing tr{ $\hat{\mathbf{J}}$ } = $\hat{J}_{11} + \hat{J}_{22}$. Note that this may seem troubling because it only depends on the diagonal elements of J and does not capture the effect of the compression on the cross-terms (which control the tilt of the ellipse). However, it does allow importance-weighting on the accuracy of the two parameters: we can use as our distortion measure a weighted trace wtr{ \mathbf{J} } = $\alpha J_{11} + (1 - \alpha)J_{22}$, where α is an importancecontrolling parameter satisfying $0 \le \alpha \le 1$. Unlike the determinant, this has the very nice property of being an additive distortion function. However, an issue of concern is the effect that compression can have on the tilt of the ellipse, which is of most concern when the FIM ellipse is highly eccentric. The following theorem shows that this is not a serious issue because the post-compression FIM ellipse will always reside inside the original FIM ellipse; thus, for a highly eccentric original ellipse, compression is not able to greatly change the orientation of the ellipse. Thus, our goal is to seek an *operational* rate-distortion method that allocates bits within the above defined transform-coding framework to minimize the reduction in wtr{J} = $\alpha J_{11} + (1-\alpha)J_{22}$ with $0 \le \alpha \le 1$ while satisfying a budget on the total number of allocated bits.

<u>Theorem</u>: For the transform coding framework outlined above, the post-compression FIM $\hat{\mathbf{J}}$ given in (4.7) has an information ellipse $\mathbf{\theta}^T \hat{\mathbf{J}}^{-1} \mathbf{\theta} = \kappa$ that lies inside the original FIM ellipse $\mathbf{\theta}^T \mathbf{J}^{-1} \mathbf{\theta} = \kappa$.

<u>Proof</u>: First we prove that for any positive definite, complex-valued $N \times N$ matrix $\mathbf{A} = \mathbf{A}_R + j\mathbf{A}_I$ there is a relationship between the ellipsoid of \mathbf{A} in C^N and the corresponding ellipsoid of matrix \mathbf{A}_R in R^N . Let $\mathbf{x} = \mathbf{x}_r + j\mathbf{x}_i$ be a vector in the complex space. By using $\mathbf{A}_R^T = \mathbf{A}_R$ and $\mathbf{A}_I^T = -\mathbf{A}_I$ and simple manipulation we see that $\mathbf{x}^H \mathbf{A} \mathbf{x} =$

$$\mathbf{x}_{r}^{T}\mathbf{A}_{R}\mathbf{x}_{r} + \mathbf{x}_{i}^{T}\mathbf{A}_{R}\mathbf{x}_{i} + 2\mathbf{x}_{i}^{T}\mathbf{A}_{I}\mathbf{x}_{r} + j(\mathbf{x}_{r}^{T}\mathbf{A}_{I}\mathbf{x}_{r} + \mathbf{x}_{i}^{T}\mathbf{A}_{I}\mathbf{x}_{i}); \text{ but } \mathbf{x}_{r}^{T}\mathbf{A}_{I}\mathbf{x}_{r} = -\mathbf{x}_{r}^{T}\mathbf{A}_{I}\mathbf{x}_{r},$$

so $\mathbf{x}_r^T \mathbf{A}_I \mathbf{x}_r = 0$ and a similar argument gives $\mathbf{x}_i^T \mathbf{A}_I \mathbf{x}_i = 0$. Thus

$$\mathbf{x}^H \mathbf{A} \mathbf{x} = \mathbf{x}_r^T \mathbf{A}_R \mathbf{x}_r + \mathbf{x}_i^T \mathbf{A}_R \mathbf{x}_i + 2\mathbf{x}_i^T \mathbf{A}_I \mathbf{x}_r,$$

which describes the ellipse in the complex space. However, if we "slice" this ellipse through the planes where the imaginary parts are zero (i.e., restrict it to only the real axes) we get the ellipse defined in R^N by the real part of matrix **A**. This states that all we need to prove is that the C^N information ellipse for $\hat{\mathbf{J}}_c = \widetilde{\mathbf{G}}^H (\boldsymbol{\Sigma} + \mathbf{Q})^{-1} \widetilde{\mathbf{G}}$ is inside the C^N information ellipse for $\mathbf{J}_c = \widetilde{\mathbf{G}}^H \boldsymbol{\Sigma}^{-1} \widetilde{\mathbf{G}}$, where the subscript "*c*" indicates the related complex version of the FIM.

Let **V** be the matrix that simultaneously diagonalizes the positive-definite matrices Σ and **Q**, that is: $\mathbf{V}^{\mathbf{H}}\Sigma\mathbf{V} = \Lambda$ and $\mathbf{V}^{\mathbf{H}}\mathbf{Q}\mathbf{V} = \mathbf{I}$, where $\Lambda = \text{diag}\{\lambda_1, \lambda_1, \dots, \lambda_N\}$ are the generalized eigenvalues of Σ and **Q**. Then

$$\hat{\mathbf{J}}_{c} = \widetilde{\mathbf{G}}^{H} \mathbf{V} \operatorname{diag}([\frac{1}{1+\lambda_{1}}, \frac{1}{1+\lambda_{2}}, \cdots, \frac{1}{1+\lambda_{N}}]) \mathbf{V}^{H} \widetilde{\mathbf{G}}$$
(4.13)

and

$$\mathbf{J}_{c} = \widetilde{\mathbf{G}}^{H} \mathbf{V} \operatorname{diag}([\frac{1}{\lambda_{1}}, \frac{1}{\lambda_{2}}, \cdots, \frac{1}{\lambda_{N}}]) \mathbf{V}^{H} \widetilde{\mathbf{G}}$$
(4.14)

Clearly, the diagonal matrices in (4.13) and (4.14) defines ellipses that are aligned with the axes in a new coordinate system, and $\hat{\mathbf{J}}_c$ and \mathbf{J}_c are just these coordinate-aligned ellipses rotated and scaled into a new coordinate system. Thus, if the ellipse from the diagonal matrix in (4.13) is inside the ellipse from the diagonal matrix in (4.14) the proof is complete. This is true if $\lambda_i \ge 0$, and, as we show next, the generalized eigenvalue problem $\Sigma \mathbf{x} = \lambda \mathbf{Q} \mathbf{x}$ has positive eigenvalues when both Σ and \mathbf{Q} are positive definite (I call this the "2-pd" generalized eigenvalue problem).

Because **Q** is positive definite it can be written as $\mathbf{Q} = \mathbf{L}\mathbf{L}^{H}$ where **L** is nonsingular. Then the generalized eigenvalue problem can be written as $\mathbf{\Sigma}\mathbf{x} = \lambda \mathbf{L}\mathbf{L}^{H}\mathbf{x}$. Now inserting $\mathbf{I} = (\mathbf{L}^{H})^{-1}\mathbf{L}^{H}$ between $\mathbf{\Sigma}$ and \mathbf{x} leads to $\mathbf{K}\mathbf{\Sigma}\mathbf{K}^{H}\mathbf{y} = \lambda\mathbf{y}$, where we have defined $\mathbf{y} = \mathbf{L}^{H}\mathbf{x}$ and $\mathbf{K} = \mathbf{L}^{-1}$. Thus, this shows that the 2-pd generalized eigenvalue problem is equivalent to a standard eigenvalue problem of the matrix $\mathbf{K} \Sigma \mathbf{K}^{H}$. Now, note that this matrix is positive definite, so its eigenvalues (which are also the generalized eigenvalues) are positive. Noting that Σ and \mathbf{Q} are positive definite completes the proof.

Thus our proposed approach is to compress the data collected at sensor #1 so as to maintain the largest weighted trace of the data-computed FIM while meeting the constraint on the rate; that is,

$$\max_{B} \left[\alpha J_{11}(\hat{\chi}_{1}) + (1 - \alpha) J_{22}(\hat{\chi}_{1}) \right] \text{ subject to } \sum_{n=1}^{N} b_{i} \le R, \qquad (4.15)$$

where $J_{ii}(\hat{\chi}_1)$ is the *ii*th element of the data-computed FIM and the maximization is done over all allocation sets *B* that satisfy the rate constraint.

4.1.1 Application to Simultaneous Estimation

To provide a clear focus on the issues as well as to ensure maximal relevance of the results, joint TDOA/DOA and TDOA/FDOA applications are chosen to make our point; More discussion is put on TDOA/FDOA application because the ideas and procedures developed for TDOA/FDOA is quite general and are applicable to a broad range of scenarios of interest to the designer of sensor network, including joint TDOA/DOA application.

4.1.1.1 Joint TDOA and DOA Estimation:

Equations (3.28)-(3.31) and (3.44)-(3.46) provide the formulas for computing the FI for the TDOA only and DOA only and show the TDOA Fisher information depends on the quadratically-frequency-weighted DFT coefficients while the DOA Fisher

information depends on the SNR of the signal. Thanks to Parsaval theorem, maximizing SNR can be done in any ON transform, which gives us the flexibility to use DFT instead of wavelet transform for DOA estimate and to perform compression for joint TDOA and DOA using TDOA-optimal ON transform. DFT transform coding for DOA requires more bits be allocated to the DFT coefficients whose magnitudes are much higher than the noise floor. In the following, after we first extend the necessary signal model in Chapter 3 for the joint TDOA/DOA case, we use the theory and algorithms in Section 4.1 to achieve the desired trade-offs between TDOA/DOA estimates.

The continuous-time signal model for two passively-received complex baseband signals at sensors S_1 and S_2 having an unknown TDOA of τ and phase shift w_s due to DOA θ is given by

$$x_{1}[m] = s[m - (t_{0} + \tau/2)] \exp[-j(w_{0} - w_{s}/2)] + w_{1}[m] \quad m = 0, 1, ..., N-1$$

$$x_{2}[m] = s[m - (t_{0} - \tau/2)] \exp[-j(w_{0} + w_{s}/2)] + w_{2}[m] \quad m = 0, 1, ..., N-1,$$
(4.16)

where s(n) is the discrete complex-baseband signal, t_0 and w_0 are unknown nuisance parameters that can not be estimated, and $w_i(t)$ is complex-baseband bandlimited white Gaussian noise with variance of σ_i^2 , with σ_1^2 assumed known. Note that bearing θ must be estimated indirectly through the phase shift w_s by $\theta = \sin^{-1}\left(\frac{\lambda w_s}{2\pi d}\right)$. Taking the DFT of \mathbf{x}_1 , we have

$$X_{1}[k] \approx S[k] \exp[-j\frac{2\pi}{N}k(t_{0} + \tau/2)] \exp[-j(w_{0} - w_{s}/2)] + W_{1}[k]$$

$$k = -N/2, -N/2 + 1, \dots, N/2 - 1$$
(4.17)

where the S[k] are the DFT coefficients (for negative and positive frequencies) of the samples of signal s(t) and $W_1[k]$ are the DFT coefficients of the noise.

Each element of the 2×2 FIM for τ and w_s are

$$J_{\tau\tau} = 2 \sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right)^2 \frac{|X_1[k]|^2}{N\sigma^2},$$
(4.18)

$$J_{\theta\theta} = \left(\frac{\lambda}{2\pi d\cos\theta}\right)^2 \frac{1}{2} \sum_{k=-N/2}^{N/2-1} \frac{|X_1[k]|^2}{N\sigma^2}, \qquad (4.19)$$

$$J_{\tau\theta} = J_{\theta\tau} = \left(\frac{\lambda}{2\pi d\cos\theta}\right)_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right) \frac{\left|X_1[k]\right|^2}{N\sigma^2},\tag{4.20}$$

Hence the pre-quantization FIM can be written as

$$\mathbf{J} = \begin{bmatrix} 2\sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right)^2 \frac{\left|X_1[k]\right|^2}{N\sigma^2} & \left(\frac{\lambda}{2\pi d\cos\theta}\right)_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right) \frac{\left|X_1[k]\right|^2}{N\sigma^2} \\ \left(\frac{\lambda}{2\pi d\cos\theta}\right)_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right) \frac{\left|X_1[k]\right|^2}{N\sigma^2} & \left(\frac{\lambda}{2\pi d\cos\theta}\right)^2 \frac{1}{2} \sum_{k=-N/2}^{N/2-1} \frac{\left|X_1[k]\right|^2}{N\sigma^2} \end{bmatrix}.$$
(4.21)

Distortion measures of FIM of \mathbf{J} with respect to the quantization can be either one of the following:

(1) **Optimize** trace(**J**)

Mathematically, we have

$$\operatorname{tr}(\hat{\mathbf{J}}) = 2\sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right)^2 \frac{\left|X_1[k]\right|^2}{N\sigma^2 + q_k^2} + \left(\frac{\lambda}{2\pi d\cos\theta}\right)^2 \frac{1}{2} \sum_{k=-N/2}^{N/2-1} \frac{\left|X_1[k]\right|^2}{N\sigma^2 + q_k^2} \quad (4.22)$$

Letting $\gamma = (\lambda / 4\pi d \cos \theta)^2$, (4.22) can be transformed to

$$\operatorname{tr}(\hat{\mathbf{J}}) = \sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right)^2 \frac{\left|X_1[k]\right|^2}{N\sigma^2 + q_k^2} + \gamma \sum_{k=-N/2}^{N/2-1} \frac{\left|X_1[k]\right|^2}{N\sigma^2 + q_k^2}, \qquad (4.23)$$

From (4.23), we can see that the value of γ varies as the different true values of θ and the choice of λ and d, and the value of γ determines the relative importance of the two summed terms for tr($\hat{\mathbf{J}}$). If we have a priori information of $\hat{\gamma}$ or we can

estimate $\hat{\gamma}$ roughly, then (4.23) can be transformed into the weighted-trace-based distortion measure as

$$\operatorname{tr}(\hat{\mathbf{J}}) = (1 - \alpha) \sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right)^2 \frac{\left|X_1[k]\right|^2}{N\sigma^2 + q_k^2} + \alpha \sum_{k=-N/2}^{N/2-1} \frac{\left|X_1[k]\right|^2}{N\sigma^2 + q_k^2}.$$
(4.24)

(2) Optimizing det(J)

Factoring out the unknown bearing parameter θ from (4.21), we have

$$\mathbf{J} = \begin{bmatrix} 1 & 0 \\ 0 & \left(\frac{\lambda}{2\pi d \cos\theta}\right) \end{bmatrix} \begin{bmatrix} 2\sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right)^2 \frac{|X_1[k]|^2}{N\sigma^2} & \sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right) \frac{|X_1[k]|^2}{N\sigma^2} \\ \sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right) \frac{|X_1[k]|^2}{N\sigma^2} & \frac{1}{2}\sum_{k=-N/2}^{N/2-1} \frac{|X_1[k]|^2}{N\sigma^2} \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & \left(\frac{\lambda}{2\pi d \cos\theta}\right) \end{bmatrix} \\ \mathbf{J}$$

(4.25)

Maximizing $det(\mathbf{J})$ is equivalent to maximizing

$$\det(\mathbf{F}) = \begin{vmatrix} 2\sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right)^2 \frac{|X_1[k]|^2}{N\sigma^2 + q_k^2} & \sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right) \frac{|X_1[k]|^2}{N\sigma^2 + q_k^2} \\ \sum_{k=-N/2}^{N/2-1} \left(\frac{\pi k}{N}\right) \frac{|X_1[k]|^2}{N\sigma^2 + q_k^2} & \frac{1}{2} \sum_{k=-N/2}^{N/2-1} \frac{|X_1[k]|^2}{N\sigma^2 + q_k^2} \end{vmatrix},$$
(4.26)

The property det(AB) = det(A)det(B) is used to get (4.26) from (4.25). We can maximize

(4.26) according to (4.9) with vector **G** in (4.5)-(4.9) as

$$\mathbf{G} = \begin{bmatrix} \frac{\partial \mathbf{s}_1}{\partial \Delta} & \frac{\partial \mathbf{s}_1}{\partial w_s} \end{bmatrix} = \begin{bmatrix} \vdots & \vdots \\ f_k S_1[n,k] & S_1[n,k] \\ \vdots & \vdots \end{bmatrix}$$

4.1.1.2 Simulations for Joint TDOA and DOA Estimation

The signal we use is a complex baseband FM signal with a pseudo-random modulating signal; a sample spectrum is shown in Figure 4.1. We assume the same SNR at S_1 and S_2 ; Both are set at 15 dB. For compression ratio (CR) values of interest, we



Figure 4.1: The spectrum of a typical FM signal used in the simulations.

evaluated the estimation errors of TDOA and DOA over 400 Monte Carlo runs. As usual, to focus on the capabilities of the transform coding and to compare compression algorithms under different designs, we performed no entropy coding, which would likely provide further improvement in the CR with no further accuracy degradation.

Simulation results are shown in Figure 4.2 and Figure 4.3. In Figure 4.2 we see the inherent tradeoff that is controlled by the choice of α , whose value controls whether the algorithm favors TDOA accuracy, DOA accuracy, or both. It is clear that optimization of the determinant of FIM can balance both TDOA accuracy and DOA accuracy and tend to approach around the operating point nearest to no compression case. Figure 4.3 illustrates how changing the compression ratio affects the TDOA and DOA accuracy under several scenarios. The advantages of using our distortion measures rather than a MSE measure are made clear from these figures.

4.1.1.3 Joint TDOA and FDOA Estimation

The continuous-time signal model for two passively-received complex baseband signals at sensors S_1 and S_2 having an unknown TDOA of τ and FDOA of v is given by



Figure 4.2: Trade-off between TDOA and DOA accuracies as α is varied for $SNR_1 = 15$ dB & $SNR_2 = 15$ dB and compression ratio (a) 3.5:1 and (b) 6:1.





Figure 4.3 Effect of compression ratio on (a) TDOA and (b) DOA performance.
$$x_{1}(t) = s(t - (t_{0} + \tau/2))e^{j(v_{0} + \nu/2)t} + w_{1}(t)$$

$$x_{2}(t) = s(t - (t_{0} - \tau/2))e^{j(v_{0} + \nu/2)t} + w_{2}(t)$$
(4.27)

The TDOA Fisher information measure is

$$J_{11} = \frac{2\pi^2}{N^2 \sigma_1^2} \sum_{m=-N/2}^{N/2} m^2 \left| S_1[m] \right|^2 , \qquad (4.28)$$

where $S_1[m]$ is the DFT of the signal at sensor S_1 and the FDOA Fisher information measure is

$$J_{22} = \frac{2\pi^2}{\sigma_1^2} \sum_{n=-N/2}^{N/2} n^2 |s_1[n]|^2 , \qquad (4.29)$$

where $s_1[n]$ is the signal data at sensor S_1 .

Equations (4.28) and (4.29) provide the formulas for computing the FI for the TDOA only and FDOA only. Similar to joint TDOA/DOA estimation, compression requirements for TDOA and FDOA are also conflicting because the TDOA Fisher information depends on the DFT coefficients (with the DFT frequencies running over both negative and positive frequencies) while the FDOA Fisher information depends on the signal samples themselves. However, unlike the joint TDOA/DOA estimation where the compression can be performed using the TDOA-only ON transform, the optimal ON transform for TDOA is the DFT whereas optimal transform for FDOA is the identity matrix. Each DFT coefficient is the combined result of all the time samples (from the beginning to the end), similarly, one time sample is the combined result of all the DFT coefficients (from the negative frequency to positive frequency). Even worse, the off-diagonal terms of the FIM depend on both frequency-domain characteristics and time-domain characteristics. It is difficult to calculate and optimize the post-compression values of all the components of the FIM jointly unless we transform the data $x_1[n]$ into a time-frequency representation where frequency resolution and time resolution are jointly provided. Wavelet packet transform is such choice because it provides not only the time-frequency resolution that is needed for our joint TODA/FDOA, but also provides the energy compactness which is important for the compression efficiency.

Given an orthonormal wavelet packet basis set $\{\psi_n\}$ with coefficients $\{c_n\}$ for the data vector \mathbf{x}_1 , we wish to select a subset Ω of coefficient indices and an allocation of bits $B = \{b_i | i \in \Omega\}$ to those selected coefficients such that the selected/quantized signal $\widetilde{\mathbf{x}}_1 = \sum_{i \in \Omega} \widetilde{c}_i \psi_i$, where $\{\widetilde{c}_i | i \in \Omega\}$ are the quantized version of the selected coefficients. The selection/quantization is to be done so as to maximize the post-compression values of wtr $\{\mathbf{J}\} = \alpha J_{11} + (1 - \alpha)J_{22}$ and det $\{\mathbf{J}\}$ while meeting a constraint on the total number of bits.

In principle, it is possible to write a mathematical equation for the coefficients and then consider the derivatives with respect to parameters, as needed for the FIM. This, however, quickly becomes intractable. Thus, we resort to exploiting the relationships among the wavelet coefficients and DFT coefficients and original time-domain samples of \mathbf{x}_1 . Although it is impossible to relate each coefficient to a single DFT coefficient or single time-domain samples, as shown in [81], [82], and [83], all wavelet coefficients in the same band have the tendency to be associated only with a certain frequency range and a block of adjacent wavelet coefficient in the same band have a tendency to correspond to the same time range. Therefore, we can divide the wavelet coefficient into small blocks, and give each block a frequency weight and time weight according to the time-frequency location characteristic of wavelets. Besides providing us the reasonable time-frequency resolution for compression for joint TDOA/FDOA estimation, operations on blocks of coefficients, instead of coding coefficients individually, brings us computational efficiency as well. Now consider that the wavelet packet coefficients c_i are grouped into M blocks, where each block contains coefficients at the same frequency and over a short contiguous temporal range. The diagonal element of the FIM can be computed (up to a multiplicative factor) as:

$$J_{11} = \sum_{j=1}^{M} \left(\frac{f_j^2 \sum_{i \in \{j \text{ block}\}} |c_i|^2}{\sigma_1^2 + q_j^2} \right)$$
(4.30)

$$J_{22} = \sum_{j=1}^{M} \left(\frac{t_j^2 \sum_{i \in \{j \text{ block}\}} |c_i|^2}{\sigma_1^2 + q_j^2} \right).$$
(4.31)

where *j* is the block index, f_j and t_j are the frequency and central time, respectively, of the *j*th block. Based on (4.30) and (4.31), the distortion measure for joint TDOA/FDOA estimate can be determined in the following:

Weighted Trace Optimization:

The weighted-trace-based TDOA/FDOA distortion measure is then

$$wtr(\widetilde{J}(\alpha)) = \alpha \sum_{j=1}^{M} \left(\frac{f_{j}^{2} \sum_{i \in \{j \text{ block}\}} |c_{i}|^{2}}{\sigma_{1}^{2} + q_{j}^{2}} \right) + (1 - \alpha) \sum_{j}^{M} \left(\frac{t_{j}^{2} \sum_{i \in \{j \text{ block}\}} |c_{i}|^{2}}{\sigma_{1}^{2} + q_{j}^{2}} \right)$$
(4.32)

Bits are allocated to the coefficients using the method of [33] to maximize (4.32) for a given rate constraint. To compare the proposed scheme with a traditional scheme, we also

allocated bits to the wavelet packet blocks to minimize MSE (2.36) under the bit constraint.

Although the developed scheme is applicable to all varieties of signals, we use a linear FM radar signal [87] to illustrate the method. A 3-level wavelet packet transform is performed and 8 subbands are produced; each subband is partitioned into 8 blocks. Moreover, to focus attention on the lossy compression performance, no entropy coding is applied after quantization.

Simulation results are shown in Figure 4.4 and Figure 4.5. In Figure 4.4 we see the inherent trade-off that is controlled by the choice of α , whose value controls whether the algorithm favors TDOA accuracy, FDOA accuracy, or balances them to achieve the closest operation to the no compression case. Figure 4.5 illustrates how changing the compression ratio affects the TDOA and FDOA accuracy under several scenarios: the results labeled "Goal Attained" illustrate the performance when the impact on TDOA and FDOA is balanced. In these figures the advantage of using our distortion measures is made clear.

Determinant Optimization:

Unlike the weighted trace, the other approach, maximizing the determinant, does not provide us the flexibility to weight the relative importance between TDOA and FDOA However, it takes into account the off-diagonal elements of **J** that the weighted trace ignores and has the advantage of maximizing the size of the FIM ellipse. It is possible for determinant optimization to automatically achieve the best tradeoff between TDOA accuracy and FDOA accuracy without heuristically deciding the weight α in the weighted trace approach. Correspondingly, the vector **G** needed to be used in (4.5)-(4.9)



Figure 4.4: Trade-off between TDOA and FDOA accuracies as α is varied for compression ratio 3:1 and $SNR_1 = 15 \text{ dB} \& SNR_2 = 15 \text{ dB}$; symbol \Box denotes the operational point (α =0.5) closest to that without compression.



Figure 4.5 (a)



Figure 4.5: Effect of compression ratio on (a) TDOA and (b) FDOA performance. A comparison is also made to the case of simply sending less data ("Length Reduced") rather than compressing the data.



Figure 4.6: Comparison between the determinant optimization method ('Area') and weighted trace method ('Premeter') and MSE for compression ratio 3:1 and $SNR_1 = 15$ dB & $SNR_2 = 15$ dB;

can be determined by

$$\mathbf{G} = \begin{bmatrix} \frac{\partial \mathbf{s}_1}{\partial \Delta} & \frac{\partial \mathbf{s}_1}{\partial \nu} \end{bmatrix} = \begin{bmatrix} \vdots & \vdots \\ f_j c_j [i] & t_j c_j [i] \\ \vdots & \vdots \end{bmatrix},$$

where $c_i[i]$ denote the wavelet coefficients residing in the jth block.

Simulation results are shown in Figure 4.6, where the weighted trace method is also given for comparison; as expected, the operating point obtained from the determinant method is the point to achieve the pareto optimal point (the closest operation to the no compression case).

4.1.2 Modified Distortion Measures to Include Off-Diagonals

For the TDOA/FDOA-based location problem the estimation of TDOA/FDOA is only a preliminary task; the real task is estimation of the geo-location of the emitter and that, as is shown below, depends on the off-diagonals of the TDOA/FDOA FIM. For simplicity we consider only 2-D geolocation.

The 2-D location estimate of the emitter $\hat{\mathbf{x}} = [\hat{x}_e, \hat{y}_e]^T$ is sought and *K* pairs of sensors are used. Each pair of sensors is able to measure both TDOA and FDOA so that we have the measurements

$$\tau_i = f_{TDOA,i}(\mathbf{x}) + w_{\tau_i}$$

$$v_i = f_{FDOA,i}(\mathbf{x}) + w_{v_i} \qquad i = 1, 2, \dots, K$$
(4.33)

where w_{τ_i} and w_{v_i} are the random TDOA and FDOA measurement errors at the *i*th pair of sensors, respectively. Any two measurements at a pair of sensor can be written as a single equation for a 2-dimensional column vector.

$$\begin{bmatrix} \tau_i \\ v_i \end{bmatrix} = \begin{bmatrix} f_{TDOA,i}(\mathbf{x}) \\ f_{FDOA,i}(\mathbf{x}) \end{bmatrix} + \begin{bmatrix} w_{\tau_i} \\ w_{v_i} \end{bmatrix}, \quad i = 1, 2, \dots, K$$
(4.34)

Because the TDOA/FDOA estimates are obtained using the maximum likelihood (ML) estimator of cross correlation, the asymptotic properties of ML estimators (2.16) says that $[w_{n_i}, w_{v_i}]$ follows the Gaussian distribution

$$\begin{bmatrix} w_{\tau_i} \\ w_{\nu_i} \end{bmatrix} \sim \mathcal{N}(0, \mathbf{J}_i^{-1}), \qquad \mathbf{\theta} = [\tau_i, \nu_i],$$

where J_i is the Fisher information matrix for the *i*th TDOA/FDOA pair. As we know, J_i is dependent only on the signal and SNR according to

$$\mathbf{J}_{i} = 2 \operatorname{Re}\left[\frac{\partial \mathbf{s}^{H}(\mathbf{\theta}_{i})}{\partial \mathbf{\theta}_{i}} \boldsymbol{\Sigma}_{i}^{-1} \frac{\partial \mathbf{s}(\mathbf{\theta}_{i})}{\partial \mathbf{\theta}_{i}}\right],$$

where Σ_i is the covariance of the AWGN at the sensors. Letting

$$\mathbf{G}_{i} = \begin{bmatrix} \frac{\partial f_{TDOA,i}}{\partial x_{e}} & \frac{\partial f_{TDOA,i}}{\partial y_{e}} \\ \frac{\partial f_{FDOA,i}}{\partial x_{e}} & \frac{\partial f_{FDOA,i}}{\partial y_{e}} \end{bmatrix}$$

it is known that the FIM of the geo-location estimation is given by (2.8)

$$\mathbf{J}_{geo} = \begin{bmatrix} \mathbf{G}_{1}^{\mathrm{T}}, \cdots, \mathbf{G}_{J}^{\mathrm{T}} \end{bmatrix} \begin{bmatrix} \mathbf{J}_{1} & \mathbf{0} & \cdots & \mathbf{0} \\ \mathbf{0} & \mathbf{J}_{2} & \ddots & \vdots \\ \vdots & \ddots & \ddots & \mathbf{0} \\ \mathbf{0} & \cdots & \mathbf{0} & \mathbf{J}_{J} \end{bmatrix} \begin{bmatrix} \mathbf{G}_{1} \\ \vdots \\ \mathbf{G}_{J} \end{bmatrix}$$

$$= \sum_{i=1}^{J} \mathbf{G}_{i}^{\mathrm{T}} \mathbf{J}_{i} \mathbf{G}_{i} .$$

$$(4.35)$$

Equation (4.35) conveys two important information: (i) the network-wide optimal compression strategy can be achieved in a decentralized way such that each pair of sensors can optimize the compression based on their own G_i without knowledge of the

operations of other pairs of sensors; (ii) besides the TDOA/FDOA accuracies, the location estimation accuracy also strongly depends on the geometry of the emitter and the sensors [66]. For the 2-D X-Y location case, the 2×2 location error covariance matrix can be replaced by $\mathbf{P} = \mathbf{J}_{geo}^{-1}$ [66]. Diagonalizing \mathbf{P}^{-1} (or \mathbf{J}_{geo}) using $\mathbf{A}^{T}\mathbf{P}^{-1}\mathbf{A} = diag\{\lambda_{1}^{-1}, \lambda_{2}^{-1}\}$ and letting $\boldsymbol{\zeta} = \mathbf{A}^{T}(\boldsymbol{\xi} - \mathbf{m})$, where **m** is the real location of target, leads to

$$R = \{ \boldsymbol{\xi} : \sum_{i=1}^{2} \frac{\zeta_i^2}{\lambda_i} \le \boldsymbol{\kappa} \}, \qquad (4.36)$$

where region *R* is the interior of the error ellipse with the *i*th principal axes length of $2(\kappa \lambda_i)^{\frac{1}{2}}$.

In most sensor-target geometries, the relative importance of TDOA and FDOA is set by the target-sensors geometry. Figure 4.4 and Figure 4.6 show how α controls the compression trade-off between the TDOA and FDOA accuracies achieved by our method. However, to choose a proper value of α requires knowledge of the target's location – which, unfortunately, is precisely what we are trying to estimate. However, we can first send a small amount of data – enough to *roughly* determine the geometry; then that rough geometry can be used to estimate an appropriate α value that would be fed back to the compressing sensor and used to compress the remainder of the signal. In fact, this can even be done to provide repeated updates of α as more data is compressed and sent; this corresponds to updating the operating point along the curve shown in Figure 4.4 and Figure 4.6. This directly leads to the following simple scheme (not the only one) where the compression can be adapted to the geometry. The scheme is shown in Figure 4.7, where circled numbers correspond to the actions of the steps described below.

Step 1: The central node determines the compression ratio and its associated operational compression points based on the requirement of energy consumption and latency. Then it randomly picks J pairs of nodes to get J measurements of TDOA and FDOA. In the beginning, one of the nodes (*i*,1), $1 \le i \le J$ sends a small length of its data to the other (*i*,2), where rough measurements of TDOAs $\{\tilde{\tau}_{d,i}\}$ and FDOAs $\{\tilde{\nu}_{d,i}\}$ are measured.

Step 2: $\{\tilde{\tau}_{d,i}\}$ and $\{\tilde{v}_{d,i}\}$ are sent to the central node, where the rough prediction of the geometry matrix *G* is estimated.

Step 3: In term of the estimated geometry matrix **G**, the central node determines the optimal α and informs the nodes (*i*,1) about this choice and the compress ratio.

Step 4: The nodes (i,1) compress the data according the compression ratio and the chosen α and send the compressed version of its sensed data to the nodes (i,2).



Figure 4.7: Geometry Adaptive TDOA/FDOA System Scheme.

Step 5: The nodes (*i*,2) measure the $\{\tau_{d,i}\}$ and $\{v_{d,i}\}$ with the received data from node (*i*,1) in Step 1 and decompressed data in Step 4. The resulting $\{\tau_{d,i}\}$ and $\{v_{d,i}\}$ are sent to the central node, where the final location of the target is determined.

Simulations were performed as follows: A target can appear anywhere in a 2-D 5x5 km² area, within which a large number of sensors is spread. Four sensors are randomly picked to form two pairs to estimate TDOA/FDOA. 256 samples are sent in Step 1 to roughly estimate the geometry. Finally, 4096 samples are compressed and shared between each pair. For comparison, the 4096 data samples are compressed using the MSE measure. In the simulation, we use the ratio of the area of circular error probable (CEP) [66] with compression to that without compression minus 1 (which corresponds to the relative increase in CEP) as the metric. CEP is estimated by $0.75\sqrt{\lambda_1 + \lambda_2}$ [66]. There were 2000 Monte Carlo runs performed for each compression ratio and the average values are shown in Figure 4.8. Over the range of compression ratio values considered, the results in Figure 4.8 show a roughly 5x improvement in CEP for our method relative to one that uses the wavelet packet transform but optimized according to MSE.

Actually, this geometry adaptive algorithm is optimized with respect to the FIM of $\{\tau_i\}$ and $\{v_i\}$ given a chosen α for each pair of sensor. Although the significant improvement is achieved over MSE as shown in Figure 4.8, there is a step left to claim that it is optimal with respect to FIM of x_e and y_e . As shown in (4.35). The correct usage of the viable "Perimeter of the FIM" idea to optimizing compression with respect geo-location to should be to compress to minimize the trace of $E\{(\hat{x}_e - x_e) + (\hat{y}_e - y_e)\} = tr(\mathbf{P}) = tr(\mathbf{J}_{geo}^{-1});$ as before, instead of minimizing $tr(\mathbf{J}_{geo}^{-1})$, we



Figure 4.8: Effect of compression ratio on CEP.

resort to maximizing tr(\mathbf{J}_{geo}) instead because of their close relationship (4.2). However, the trace of FIM \mathbf{J}_{geo} depends on the off-diagonal elements of the TDOA/FDOA FIM as the following shows. Let for notational purposes

$$\mathbf{G}_{i} = \begin{bmatrix} G_{i,11} & G_{i,12} \\ G_{i,21} & G_{i,22} \end{bmatrix} \qquad \mathbf{J}_{i} = \begin{bmatrix} J_{i,11} & J_{i,12} \\ J_{i,21} & J_{i,22} \end{bmatrix}$$

then

$$\operatorname{tr}\left\{\mathbf{J}_{geo}\right\} = \sum_{i=1}^{J} \left(G_{i,11}^{2} + G_{i,12}^{2}\right) J_{i,11} + \left(G_{i,21}^{2} + G_{i,22}^{2}\right) J_{i,22} + 2\left(G_{i,11}G_{i,21} + G_{i,12}G_{i,22}\right) J_{i,12}, \quad (4.37)$$

where the first two terms in the numerator of each term of the sum depend on the diagonal elements of the TDOA/FDOA FIM but the last term depends on the off-

diagonal elements. If we assume that TDOA and FDOA are uncoupled, i.e., $J_{i,12} = 0$, it is obvious that the algorithm is optimal and α can be set equal to $(G_{i,11}^2 + G_{i,12}^2)/(G_{i,11}^2 + G_{i,12}^2 + G_{i,21}^2 + G_{i,22}^2)$. However, as shown in Appendix E, the assumption that TDOA and FDOA are uncorrelated is generally not met. And although the above algorithm provides significant improvement over MSE, there might be further potential to go further by maximizing the trace of the geo-location FIM. In doing so, we need to consider the off-diagonal elements of the TDOA/FDOA FIM and we need to modify the distortion measure which should be based on (4.37), and that requires knowledge of the FIM cross-terms for the TDOA/FDOA.

We can conjecture that the cross-terms can be numerically approximated as

$$J_{12} = J_{21} \approx \sum_{j=1}^{M} \left(\frac{t_j f_j \sum_{i \in \{j \text{ block}\}} |c_i|^2}{\sigma_1^2 + q_j^2} \right)$$
(4.38)

However, (4.38) can not be able to compute the off-diagonal terms J_{12} and J_{21} correctly by using the wavelet packet transform because each channel in this wavelet packet transform has both positive and negative frequency content (see Appendix F where such filter banks are called "two-sided" filter banks). This is not a problem for evaluating the weighted trace of FIM because of the squared f_k^2 and t_n^2 in (4.30) and (4.31). But for (4.37), with its $t_n f_k$ term, it is crucial that we have a so-called "one-sided" filter bank that has individual positive and negative frequency channels; such filter banks have complexvalued impulse responses for the channels. After a thorough search of the literature it was determined that no directly usable results were readily available for the design of onesided, complex-valued, orthogonal, perfect-reconstruction filter banks; Focusing on showing that (4.37) indeed provides the optimal solutions and how (4.38) fit into the optimization of (4.37) we resorted to a "quick and dirty" method to develop such filters– the result is given in Appendix F. The advantage of this approach is that existing routines for designing and implementing orthogonal PR filter banks can be used with minor modifications.

To test this filter bank we needed a signal that would impinge on each channel but for which it would still be possible to *analytically* compute the FIM elements. Then we could apply the filter bank and equations (4.30) (4.31) and (4.38) numerically compute the FIM elements and then compare the results to the analytically derived theoretical results for this signal. A good choice for this signal is a linear chirp signal. The derivation of the FIM for a chirp signal is given in Appendix E and the result is

$$\mathbf{J} = \left(2T^2 \sum_{n} n^2\right) \begin{bmatrix} \alpha^2 & \alpha \\ \alpha & 1 \end{bmatrix}.$$
 (4.39)

where α is the sweep rate of the chirp signal. The comparison of the numerical and theoretical results for the case of using a complex filter bank are shown in Figure E.1, where it is seen that the numerical results are not very accurate: in the top plot the numerical results in red should be a constant, in the middle plot the numerical results in red should follow the blue curve of α^2 , and in the bottom plot the numerical results in red should follow the blue curve of α . This was disappointing and distressing: if we are unable to accurately assess the FIM elements we can not use them to drive the bit allocation process.

To demonstrate that it *is* at least possible to evaluate the FIM from signal data, the short-time Fourier transform (STFT) was investigated as a replacement for the complex-



Figure 4.9: Our compression framework using parallel auxiliary STFT processing to evaluate the FIM elements.

valued orthogonal PR filter bank. The results obtained using STFT are given in Figure E.2; notice that the numerical results nearly perfectly overlay the theoretical (only "nearly" because there are some slight differences when the plots are zoomed). Although the STFT enables quite accurate evaluation of the FIM, the STFT is not well-suited for compression due to the fact that it is not an orthogonal representation. Thus, we could use it simply as an auxiliary parallel mechanism to allow evaluation of the FIM for bit allocation to the wavelet packet transformed coefficients, as shown in Figure 4.9. Each sample coming out of the filter bank represents a known time-frequency region and the FIM elements can be evaluated over this time-frequency region using the STFT to provide the FIM evaluation for the filter bank sample. This is very similar to how audio compression schemes like the mp3 standard use a filter bank to create the values that get allocated to the various filter bank samples.

4.1.3 Simulations of Compression Using Parallel Auxiliary STFT Processing

The idea and implementation of using parallel auxiliary STFT processing shown in Figure 4.9 will be further exploited in this section through simulations. Two typical geometries of the two pair of sensors and emitters are chosen to illustrate the improvement obtained by accurate assessment of FIM. These two geometries are shown in Figure 4.10 (a) and Figure 4.11 (a), respectively. Figure 4.10 (a) corresponds to the case when TDOA and FDOA are both important for the location of emitter, while Figure 4.11 (a) corresponds to the case when FDOA is more important for one pair of sensor while both TDOA and FDOA are important to the other pair. The Jacobian matrix G_1 and G_2 in (4.36) for Figure 4.10 are²:

$$\mathbf{G}_{1} = \begin{bmatrix} 1.8314 & 1.5157 \\ 2.1728 & 0.8425 \end{bmatrix} \text{ and } \mathbf{G}_{2} = \begin{bmatrix} -3.255 & 2.4777 \\ 2.1224 & 2.6003 \end{bmatrix}.$$

Similarly, the Jacobian matrix G_1 and G_2 in (4.36) for Figure 4.11 are

$$\mathbf{G}_{1} = \begin{bmatrix} 0.3372 & 0.4077 \\ -1.1369 & -0.3697 \end{bmatrix} \text{ and } \mathbf{G}_{2} = \begin{bmatrix} -0.9176 & -2.1011 \\ 1.3836 & -1.6357 \end{bmatrix}$$

Simulation results are shown in 4.10 (b) and 4.11 (b), where the trace of the practical error covariance matrix of location is used to compare different algorithms under investigation. STFT-WP FIM is the wavelet packet transform coding that uses parallel auxiliary STFT to evaluate the FIM while WP FIM is the wavelet packet transform coding that directly use equations (4.30), (4.31) and (4.3) to compute the elements of the FIM. Clearly, STFT-WP FIM is slightly better than WP FIM while both achived

² (\mathbf{G}_1 and \mathbf{G}_2 incorporate the mismatched units factor (TDOA error is in the unit of ns and FDOA error is in the unit of Hz) or other factors determining the functions in (4.33). (See [66] and [84] for details)

significant improvement over the wavelet transform coding based on MSE. Results marked wtr FIM represent the performance of the wavelet packet transform coding that is based on the weighted trace distortion measure in (4.32). The wtr FIM ignores the off-diagonal element $J_{i,12}$ and chooses α as $(G_{i,11}^2 + G_{i,12}^2)/(G_{i,11}^2 + G_{i,12}^2 + G_{i,21}^2 + G_{i,22}^2)$. According to the Schwarz inequality, $(G_{i,11}^2 + G_{i,12}^2)J_{i,11} + (G_{i,21}^2 + G_{i,22}^2)J_{i,22}$ is greater than or equal to $2(G_{i,11}G_{i,21} + G_{i,12}G_{i,22})J_{i,12}$, which means that the on-diagonal elements of J_{geo} dominant the off-diagonal elements and it might be good enough to use only wtr FIM because it is much simpler than STFT-WP FIM. This is verified by Figure 4.10. (b) and Figure 4.11. (b) where wtr FIM is very close to the optimal performance represented by STFT-WP FIM.

4.2 Data Compression for Parameter Estimation with Detection

There might be a scenario where sensor S_2 uses its own data and the decompressed data from S_1 to first detect the presence of a signal and then estimates the parameter of interest. We can derive the Chernoff-distance based distortion. However, based on the framework we used for multiple estimation tasks, it is not difficult to derive the following Fisher-information-Matrix-Chernoff-Distance-based distortion measure to guide the compression algorithm:

$$\max_{B} \left[\beta \times \psi(\mathbf{J}(\mathbf{x}_{1})) + (1 - \beta) \times \mu_{s}(\mathbf{x}_{1}) \right] \quad subject \ to \ \sum_{n=1}^{N} b_{i} \le R \quad , \tag{4.40}$$

N 7

where ψ represents any form of the FIM measures in 4.1 and β is a parameter used to control the relative importance of estimation accuracy and detection error.



Figure 4.10 (b)

Figure 4.10, (a) Geometry among the two pair of sensors and emitter, "*" represents the emitter and "o" represent the sensors, and "---" denotes the wireless communication channel between each pair of sensors. (b) The comparison of difference compression algorithms on the trace of error covariance of emitter

location.



Figure 4.11, (a) Geometry among the two pair of sensors and emitter, "*" represents the emitter and "o" represent the sensors, and "---" denotes the wireless communication channel between each pair of sensors.
(b) The comparison of difference compression algorithms on the trace of error covariance of emitter location.

For example, consider detecting the presence of a signal and then estimating TDOA between a pair of sensors in the presence of AWGN. The distortion measure for detection happens to coincide with the distortion measure for DOA. In other words, the compression algorithm for DOA should be almost identical to that for detection except for slight differences. Therefore the simulations used for testing the joint TDOA/DOA application can be applied to the joint TDOA and detection without any changes except that we replace weight α with β . Due to this reason and the reason that the usage of (4.39) will be further discussed later in the compression for sequential inference task application, we will postpone the simulations until the next chapter.

CHAPTER 5

Data Compression for Sequential Task of Multiple Inference Quantities

Recall the geometry-adaptive compression algorithm that we developed for joint TDOA/FDOA estimation, it requires that either a small amount of initial data or very coarse compressed data to be shared first to get a rough estimate of the emitter's location, if the rough estimate is inside the predicted error surface, it is possible to determine the proper TDOA/FDOA trade-off for compressing and sharing the remaining data and refining the rough estimate to a much more accurate estimate that satisfies the requirement of the sensor system. On the other hand, if the rough estimate is far away from our projected error surface, the fusion sensor can determine that the sensor has not enough valuable information and no further sharing is needed. A similar scenario is the sequential detect-then-estimate problem, where some compressed data is sent for the detection task; once detected, the fusion center will require more data sharing among the sensors to complete some estimation tasks. From these two examples, we can see that in sensor networks, the inference tasks, rather than occurring simultaneously, are more likely to happen sequentially for achieving the maximal information gain under the energy and time constraint. Generally, the carrying out of major inferences of interest could be preceded by conducting some "preliminary" inference which offers the central information fusion center rough information about the process under investigation, since in the complete absence of prior information, global optimization of the inference is impossible. These preliminary inference results can be either deciding if the target is present or roughly estimating of the geometry among the sensors and target. Then, after estimates of the parameters are found, it is necessary to verify whether or not the accuracy of the ultimate inference task at hand is sufficient, then depending on the circumstances of sensor networks, the information fusion center must either stop or design supplementary inference tasks for estimating the entire collection of parameters or some group of them which are of more interest. As we know, the compression requirements for the different inference tasks at each sequential stage are different and, usually, conflicting. Therefore, optimized compression to handle sequential inference tasks requires what we call "task-embedded compression": the transmitting sensor constructs the optimal task-embedded data stream to send only the data needed to supplement the already-delivered data for optimal processing for the current task at each stage. Although task-embedded compression is similar to the traditional embedded coding compression technique, they are quiet different in principle since the data stream of traditional embedded coding is constructed based on a single task (reconstruction) and cannot be adapted for different requirements for different tasks. In the following, we will first use a simple detection-then-TDOA sequential inference problem to illustrate our task-embedded compression approach, then generalization to the general case is given.

5.1 Simple Detection-Then-TDOA Problem

Consider the scenario where multiple sensors are deployed to detect and then locate RF emitters (e.g., communication or radar transmitters). At first the sensors would share their collected data for the purpose of detecting when they have jointly intercepted a

common signal. After detecting the presence of an emitter, data is then shared among the sensors to estimate the emitter's position using the TDOA method. In these two sequential stages, our task-embedded compression can be applied as follows: (i) the data stream that is shared during the detection phase is optimally compressed for detection, then (ii) the additional data "layer" needed to optimally estimate TDOA is sent.

For Gaussian noise, the compression distortion measures for the detection task are (3.66)-(3.67) that depend on the SNR of post-compressed data while the distortion measures for TDOA are based on its FI in (3.28)-(3.31) that depend on the quadraticallyfrequency-weighted DFT coefficients. Suppose that two bit rates $R_{D,1}$ and $R_{E,1}$ for the detection task and the estimation task are given, respectively. We can think of these two rates as specifying corresponding resource consumption, resource-for-detection $U_{D,1}$ and resource-for-estimation $U_{E,1}$, respectively, where the resource consumption means some combination of energy and time whose exact form is up to different network design criteria. Suppose further that optimizing the two tasks independently for their given allocations yields a corresponding probability-of-miss $(P_{M,1})$ and estimation error standard deviation $(\sigma_{E,1})^{1}$. We can visualize the scenario in terms of these performance measures as shown in the "conceptual" plots shown in Figure 5.1. The top plot on the left in this figure shows the magnitude-squared DFT of the signal and the noise floor... to maximize SNR for the detection task, the most important data is the DFT coefficients that have magnitude-squared values above the noise floor. Suppose that the region shaded in blue in the top plot is the optimally selected DFT coefficients for detection given the rate

¹ We use probability-of-miss here (rather than the related probability-of-detection) so that the performance measures for detection and estimation each have the characteristic of "bigger is bad, smaller is good".



Figure 5.1: Sequential Tasks with the Initial Trade-Off.

constraint $R_{D,1}$. The bottom plot on the left in the figure shows the quadraticallyweighted magnitude-squared DFT that drives the selection process for the TDOA-driven compression. Suppose that the region shaded in solid blue in the bottom plot is the optimally selected DFT coefficients for estimation given the rate constraint $R_{E,1}$. Note that the part that was previously sent for detection (shown in the bottom plot with stripedblue shading) is already at S_2 so it can be used for TDOA processing. The performance trade-off is shown in the plot on the right of the figure. We take this performance as a "baseline scenario" and see what else may be accomplished with respect to this baseline. On the other hand, it should be noticed that an inherent tradeoff exists between different stages. For example, if it is determined that the resource-for-estimate $U_{E,1}$ is too much to meet other important design requirement and has to be decreased while we want to keep the estimation accuracy about the same and we wish to keep the resource-for-detection $U_{D,1}$ about the same. This can only be done by forcing the detection stage to send data



Figure 5.2: Sequential Tasks with a New Trade-Off.

that would ultimately be helpful in the estimation stage with the consequence of the degradation detection performance (an increase in the Probability-of-Miss). The top plot of Figure 5.2 shows a heuristic way for this case where we remove some of the previously selected DFT coefficients and include higher frequencies in the detection selection set in the detection stage. The new chosen DFT coefficients should take about the same number of bits (actually it would likely take fewer bits because of the smaller amplitudes at these higher frequencies... but we'll ignore that here). The vertical dashed lines are the same as in the first scenario and simply show that we are now sending higher frequency info during detection than we did before. In the bottom plot of Figure 5.2, the solid yellow shows the new data that we need to send to support estimation... the striped-yellow shaded data is what was already there from the detection task.

So the resource-for-estimation is decreased while the location accuracy will be about the same. The detection performance will be worse but the resource-for-detection will be about the same. This performance trade-off is shown on the right side of Figure 5.2, where blue points indicate the initial trade-off scenario and the yellow points indicate the new trade-off scenario. Note that the total resource for these two successive tasks has also been reduced. (4.40) gives us the standard optimization methods for us to make tradeoffs between detection and estimation and we can use it here in the detection stage. This shows that there is a fundamental similarity between the simultaneous task case and the sequential task case, compression algorithms for simultaneous algorithm (4.12), (4.15) and (4.39) can be easily modified to give the following sequential algorithm.

5.1.1 An Algorithm for Sequential Detect-Then-TDOA

Detection stage 1: Maximizing SNR for the Detection Task according to (4.40)

$$Maximize \left[\beta \times \sum_{n=-N/2}^{N/2} \frac{n^2 |X[n]|^2}{\sigma^2 + q^2(b_{n_1})} + (1 - \beta) \times \sum_{n=-N/2}^{N/2} \frac{|X[n]|^2}{\sigma^2 + q^2(b_{n_1})} \right] \text{ subject to } \sum_{n=0}^{N-1} b_{n_1} \le R_D$$
(5.1)

where X[n] is the DFT of the data, $q^2(b_{n_1})$ is the quantization noise power as a function of the number of bits b_{n_1} allocated to that coefficient, R_D is the rate for Stage 1, and the parameter β controls the tradeoff between TDOA estimation and detection. The first summation term in (5.1) measures the importance of a bit for TDOA estimation while the second summation measures the importance of the bit for detection. Thus, setting $\beta = 0$ causes this allocation to be done with no consideration of the Stage 2 task of TDOA estimation; however, increasing β forces more consideration of the subsequent Stage 2 task. Thus, instead of achieving trade-offs by the "discard-and-move" operations described above, here the trade-off is achieved by β .

Stage 2: Maximizing Fisher Information for the TDOA Estimation Task According to (3.28)

In this stage only TDOA is of interest and its accuracy needs to be refined. Thus we will maximize

$$\underset{\{b_{n_2}\}}{Maximize} \left[\sum_{n=-N/2}^{N/2} \left(\frac{n^2 |S[n]|^2}{\sigma^2 + q^2 (b_{n_1} + b_{n_2})} \right) \right] \text{ subject to } \sum_{n=0}^{N-1} b_{n_2} \le R_E$$
(5.2)

where R_E is bit budget for the stage 2, and the $\{b_{n1}\}$ are the bits already allocated in the first stage.

5.1.2 Simulation Scenario

In terms of practical coding, unlike the simultaneous case, quantizers must be changed to dead-zone uniform embedded quantizer in order to support the task-embedded data stream between the sensors. Moreover, for the ease of implementation and testing, we will focus on a trade-off scenario that is an alternative to that shown in Figure 5.2; the case considered is shown in Figure 5.3 where the resourses for different tasks are held fixed and only performance is changed during the trade-off: performance points move along the vertical dotted lines. However, it should be pointed out that once this trade-off is obtained, estimation bits could be discarded to decrease U_E and increase $\sigma_{E,2}$; thus, this new fixed-task-resource scenario can actually be made equivalent to the fixed-estimation-accuracy/reduced-estimation-time scenario described above.



Figure 5.3: Trade-off for fixed task resources .

For the simulation results below we imposed the conditions:

- During the first detection stage, in order to satisfy a resource constraint U_D , the data is compressed with a compression ratio of 8:1.
- During the second estimation stage, in order to satisfy a time constraint U_E ,

Stage 2 bits are added such that the total compression ratio is 5:1.

As illustrated in Figure 5.4, the variation of the parameter β enables the tradeoffs for the detection performance during the first detection stage and the location accuracy during the second estimation stage. Actual simulation results are shown in Figure 5.5, where the actual probabilities in the detection stage were not evaluated due to the fact that such evaluations would take tremendous time. The results in Figure 5.5 show the various SNR- σ_E points that can be achieved for fixed U_D and U_E ; the points in the upper right corner favor detection while those in the lower left corner favor estimation. Note that the values of the post-compression detection-stage SNR were quite low; this is due to the large compression ratio that was imposed. However, detection accuracy is not only dependent on post-compression SNR but also on the size of data; any increase in SNR indicates a significant decrease in error probability for large sample sizes (3.55). In another point of view, it should be noted that this SNR is increased via the crosscorrelation by the time-bandwidth product, which generally is large enough to pull such low SNR values up to an easily detectable level [5][69]. The specific operating values shown in Figure 5.5 are of less importance than the fact that the simulation results verify that the algorithm achieves a curve of trade-off points as expected.



Figure 5.4: Conceptual illustration of the trade-off accomplished via choice of the β parameter. Where $SNR_1 = SNR_2 = 15$ dB.



Figure 5.5: Simulation results illustrating the achieved trade-offs.

5.2 General Sequential Data Compression Algorithm

Based on the detection-then-TDOA example, the general idea of task-embedded compression is very clear. In the following, we extend it into the generalized case. Without loss of generality, we assume there are only two stages and the ultimate goal is to estimate the unknown vector $\mathbf{\theta} = [\theta_1, \dots, \theta_p]$ or unknown $1 \times r$ vector $\mathbf{a} = \mathbf{g}(\mathbf{\theta})$.

First Inference Stage (Detection or Rough Estimation):

In this stage, we either don't know if the target is present or we have no priori geometry to compress data optimally. Therefore, responding to different processes under investigation, optimization of the compression algorithm can be any one of the following three under the constraint that the summation of the bit allocation vector $B_1 = [b_{n1}, \dots, b_{n2}]$ is not greater than R_1 .

1. Maximize the Chernoff distance $\mu^{s}(\mathbf{x}_{1}, B_{1})$ for the detection purpose, i.e.,

$$\max_{B_{1}} \boldsymbol{\mu}^{s}(\mathbf{x}_{1}, B_{1}) = \sum_{n=0}^{N-1} \frac{|S_{n}|^{2}}{\sigma^{2} + q^{2}(b_{n_{1}})}$$
(5.2)

where S_n is the ON transform of \mathbf{x}_1 .

2. Maximize the information utility function ψ of FIM for the rough estimation purpose. ψ could be either determinant or trace² as

$$\max_{B_{1}} \det(\hat{\mathbf{J}}_{1,\theta}(\mathbf{x}_{1}, B_{1})) = \begin{vmatrix} \hat{J}_{11}(\mathbf{x}_{1}, B_{1}) & \cdots & \hat{J}_{1p}(\mathbf{x}_{1}, B_{1}) \\ \vdots & \ddots & \vdots \\ \hat{J}_{p1}(\mathbf{x}_{1}, B_{1}) & \cdots & \hat{J}_{pp}(\mathbf{x}_{1}, B_{1}) \end{vmatrix}$$

$$= \begin{vmatrix} \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{11}}{\sigma^{2} + q^{2}(b_{n1})} & \cdots & \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{1p}}{\sigma^{2} + q^{2}(b_{n1})} \\ \vdots & \ddots & \vdots \\ \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{p1}}{\sigma^{2} + q^{2}(b_{n1})} & \cdots & \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{pp}}{\sigma^{2} + q^{2}(b_{n1})} \end{vmatrix}$$
(5.3)

$$\max_{B_1} \operatorname{tr}(\hat{\mathbf{J}}_{1,\mathbf{\theta}}(\mathbf{x}_1, B_1)) = \sum_{i=1}^p \hat{J}_{ii}(\mathbf{x}_1, B_1) = \sum_{i=1}^p \sum_{n=0}^{N-1} \frac{\Gamma_n^n}{\sigma^2 + q^2(b_{n1})}$$
(5.4)

² We might only need to maximize determinant or trace of a sub-matrix of the Fisher information matrix, corresponding to certain inference task of interest. *Although weighted trace can be used too, due to the lack of information to determine the weight* α , we have to let us stay with the assumption that all the inference tasks are equally important at the inference stage 1

where
$$\Gamma_n^{ij} = \frac{\partial S_n}{\partial \theta_i} \frac{\partial S_n}{\partial \theta_j}$$
.

Maximize the weighted combination of information utility function ψ of FIM and Chernoff distance μ^s(x₁), i.e., βψ(Ĵ_{1,θ}(x₁, B₁)) + (1 - β)μ^s(x₁, B₁) for situation where the first stage is forced to send data that would ultimately be helpful in the second stage.

Second Inference Stage (estimation refinement):

In this stage, central fusion sensor could determine the optimal strategy to do the subsequent compression to refine the estimation given the known the presence of target or priori knowledge for the relative importance of multiple estimate for a specific task from the first stage. The compression can take one of the following two cases under the constraint that the sum of the supplemental bit allocation vector $B_2 = [b_{12}, ..., b_{n2}]$ is not greater than R_2 :

1. Only a single θ is of interest, we need to

$$\max_{B_1} \operatorname{tr}(\hat{\mathbf{J}}_{2,\theta}(\mathbf{x}_1, B_1)) = \sum_{i=1}^p \alpha_i \hat{J}_{ii}(\mathbf{x}_1, B_1) = \sum_{i=1}^p \sum_{n=0}^{N-1} \frac{\alpha_i \Gamma_n^{ii}}{\sigma^2 + q^2(b_{n1})} \quad ; \quad (5.5)$$

2. The whole set parameter $\boldsymbol{\theta} = [\theta_1, \dots, \theta_p]$ are of interest, we need to

$$\max_{B_{2}} \operatorname{Trace} \left(\hat{\mathbf{G}}^{T} \begin{bmatrix} \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{11}}{\sigma^{2} + q^{2}(b_{n2} + b_{n1})} & \cdots & \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{1p}}{\sigma^{2} + q^{2}(b_{n2} + b_{n1})} \\ \vdots & \ddots & \vdots \\ \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{p1}}{\sigma^{2} + q^{2}(b_{n2} + b_{n1})} & \cdots & \sum_{n=0}^{N-1} \frac{\Gamma_{n}^{pp}}{\sigma^{2} + q^{2}(b_{n2} + b_{n1})} \end{bmatrix} \hat{\mathbf{G}} \right), \quad (5.6)$$

where $\hat{\mathbf{G}}$ is the $r \times p$ Jacobian matrix $\hat{\mathbf{G}} = \partial g(\theta) / \partial \theta$ estimated from the first stage.

Two more simulations are performed to exploit this task-embedded compression algorithm. The first simulation is for the sequential multiple parameter estimation, where in the first stage the determinant optimization algorithm is used to get a rough estimate of TDOA/FDOA/DOA according to (5.3). Based on which, one of TDOA and FDOA and DOA estimate is determined as the dominant parameter that needs refining. Then in the second stage, subsequent compression for any one of TDOA or FDOA or DOA is performed according to (5.5). The determinant method is chosen for the first stage because it is more likely to produce the best tradeoff among three different location methods, TDOA, FDOA, and DOA, which provide us with a good initial estimate of the parameter when we have no prior information of the location of the emitter. The compression ratio imposed on the first stage is 8:1, and bits are added in the second stage until the total compression ratio is 5:1. Pre-compression SNR at the two sensors (SNR₁) and SNR_2) are both set as 15 dB. The effectiveness of this method is shown in Table 1. The first three rows of Table. 1 correspond to roughly estimating TDOA/FDOA/DOA in the first stage by maximizing determinant of FIM and refining only one single parameter TDOA, FDOA and DOA in the second stages, respectively. The first three columns are the error standard deviation for the rough estimation of TDOA (ns), FDOA (Hz) and DOA (m °C) in the first stage, respectively and the last three columns are the final estimation error standard deviation for TDOA, FDOA and DOA in the refining second stage, where the shaded and underlined shows the final estimation error standard deviation of the parameter wanted to be refined in the second stage. The estimation accuracies for TDOA/FDOA/DOA without any compression are also shown in Table.1 for the comparison purpose. We can see that at the second stage, the accuracy of the

parameter of interest is very close to the case where no compression is performed. The scheme of second simulation is same to the first one except that the detection is the required task in the first stage and compression for maximizing SNR is performed instead in this stage. The results are shown in Table 2.

	First Stage Compression				Second St	Second Stage Compression		
	TDOA (S1)	FDOA (S1)	DOA (S1)		TDOA (S2)	FDOA (S2)	DOA (S2)	
det(FIM)	1.7655	3.0168	0.10761	TDOA only	0.61107	1.8143	0.068	
det(FIM)	1.8732	3.1543	0.10428	FDOA only	0.82308	1.5719	0.068	
det(FIM)	1.8466	3.0799	0.1058	DOA only	0.825	1.8035	<u>0.059</u>	
					TDOA	FDOA	DOA	
					(w/o)	(w/o)	(w/o)	
					0.51755	1.3621	0.0501	

 Table 5.1, Sequential Example, Rough Multiple Estimate - Then- Refine Estimation

 Table 5.2, Sequential Example, Detection-Then - Estimation

	First Stage Compression				Second Stage Compression		
	TDOA (S1)	FDOA (S1)	DOA (S1)		TDOA (S2)	FDOA (S2)	DOA (S2)
Detect	1.6472	3.1896	0.1005	TDOA only	<u>0.6028</u>	1.8306	0.0652
Detect	1.6388	3.2728	0.1000	FDOA only	0.8132	<u>1.5575</u>	0.0667
Detect	1.6132	3.1146	0.0973	DOA only	0.843	1.7125	<u>0.0632</u>

CHAPTER 6: Conclusion

This dissertation addresses using data compression as the tool to improve the energy efficiency and timeliness of wireless sensor networks for inference tasks (estimation/detection). Firstly, investigations are performed to assert that compression can be deployed as a powerful and effective tool for the optimal trade-offs between rate, energy, and accuracy in the wireless sensor network; Secondly, instead of using the traditional MSE distortion measure that is only weakly related to the quality of estimation and detection tasks, we use distortion measures based on the Fisher information for estimation tasks and the Chernoff-distance for detection; The new developed distortion measures are able to better evaluate the underlying information embedded in the collected data for estimation/detection. As demonstrated by exhaustive examples of practical interest, the compression algorithms based on our proposed distortion significantly outperform those based on MSE. Finally, although the sequential and simultaneous inference tasks are typical in wireless sensor network applications, no other researchers have ever addressed compression for these two issues for wireless sensor networks, and even for networks of several macrosensors. We develop the theories and compression algorithms for the general simultaneous estimation problem and joint estimation and detection problem as well as the theories and algorithms for "task-embedded compression" to support sequential multiple inference tasks. In all, this dissertation sets up the whole compression structure for inference tasks in wireless sensor networks and the results can be used as a key component of and incorporated into the task-driven network management and sensor management.

Despite these accomplishments, there are still some details insides our framework for which some further work are needed:

- 1. In Chapter 3, although the general FI and Chernoff distance of quantized noisy measurements are formulated and are applicable to all types of noise and quantizers, in this dissertation, for the purpose of simplicity, we limit our focus mainly on the case when the noise is Gaussian noise and the quantizer is uniform quantizer due to their widely usage in literature and propose the corresponding approximation model, which is used as a replacement for the true model in Chapter 4 and 5. Further exploitation of other type of noise will definitely further strengthen our compression structure.
- 2. For compressions for estimation, in general the FIM can depend on the parameter vector values, although in this dissertation, we limit ourselves to some practical estimation problems the FIM does not depend on the parameter values or at least the FIM can be separated into the part containing only the unknown-parameter-centric part and the data-centric part containing to which our "maximize Fisher information" approach can still be applied. However, there may exist the case where the Fisher information depends on the parameters which can not be pulled out of the data as in (3.43). We have proposed the correct approach for this problem but have not yet fully explored its application. Thus, some further research is needed to determine if there are estimation problems for which the "maximize Fisher information" approach

can not be directly applied and then demonstrate how to attack them using the minimax approach.

- 3. Regarding to the geometry adaptive algorithm for joint TDOA/FDOA estimation, although we demonstrate that we can use parallel auxiliary STFT to access the FIM elements and help the algorithm to evaluate the right relative importance of wavelet packet coefficients to joint TDOA/FDOA estimation. To improve the computational efficiency, Further works is still needed to develop filter banks that are well-suited to both compression and FIM evaluation
- 4. There are still some small details surrounding the detection problem, such as instead of real probability of detection, we use the Chernoff distance as the indication of detection. Although this is true from the theory, further work is still needed to verify it.
- 5. Although a lot of effort is paid to choose and improve the computation efficiency of compression algorithms to achieve maximal saving of time and energy, further examination of the computational and implementation aspects of the algorithm is still needed, and hardware implementation is desired in the future.
APPENDICES

APPENDIX A

A.1 Fisher Information Separation

Theorem A.1: Assume that we have a set of data **x** whose PDF $p(\mathbf{x}; \theta)$ depends on a parameter θ to be estimated. If the data can be split into two independent sets $\mathbf{x}_1 \sim p_1(\mathbf{x}_1, \theta)$ and $\mathbf{x}_2 \sim p_2(\mathbf{x}_2, \theta)$; then (due to independence) $p(\mathbf{x}, \theta) = p_1(\mathbf{x}_1, \theta)p_2(\mathbf{x}_2, \theta)$. Let the individual PDFs satisfy the "regularity" condition (2.3)

$$E\left[\frac{\partial \ln p_k(\mathbf{x}_k;\theta)}{\partial \theta}\right] = 0, \quad k = 1,2$$
 (A.1.1)

Let *J* be the joint Fisher information of the complete data set **x** and let J_k be the Fisher information of the k^{th} subset \mathbf{x}_k ; then $J = J_1 + J_2$. Or more equivalently,

$$E\left[\frac{\partial \ln p(\mathbf{x}_{1}, \mathbf{x}_{2}; \theta)}{\partial \theta} \frac{\partial \ln p(\mathbf{x}_{1}, \mathbf{x}_{2}; \theta)}{\partial \theta}^{H}\right]$$
$$= E\left[\frac{\partial \ln p_{1}(\mathbf{x}_{1}; \theta)}{\partial \theta} \frac{\partial \ln p_{1}(\mathbf{x}_{1}; \theta)}{\partial \theta}^{H}\right] + E\left[\frac{\partial \ln p_{2}(\mathbf{x}_{2}; \theta)}{\partial \theta} \frac{\partial \ln p_{2}(\mathbf{x}_{2}; \theta)}{\partial \theta}^{H}\right]$$
(A.1.2)

<u>Proof</u>: Because \mathbf{x}_1 and \mathbf{x}_2 are independent, we have $p(\mathbf{x}_1, \mathbf{x}_2; \theta) = p_1(\mathbf{x}_1; \theta) p_2(\mathbf{x}_2; \theta)$ and then

$$E\left[\frac{\partial \ln p(\mathbf{x}_{1}, \mathbf{x}_{2}; \theta)}{\partial \theta} \frac{\partial \ln p(\mathbf{x}_{1}, \mathbf{x}_{2}; \theta)}{\partial \theta}^{H}\right] = E\left[\frac{\partial \ln p_{1}(\mathbf{x}_{1}; \theta)}{\partial \theta} \frac{\partial \ln p_{1}(\mathbf{x}_{1}; \theta)}{\partial \theta}^{H}\right] + E\left[\frac{\partial \ln p_{2}(\mathbf{x}_{2}; \theta)}{\partial \theta} \frac{\partial \ln p_{2}(\mathbf{x}_{2}; \theta)}{\partial \theta}^{H}\right] + 2E\left[\frac{\partial \ln p_{1}(\mathbf{x}_{1}; \theta)}{\partial \theta} \frac{\partial \ln p_{2}(\mathbf{x}_{2}; \theta)}{\partial \theta}^{H}\right]$$
(A.1.3)

Due to independence the third term on the right hand side of (A.1.3) can be written as

$$E\left[\frac{\partial \ln p_1(\widetilde{\mathbf{x}}_1;\theta)}{\partial \theta}\frac{\partial \ln p_2(\widetilde{\mathbf{x}}_2;\theta)^H}{\partial \theta}\right] = E\left[\frac{\partial \ln p_1(\widetilde{\mathbf{x}}_1;\theta)}{\partial \theta}\right] \times E\left[\frac{\partial \ln p_2(\widetilde{\mathbf{x}}_2;\theta)^H}{\partial \theta}\right].$$
 (A.1.4)

and each term on the right hand side of (A.1.4) is zero due to the assumed regularity condition. Thus, (A.1.3) reduces to (A.1.2) and the claimed result is proved.

A.2 Chernoff Distance Separation

Theorem A2: Consider a problem consisting of deciding between two hypothesis H_0 and H_1 based on a set of observed random vectors **x**. Hypotheses H_0 and H_1 are (no target) and (target present) hypothesis, respectively, and are represented by

$$H_0: \mathbf{x} \sim p_0(\mathbf{x})$$

$$H_1: \mathbf{x} \sim p_1(\mathbf{x})$$
(A.1.5)

If the data can be split into two independent sets \mathbf{x}_1 and \mathbf{x}_2 , each of which follows

$$\begin{aligned} H_0 &: \mathbf{x}_i \sim p_{0,i}(\mathbf{x}_i), \quad i = 1,2 \\ H_1 &: \mathbf{x}_i \sim p_{1,i}(\mathbf{x}_i), \quad i = 1,2 \end{aligned}$$
 (A.1.6)

Let $\mu_s(p_0(\mathbf{x}), p_1(\mathbf{x}))$ be the joint Chernoff distance of the complete data set \mathbf{x} and let $\mu_s(p_{0,k}(\mathbf{x}_k), p_{1,k}(\mathbf{x}_k))$ be the Chernoff distance of the k^{th} subset \mathbf{x}_k ; then $\mu_s(p_0(\mathbf{x}), p_1(\mathbf{x})) =$

 $\mu_s(p_{0,1}(\mathbf{x}_1), p_{1,1}(\mathbf{x}_1)) + \mu_s(p_{0,2}(\mathbf{x}_2), p_{1,2}(\mathbf{x}_2))$. Or more equivalently,

$$-\ln \int p_0(\mathbf{x})^{1-s} p_1(\mathbf{x})^s d\mathbf{x} = -\ln \int p_{0,1}(\mathbf{x}_1)^{1-s} p_{1,1}(\mathbf{x}_1)^s d\mathbf{x}_1 - \\ \ln \int p_{0,2}(\mathbf{x}_2)^{1-s} p_{1,2}(\mathbf{x}_2)^s d\mathbf{x}_2$$
(A.1.7)

<u>Proof</u>: Because \mathbf{x}_1 and \mathbf{x}_2 are independent, we have $p_0(\mathbf{x}) = p_{0,1}(\mathbf{x}_1)p_{0,2}(\mathbf{x}_2)$ and $p_1(\mathbf{x}) = p_{1,1}(\mathbf{x}_1)p_{1,2}(\mathbf{x}_2)$ then

$$-\ln \int p_0(\mathbf{x})^{1-s} p_1(\mathbf{x})^s d\mathbf{x} = -\ln \left(\int \int p_{0,1}(\mathbf{x}_1)^{1-s} p_{1,1}(\mathbf{x}_1)^s p_{0,2}(\mathbf{x}_2)^{1-s} p_{1,2}(\mathbf{x}_2)^s d\mathbf{x}_1 d\mathbf{x}_2 \right)$$

= $-\ln \int p_{0,1}(\mathbf{x}_1)^{1-s} p_{1,1}(\mathbf{x}_1)^s d\mathbf{x}_1 \int p_{0,2}(\mathbf{x}_2)^{1-s} p_{1,2}(\mathbf{x}_2)^s d\mathbf{x}_2$
= $-\ln \int p_{0,1}(\mathbf{x}_1)^{1-s} p_{1,1}(\mathbf{x}_1)^s d\mathbf{x}_1 - \ln \int p_{0,2}(\mathbf{x}_2)^{1-s} p_{1,2}(\mathbf{x}_2)^s d\mathbf{x}_2$

(A.1.8)

Thus the claimed result is proved.

APPENDIX B

B. 1: Fisher information of scalar quantized data

Let $r_n = Q_{bn}(\xi_n(\theta) + \omega_n)$, and p_{ω_n} denote the probability density function (PDF) of ω_n . The thresholds of the scalar quantization of χ_n $\{t_i\}$ is illustrated in Figure 2.3, where $M = 2^{b_n}$. The probability of r_n falling into the interval $(-\infty, t_1]$ can be thereby calculated as

$$f_{r_n}^{q_0}(b_n;\theta) = \operatorname{Prob}(r_n = q_0 \mid \chi_n)$$

= $\operatorname{Prob}(\xi_n(\theta) + \omega_n < t_1) = \operatorname{Prob}(\omega_n < t_1 - \xi_n(\theta)),$ (B.1)
= $\int_{-\infty}^{t_1 - \xi_n(\theta)} p_{\omega_n}(t) dt = \eta_{\omega_n}(t_1 - \xi_n(\theta))$

where η_{ω_n} is the cumulative distribution function for PDF of ω_n and is defined as

$$\eta_{\omega_n}(x) = \int_{-\infty}^{x} p_{\omega_n}(t) dt \,. \tag{B.2}$$

Similarly, the probability of r_n falling into the interval $[t_{M-1}, \infty)$ are calculated as

$$f_{r_n}^{q_{M-1}}(b_n;\theta) = \operatorname{Prob}(\xi_n(\theta) + \omega_n > t_{M-1})$$

= $\operatorname{Prob}(\omega_n > t_{M-1} - \xi_n(\theta))$ (B.3)
= $\int_{t_{M-1}-\xi_n(\theta)}^{\infty} p_{\omega_n}(t)dt = 1 - \eta_{\omega_n}(t_{M-1} - \xi_n(\theta))$

Except for these two end points, all other probability in the interval $I_j = [t_j, t_{j+1}]$ can be calculated as

$$f_{r[k]}^{q_{j}}(b_{n};\theta) = \operatorname{Prob}(r_{n} = q_{j} \mid \chi_{n}) = \operatorname{Prob}(t_{j} < \xi_{n}(\theta) + \omega_{n} < t_{j+1})$$

$$= \operatorname{Prob}(t_{j} - \xi_{n}(\theta) < \omega_{n} < t_{j+1} - \xi_{n}(\theta))$$

$$= \int_{t_{j} - \xi_{n}(\theta)}^{t_{j+1} - \xi_{n}(\theta)} p_{\omega_{n}}(t) dt$$

$$= \eta_{\omega_{n}}(t_{j+1} - \xi_{n}(\theta)) - \eta_{\omega_{n}}(t_{j+1} - \xi_{n}(\theta))$$
(B.4)

Applying Leibnitz's rule, (B.1) yields

$$\frac{\partial \left(\eta_{\omega_n}(t_1 - \xi_n(\theta))\right)}{\partial \theta} = -p_{\omega_n}(t_1 - \xi_n(\theta))\frac{\partial \left(\xi_n(\theta)\right)}{\partial \theta}.$$
(B.5)

(B.5) makes it possible to get the closed form of the Fisher information of any scalar quantized data. With it, the derivative of (B.3) and (B.4) with respect to θ are

$$\frac{\partial \left(f_{r_n}^{q_{M-1}}(b_n;\theta)\right)}{\partial \theta} = p_{\omega_n}(t_{M-1} - \xi_n(\theta))\frac{\partial \left(\xi_n(\theta)\right)}{\partial \theta},\tag{B.6}$$

$$\frac{\partial \left(f_{r_n}^{q_j}(b_n;\theta)\right)}{\partial \theta} = -\left(p_{\omega_n}(t_{j+1} - \xi_n(\theta)) - p_{\omega_n}(t_j - \xi_n(\theta))\right) \frac{\partial \left(\xi_n(\theta)\right)}{\partial \theta}.$$
 (B.7)

Plug these results into (3.12), the Fisher information of b_n bit quantized coefficient χ_n

is

$$\begin{split} \widetilde{J}_{r_{n}}(b_{n},\theta) &= \left(\frac{\partial(\xi_{n}(\theta))}{\partial\theta}\right)^{2} \left[\frac{p_{\omega_{n}}^{2}(t_{1}-\xi_{n}(\theta))}{\eta_{\omega_{n}}(t_{1}-\xi_{n}(\theta))} + \sum_{q=1}^{M-2} \frac{\left(p_{\omega_{n}}^{2}(t_{q+1}-\xi_{n}(\theta))-p_{\omega_{n}}^{2}(t_{q}-\xi_{n}(\theta))\right)^{2}}{\eta_{\omega_{n}}(t_{q+1}-\xi_{n}(\theta))-\eta_{\omega_{n}}(t_{q}-\xi_{n}(\theta))} \\ &+ \frac{p_{\omega_{n}}^{2}(t_{M-1}-\xi_{n}(\theta))}{1-\eta_{\omega_{n}}(t_{M-1}-\xi_{n}(\theta))}\right] \end{split}$$
(B.8)

B. 2, Approximation model

Consider

$$\mathbf{v} = \mathbf{\omega}_1 + \mathbf{\varepsilon} \tag{B.9}$$

where **v** is the $N \times 1$ vector of the summation of receiver noise and quantization noise, and ε is the $N \times 1$ vector of quantization noise ε_n which is modeled as white and uniform in the interval $[-\Delta/2, \Delta/2]$, and ω_1 is the $N \times 1$ vector of the white Gaussian noise coefficients ω_n with zero mean and variance σ^2 . Now, our question becomes how to analyze the PDF of v_n and approximate it properly. As we know, the probability density function of the sum of independently distributed random variables is given by the convolution of their respective probability density functions, i.e.

$$p_{\mathbf{v}}(x) = p_{\omega_{\mathbf{i}}}(x) * p_{\varepsilon}(x). \tag{B.10}$$

The characteristic function is very useful for determining the probability density function of a sum of independent random variables. The characteristic function of a random variable X is defined as

$$C(w) = \int_{-\infty}^{\infty} p(x)e^{jwx}dx .$$
 (B.11)

That is, C(w) is the complex conjugate of Fourier transform of the probability density function p(x).

The result (B.10) can be related to the characteristic function. It is well known in Fourier transform theory that the transform of a convolution is the product of the respective Fourier transforms. Denote the characteristic functions of $p_v(x)$, $p_{\omega_1}(x)$, and $p_{\varepsilon}(x)$ by $C_v(w)$, $C_{\omega_1}(w)$, and $C_{\varepsilon}(w)$ respectively. Then it follows that

$$C_{\mathbf{v}}(w) = C_{\boldsymbol{\omega}_{1}}(w)C_{\boldsymbol{\varepsilon}}(w). \tag{B.12}$$

The probability density function of v is then determined by the inverse Fourier transform of $C_v(w)$.

The characteristic function for a normal distribution with mean u and variance σ^2 is

$$C(w) = \exp(juw - \sigma^2 w^2/2)$$
 (B.13)

The characteristic function for a uniform distribution on the interval [a,b] is

$$C(w) = \frac{2}{(b-a)w} \sin(\frac{1}{2}(b-a)w) \exp(j(a+b)w/2)$$
 (B.14)

Hence, we can calculate $C_{v}(w)$ by

$$C_{\nu}(w) = C_{\varepsilon}(w)C_{\varepsilon}(w) = \frac{\sin(\Delta w/2)}{\Delta w/2}\exp(-\sigma^2 w^2/2).$$
(B.15)

The inverse Fourier transform of (B.15) is very complicate, but we can approximate $C_{v}(w)$ as a Gaussian distribution

$$\widetilde{C}_{v}(w) = \exp(-\sigma_{v}^{2}w^{2}/2), \ \sigma_{v}^{2} = \operatorname{var}(X+Y) = \sigma^{2} + \Delta^{2}/12$$
 (B.16)

under the condition that the null-to-null bandwidth of the characteristic function $C_q(w) 4\pi/\Delta$ is larger than $6/\sigma$, within which more than 99.7% of the area of the characteristic function $C_n(w)$ is included.

Here we give some numerical results that show the effectiveness of the Gaussian approximation used for the FI with multi-bit quantization. We assume that we have a set of noise-free coefficients that lie in the range $\pm A$. A noisy version of the coefficients having additive Gaussian noise of variance σ^2 is quantized to *b* bits using a mid-step uniform quantizer with quantization cell size given by

$$\Delta = \frac{\sqrt{A^2 + \sigma^2}}{2^{B-1} - 1/2} \,. \tag{B.17}$$

By using our approximation condition, we can get how many bits are at least needed in order for the approximation (B.16) to be valid.

$$4\pi / \Delta \ge 6 / \sigma$$
$$\Rightarrow \Delta \le 2\pi\sigma / 3$$

$$\Rightarrow \frac{\sqrt{A^2 + \sigma^2}}{2^{B-1} - 1/2} \le \frac{2\pi\sigma}{3}$$

$$\Rightarrow B \ge \log_2 \left(\frac{3\sqrt{PSNR + 1}}{2\pi} + \frac{1}{2}\right) + 1$$
 (B.18)

Under the inequality (B.18), we need $B \ge 2.7$ bits at 15 dB and $B \ge 2$ quantizer at PSNR=10 dB (Uniform quantizer cannot quantize the signal with fractional bits, therefore 3 bits is needed at PSNR=15 dB and 2 bits is needed at PSNR=10dB). If we loosen the approximation condition of approximation (B.11) a little bit, i.e., instead of constraining the mainlobe of $C_q(w)$ must contain 99.7% energy of $C_n(w)$, we permit some energy leakage of $C_n(w)$ and let the mainlobe of $C_q(w) 4\pi/\Delta$ include 95 energy of $C_n(w)$ (that is, $4\pi/\Delta$ is larger than $4/\sigma$). In the later case, only a 2-bit quantizer is needed at 15dB. The results provide motivation that at least for low PSNR the approximation seems to be valid even down to the lowest number of bits for which it is applied. For higher PSNR the approximations at 2 bits will be poorer;

Next, we plot the characteristic function to support the approximation model (B.16). As illustrated in Figure B.1, the approximation is perfect when the quantizer is 3 bits and PSNR is 15 dB. For 2-bit quantizer at PSNR=15dB, although there exists a little bit of difference between the approximation and the true distribution, but it is still a pretty good approximation. If the PSNR decreases to 10 dB, the condition for perfect performance of the approximation model only needs 2-bit quantizers.

The validity of this approximation can be further tested analytically by comparing the Fisher information under the approximation with its true expression in (B.8). For b_n bit uniform quantization, the thresholds are $\{-m\Delta_n, -(m+1)\Delta_n, ..., m\Delta_n\}$, where



Figure B.1 (C)

Figure B.1: Comparison of characteristic function between true model and approximation model (A)3-bit quantizer at PSNR=15dB; (B) 2-bit quantizer at PSNR=15dB; (C) 2-bit quantizer at PSNR=10 dB.

m = M/2 and Δ_n is the step size of quantizing $\chi_n(\theta)$. The closed form of Fisher information is:

$$J_{r_{n}}^{b_{n}}(\theta) = \frac{1}{2\pi\sigma^{2}} \left(\frac{\partial(\xi_{n}(\theta))}{\partial\theta}\right)^{2} \left[\frac{\exp\left(-\left(\frac{-M\Delta_{n}-\xi_{n}(\theta)}{\sigma}\right)^{2}\right)}{\Phi\left(\frac{-M\Delta-\xi_{n}(\theta)}{\sigma}\right)} + \sum_{q=1}^{M-2} \left(\frac{\exp\left(-\frac{1}{2}\left(\frac{(j+1)\Delta_{n}-\xi_{n}(\theta)}{\sigma}\right)^{2}\right) - \exp\left(-\frac{1}{2}\left(\frac{j\Delta_{n}-\xi_{n}(\theta)}{\sigma}\right)^{2}\right)\right)^{2}}{\Phi\left(\frac{(j+1)\Delta_{n}-\xi_{n}(\theta)}{\sigma}\right) - \Phi\left(\frac{j\Delta_{n}-\xi_{n}(\theta)}{\sigma}\right)} + \frac{\exp\left(-\left(\frac{M\Delta_{n}-\xi_{n}(\theta)}{\sigma}\right)^{2}\right)}{1 - \Phi\left(\frac{M\Delta_{n}-\xi_{n}(\theta)}{\sigma}\right)}\right], \quad (B.19)$$
$$= \left(\frac{\partial(\xi_{n}(\theta))}{\partial\theta}\right)^{2} I_{t}(\xi_{n}(\theta), \Delta_{n}, \sigma)$$

where $\Phi(x) = \eta_{\omega_n}(x/\sigma) = \int_{-\infty}^{x} \frac{1}{\sqrt{2\pi}} \exp(-\frac{1}{2}t^2) dt$. The expression of the Fisher

information under the approximated model is

$$\widetilde{J}_{r_n}^{b_n}(\theta) = \left[\frac{\partial \xi_n(\theta)}{\partial \theta}\right]^2 \frac{1}{\sigma^2 + \frac{\Delta_n^2}{12}} = \left[\frac{\partial \xi_n(\theta)}{\partial \theta}\right]^2 \mathbf{I}_a(\Delta_n, \sigma).$$
(B.20)

and the Fisher information without any quantization is given by

$$J(\theta) = \left[\frac{\partial \xi_n(\theta)}{\partial \theta}\right]^2 \frac{1}{\sigma^2} = \left[\frac{\partial \xi_n(\theta)}{\partial \theta}\right]^2 \mathbf{I}(\sigma).$$
(B.21)

Notice that (B.19) and (B.20) and (B.21) all have the product factor $[\partial \xi_n(\theta)/\partial \theta]^2$ in common. In order to see the difference between the true model (B.19) and the approximate model (B.20), and the degradation effect on Fisher information by quantization, we only need to compare the factor $I_t(\xi_n(\theta), \Delta_n, \sigma)$ in (B.19), $I_a(\Delta_n, \sigma)$ in (B.20) and $I(\sigma)$ in (B.21) under different values of $\xi_n(\theta)$ and Δ_n^2 and σ . In the following numerical results, we assume that $\xi_n(\theta)$ varies in the range from -1 to 1, σ is determined by different PSNR and $\Delta_n = (2+2\sigma)/2^{b_n}$, the values of $I_t(\xi_n(\theta), \Delta, \sigma)$ and $I_a(\Delta, \sigma)$ and $I(\sigma)$ are computed and shown in Figure B.2. When PSNR=10 dB, the approximation model fits well with the true model when the number of bits is greater than 1 bit. When

PSNR is increased to 15 dB, the number of bits required for the approximation model to fit well with the true model is increased to 2 bits. Figure B.2 not only further verify the fact that quantization decreases the Fisher information since $I_t(\xi_n(\theta), \Delta, \sigma)$ is always less than $I(\sigma)$, but also illustrates there is some limitation for us to use the approximation to replace the true model. We can only use the approximate model when the bits of quantizer is greater given threshold which is dependent on the variance of noise σ . This thresholds for different σ is tabulated in Table B.1.

PSNR Bits	5	10	15	20	25	30
1	x	x	x	x	x	x
2	\checkmark	\checkmark	x	x	x	x
3	\checkmark	\checkmark	\checkmark	x	x	x
4	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	x
5	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark

Table B.1, Fitness of the approximation model with the true model, $\sqrt{}$ means fit while X means not

B.3. Chernoff distance of scalar quantized data.

Proceeding as before, we can get the close form of Chernoff distances of scalar quantized data. At the two end points, the probability $f_{r_n}^{\hat{x}_0}[b_n; H_0]$ and $f_{r_n}^{\hat{x}_{M-1}}[M; H_0]$ can be calculated by:

$$f_{r_n}^{\hat{x}_0}(b_n; H_0) = \operatorname{Prob}(\omega_n < t_1)$$

=
$$\int_{-\infty}^{t_1} p_{\omega}(t) dt = \eta(t_1)^{\prime}$$
(B.22)

and



(a) PSNR = 10 dB



(b) PSNR = 15 dB

Figure B.2: Comparison of value $I_t(\xi_n(\theta), \Delta_n, \sigma)$ and $I_a(\Delta, \sigma)$ under different values of $\xi_n(\theta)$ and Δ_n^2 when (a) PSNR =10 dB, (b) PSNR=15 dB.

$$f_{r_n}^{\hat{x}_{M-1}}(b_n; H_0) = \operatorname{Prob}(\omega_n > t_{M-1})$$

= $\int_{t_{M-1}}^{\infty} p_{\omega}(t) dt = 1 - \eta(t_{M-1})$, (B.23)

respectively. On the other hand, $f_{r_n}^{\hat{x}_0}[b_n; H_1]$ and $f_{r_n}^{\hat{x}_{M-1}}[b_n; H_1]$ are same to (B.1) and (B.3) respectively. The PMF $f_{r_n}^{\hat{x}_q}[b_n; H_0]$ within each quantizer bin is

$$f_{r_n}^{\hat{x}_j}(b_n; H_0) = \operatorname{Prob}(t_j < \omega_n < t_{j+1}) = \int_{t_j}^{t_{j+1}} p_{\omega}(t) dt$$

= $\eta(t_{j+1}) - \eta(t_j)$ (B.24)

whereas $f_{r_n}^{\hat{x}_j}[b_n;H_1]$ is equal to (B.4).

Plugging (B.1), (B.3), (B.4), (B.22), (B.23) and (B.24) into (3.xx), we can write the closed form of Chernoff distance given the quantized sample r_n is:

$$\boldsymbol{\mu}_{r_{n}}^{M}(s) = \ln \left[\eta(t_{1})^{l-s} \eta(t_{1} - \xi_{n})^{s} + \sum_{j=1}^{M-2} (\eta(t_{j+1}) - \eta(t_{j}))^{l-s} \times (\eta(t_{j+1} - \xi_{n}) - \eta(t_{j} - \xi_{n}))^{s} + (1 - \eta(t_{M-1}))^{l-s} (1 - \eta(t_{M-1} - \xi_{n}))^{s} \right]$$
(B.25)

APPENDIX C

Modification of standard image compression methods for pattern recognition

There is also a need for compressing images in timely and energy-saving fashion for distributed image processing applications in wireless sensor networks. One of the most important ones is for pattern recognition (PR) where the images must be transferred from one node to the other before the target recognition task (decision of the presence of target and estimation of its positions) can be completed. Due to the reason that the image usually contains larger amount of data than one-dimensional data, very heavy lossy image compression must be employed in order to accomplish the image transfer within the energy and time requirements of the wireless sensor network. A common approach is to use standard image compression such as JPEG or the wavelet-based algorithm (used in JPEG 2000 [10],[11]) to compress the image. The compressed image is then transferred, decompressed, and processed by the specified algorithms for recognition. Walls and Mahalanobis [55] have studied the effects of the standard wavelet-based compression on the performance of an ATR system that uses the maximum average correlation height (MACH) filter and the distance classifier correlation filter (DCCF). It was shown that these two correlation filters can recognize patterns in IR and SAR imagery at high compression rates. Sims [56] discussed some data compression issues in automatic target recognition and introduced a simple spatial signature metric for distortion quantification of the decompressed data. As suggested in Chapter 3, although it is convenient to apply standard image compression techniques directly in target recognition applications, they may not be able to achieve the best performance. This is because the standard image compression algorithms aims at retaining image fidelity in terms of perceptual quality and mean squared error (MSE) rather than preserving spectrally significant information for PR. In this appendix, we apply the same idea that we proposed in Chapter 3 and develop a new distortion measure to replace the perceptual quality and MSE for compressing images for PR, and show the new algorithm can help achieve better compression performance at the same recognition accuracy as standard compression algorithms, or enhance the recognition performance at the same compression performance as standard compression algorithms.

Different from the optimization formulation that we proposed in the body of the dissertation, here we simply modify the standard compression algorithms so they can be used for pattern recognition to achieve both higher compression ratio and enhanced recognition performance. The reasons for this is two-fold: (1) provide a syntax-compatible interface, where only minimum changes are needed at the encoder and nothing is changed at the decoder, which can blindly decode the bit-stream by assuming that it is a standard compressed image. This will give a lot of leverage to the designer of a sensor network because only the transmitting sensor needs modifying while the rest is left untouched; (2) benefit from the existing advantages that standard image compression algorithms provide (such as simple hardware implementation, progressive transmission, optimal truncation, etc). For example, the option of progressive transmission is important for the sequential inference tasks (see Chapter 5). For example, in the beginning, only bits carrying relevant information for pattern recognition is transmitted, if requested, in the following, more bits carrying reconstruction information will be transmitted. The

proposed algorithms can also incorporate the operations of denoising, edge enhancement, and compression in one encoding process, thus saving processing time.

C.1 Standard Image Compression Algorithms

JPEG and wavelet-based compression are the two most widely used lossy image compression algorithms. We briefly review their principles and point out that they are not optimal compression algorithms for PR.

The JPEG compression is based on the discrete cosine transform (DCT), which has high energy compactness capability and a fast algorithm. An image to be compressed is first divided into 8×8 pixel blocks, which are then transformed using a DCT to the frequency domain. The DCT transform gets low frequencies into the upper left corner of each block, and high frequencies into the lower right corner. Many image blocks have significant coefficients only at low frequencies and thus in the upper left of each block. JPEG is based on the fact that our visual perception is typically less sensitive to highfrequency than to low frequency variations. To minimize the visual degradation of the

coded images, JPEG performs a quantization to keep low-frequency coefficients and discard high-frequency coefficients by dividing the DCT coefficients by the components of a quantization matrix as shown in the following:

$$\hat{T}_{ij} = round\left(\frac{T_{ij}}{q \cdot Q_{ij}}\right),$$
 (C.1)

where T_{ij} and \hat{T}_{ij} are the DCT and corresponding quantized coefficients at location (*i*,*j*), respectively; Q_{ij} are the components of the quantization matrix

$$Q = \begin{bmatrix} 16 & 11 & 10 & 16 & 24 & 40 & 51 & 61 \\ 12 & 12 & 14 & 19 & 26 & 58 & 60 & 55 \\ 14 & 13 & 16 & 24 & 40 & 57 & 69 & 56 \\ 14 & 17 & 22 & 29 & 51 & 87 & 80 & 62 \\ 18 & 22 & 37 & 56 & 68 & 109 & 103 & 77 \\ 24 & 35 & 55 & 64 & 81 & 104 & 113 & 92 \\ 49 & 64 & 78 & 87 & 103 & 121 & 120 & 101 \\ 72 & 92 & 95 & 98 & 112 & 100 & 103 & 99 \end{bmatrix};$$

and q is the quality factor that controls the compression ratio. As can be seen from the quantization matrix, the quantization steps are small in the upper left (low frequencies), and large in the lower right (high frequencies). The result of the division is that many high-frequency coefficients become zero, and the low-frequency coefficients undergo only minor adjustment. After quantization, each quantized block is converted into a 64element vector using the zigzag order and coded by a run-length encoding method. Finally, a lossless entropy coding is performed to get further compression. For decompression of a JPEG image, the reverse process is performed to reconstruct the image from the compressed stream. Wavelet-based image compression has been demonstrated to provide better compression performance than the JPEG at high levels of compression, and was adopted in the new image compression standard JPEG 2000 [10],[11]. The famous embedded zero-tree wavelet (EZW) [10] algorithm first performs the wavelet transformation on the image to be compressed. The wavelet transformation decomposes the input image into a set of subbands of various resolutions. Figure C.1 shows a three-level wavelet decomposition of an image, where L and H denote the lowpass and high-pass approximations, respectively. The coarsest subband is a low-pass approximation of the original image, and the other subbands are finer-scale refinements.

LL	HL	ш	
LH	нн		ш
L	H	HH	
	Lł	4	НН

Figure C.1 Three-level wavelet decomposition

The EZW then scans wavelet coefficients subband by subband, and calculates the zero trees by marking the coefficients appropriately according to their significance with respect to a given threshold. The zero tree is an efficient method to represent the wavelet structure and to quantize wavelet coefficients by discarding the coefficients of the same orientation and same spatial location at finer scales. Arithmetic coding is finally used to entropy code the quantized coefficients. Due to the multi-resolution nature of wavelet decomposition, higher compression ratio can be achieved and more high frequency components can be saved. Despite its differences from JPEG, EZW nonetheless is based on the same principle of discarding the high-frequency coefficients.

From the above review, we can see that the standard compression algorithms were developed for visualization, so they always preserve low frequency and discard highfrequency coefficients. When applied to compress images for PR, they are not optimal in terms of the achievable compression ratio and recognition performance, because for PR high frequencies play a much more important role in discrimination than low frequencies. Edge enhancement (removing low frequencies) has been proven to be very effective in improving a PR system's discrimination capability. Therefore, if we eliminate lowfrequency components during compression, it will have several benefits. First, it can enhance the discrimination performance important for pattern recognition. Second, it makes it possible to integrate edge enhancement and denoising into the same compression process and thus reduce computational complexity. Finally, it can increase the compression ratio significantly. This is because the magnitudes of low-frequency coefficients tend to be much larger than that of high-frequency, and we only need to encode middle-to-high frequency coefficients of relatively small magnitude.

C.2 Distortion Criterion for Compressing Images in PR

Let r(x, y) denote the target image and s(x, y) be the observed input scene that contains the reference image r(x, y) and additive input noise n(x, y). When the target is located at unknown coordinates (x_0, y_0) , the input scene can be written as

$$s(x, y) = r(x - x_0, y - y_0) + n(x, y), \qquad (C.2)$$

The noise n(x, y) is assumed to be zero-mean additive Gaussian noise. Let h(x, y) be the recognition FIR filter being used. The resulting output c(x, y) is given by

$$c(x,y) = \iint s(m,n)h(m+x,n+y)dmdn.$$
(C.3)

In particular, h(x, y) is the MSF when h(x, y) = r(x, y) and is optimal in the sense of producing the maximum output SNR:

$$SNR = \frac{\iint |R(u,v)|^2 du dv}{N_0}$$

where N_0 is the noise spectral density, R(u,v) is the Fourier transform of r(x,y), and u and v are frequencies respectively along the x axis and y axis. If h(x,y) = r(x,y), (C.3) gives the correlation-based pattern recognition (CRP). CPR has been extensively

investigated due to its several advantages over other pattern recognition methods, which include fast optical implementation, shift invariance, relative senor independence, and no need of target segmentation. Therefore, as the benchmark, in the simulations, CPR is used as the optimal pattern recognition method. Usually, correlators must be able not only to identify the objects in an input image but also to locate the desired target in the input scene. Ideally, c(x, y) should have a large peak located at $(x, y) = (x_0, y_0)$ and small values elsewhere. However, the noise in the input scene can obscure the peak and/or shift it from its nominal position. According to the theory of parameter estimation, the shifted coordinates, denoted by (\hat{x}_0, \hat{y}_0) , can by estimated using maximum a posteriori (MAP) estimate,

$$(\hat{x}_0, \hat{y}_0) = \arg\max_{(x_0, y_0)} \left\{ \iint s(x, y) r(x - x_0, y - y_0) \, dx \, dy + N_0 \ln(p(x_0, y_0)) \right\}$$
(C.4)

where $p(x_0, y_0)$ is the coordinate *a priori* distribution density. This is because the estimate having the minimum probability of error is the one that has maximum a posteriori probability that the coordinate has a particular value given the input signal s(x, y). It can be proven that the Cramer Rao Lower Bound on the variances of the location errors in x- and y- directions are given by [54]

$$\sigma_x^2 \ge \frac{c_x \times N_0}{4\pi^2 \iint u^2 |R(u,v)|^2 du dv},$$
(C.5)

$$\sigma_y^2 \ge \frac{c_y \times N_0}{4\pi^2 \iint v^2 |R(u,v)|^2 du dv} \,. \tag{C.6}$$

where R(u, v) is the Fourier transform of r(x, y), and u and v are the frequencies along the *x*-axis and *y*-axis. c_x and c_y are some constants. This view drives us to choose the 2D DFT as our transform to allow discarding less useful frequency components and guide the quantization of the rest of coefficients. Using the 2-D discrete Fourier transform on (C.2) and using similar procedures as in Section 3.1, the Fisher information for the location estimation in x- and y- directions are

$$J_{x_0} = \frac{2\pi^2}{M^2} \sum_{\nu=0}^{M-1} \sum_{u=0}^{N-1} \frac{u^2 |R(u,\nu)|^2}{N_0},$$
 (C.7)

$$J_{y_0} = \frac{2\pi^2}{N^2} \sum_{\nu=0}^{M-1} \sum_{u=0}^{N-1} \frac{\nu^2 |R(u,\nu)|^2}{N_0},$$
 (C.8)

where R(u,v) is the Fourier transform of s(x, y), whose size is $M \times N$, and $q^2(u,v)$ is the variance of the quantization noise on the Fourier coefficient at (u,v). From (C.7) and (C.8), we can see that the accuracy of locating the reference image in the received image depends more on the high frequency components than low frequency components, It also can be seen from (C.7) and (C.8) that the high frequency coefficients in R(u,v) should be weighted more heavily than the low frequency coefficients. This implies that high frequency components play a more important role than low frequency components in PR. So for compressing images for PR, more weight should be given to the important high frequency components instead of to the low frequencies.

Since maximizing J_{x_0} and J_{y_0} can achieve improved location accuracy in pattern recognition, (C.7) and (C.8) can be used as criteria for image compression in CPR. Applying the approximate model in Chapter 3, if we replace R(u,v) with S(u,v), and treat the x and y directions equally important, we can add F_x^2 and F_y^2 to form a single tractable and data centric distortion measure,

$$\hat{J}(s) = \hat{J}_{x_0}(s) + \hat{J}_{x_0}(s)$$

= $2\pi^2 \sum_{\nu=0}^{M-1} \sum_{u=0}^{N-1} \frac{\left(u^2 |S(u,\nu)|^2 / M^2 + \nu^2 |S(u,\nu)|^2 / N^2 \right)}{N_0 + q^2(u,\nu)}.$ (C.9)

We can use the same numerical optimization method as in Chapter 3 to design an optimal transform coder according to (C.9) under the bit constraint R. However, here we simply modify the standard compression algorithms to reflect this criterion so they can be more suitable for image PR. For example, for JPEG compression, we only change their quantization step and leave the other processing steps untouched. The advantage of doing such is that we are still able to use the standard decoders to decompress the images compressed by the proposed algorithms.

C.3 Modifying the Standard Compression Algorithms for CPR

We first modify the JPEG so it can be more suitable for CPR. The DCT is closely related to the discrete Fourier transform (DFT). Actually, DCT coefficients y_q of any *m*point sequence **x** are related to the 2m-point DFT coefficients \tilde{y}_q of the symmetrically extended sequence, \tilde{x} , by a constant scale factor, which is signal dependent; i.e.,

$$y_q = \left(c_q \sqrt{\frac{m}{2}} e^{-j\frac{\pi}{2m}q}\right) \breve{y}_q, \qquad q = 0, 1, \cdots, m-1 \quad .$$
 (C.10)

where $c_q = \sqrt{1/m}$ when q = 0, and $c_q = \sqrt{2/m}$ otherwise.

The numerator of the summation (C.7) and (C.8) can be approximated by

$$\sum_{j=0}^{M-1N-1} \sum_{i=0}^{N-1} (i^2 + j^2) T_{ij}^2 , \qquad (C.11)$$

for each of the 8×8 DCT blocks within a scaling factor that has no effect on the compression. For the criterion obtained in Section C.2, we know that high-frequency

DCT coefficients should be weighted and retained with refined quantization, whereas low-frequency DCT coefficients should be discarded or quantized coarsely. This can be done by carefully selecting a quantization array. For simplicity, we rotate the JPEG's quantization matrix by 180 deg and use it as our quantization matrix:

$$Q^* = \begin{bmatrix} 99 & 103 & 100 & 112 & 98 & 95 & 92 & 72 \\ 101 & 120 & 121 & 103 & 87 & 78 & 64 & 49 \\ 92 & 113 & 104 & 81 & 64 & 55 & 35 & 24 \\ 77 & 103 & 109 & 68 & 56 & 37 & 22 & 18 \\ 62 & 80 & 87 & 51 & 29 & 22 & 17 & 14 \\ 56 & 69 & 57 & 40 & 24 & 16 & 13 & 14 \\ 55 & 60 & 58 & 26 & 19 & 14 & 12 & 12 \\ 61 & 51 & 40 & 24 & 16 & 10 & 11 & 16 \end{bmatrix}$$

This rotated matrix can perform fine quantization for high frequency DCT coefficients and coarse quantization for low-frequency DCT coefficients, because the quantization steps are small at the lower right and large at the upper left corner. Thus the modified JPEG for CPR involves the following steps. After DCT transformation, each DCT block is weighted by a weight matrix and then quantized using the 180-deg-rotated quantization matrix Q^* . That is, (C.11) is modified as

$$\hat{T}_{ij} = round \left(\frac{T_{ij} \times W_{ij}}{q \cdot Q_{ij}^*}\right) \times mask , \qquad (C.12)$$

where $W_{i,j}$ are the components of the weight matrix and should be $\sqrt{(i^2 + j^2)}$ according to (C.12). Considering that the quantized coefficients are scanned in a zigzag order, the coding can also be simplified by changing the weight matrix to the 8×8 zigzag index matrix

$\begin{bmatrix} 0 \end{bmatrix}$	1	5	6	14	15	27	28	
2	4	7	13	16	26	29	42	
3	8	12	17	25	30	41	43	
9	11	18	24	31	40	44	53	
10	19	23	32	39	45	52	54	,
20	22	33	38	46	51	55	60	
21	34	37	47	50	56	59	61	
35	36	48	49	57	58	62	63	

where again high-frequency coefficients receive more weight than the low-frequency ones. The *mask* in (C.12) is two-fold: (1) used to further improve the efficiency of entropy coding (all-ones mask will simply do nothing); (2) combines noise tolerance and discrimination capability. Since the input image to be compressed contain noise (we use S(u,v) to replace R(u,v) in (C.9)) and non-zero extremely high frequency components are usually just noise, keeping those frequency components will jeopardize the noise tolerance capability of the decompressed image. Therefore compression algorithm for PR should be bandpass in nature to combine noise tolerance and discrimination capability and should retain middle frequencies instead of extremely high frequencies during compression. In summary, a mask can be used to delete those nonzero extremely high-and low-frequency coefficients to improve coding efficiency and also remove noise. An example of such a mask is shown below:

$\begin{bmatrix} 0 \end{bmatrix}$	0	0	0	0	0	1	1
0	0	0	0	1	1	1	0
0	0	0	1	1	1	0	0
0	0	1	1	1	1	0	0
0	1	1	1	1	0	0	0
1	1	1	1	0	0	0	0
1	1	1	0	0	0	0	0
1	1	0	0	0	0	0	0

where only coefficients of middle frequencies from 16 to 39 are encoded. The mask could be determined either by optimizing some measure or by extensive experimentation. The latter was used in this paper.

In the final stage of compression, the obtained quantized coefficients from (C. 11) are entropy-encoded. Since only middle-frequency (and some high-frequency) components are coded and they have much smaller magnitudes than low-frequency components, a higher compression ratio than the standard JPEG can be obtained. Although multiplying DCT coefficients with a weight matrix will increase the magnitudes of the quantized coefficients, this can be alleviated by increasing the corresponding quality factor. It is to be noted that the quality of an image in PR has a different meaning from the quality in human visualization. Therefore, the quality factors (C.12) and (C.2) are not the same.

We can modify the wavelet-based image compression algorithms as well for PR in a similar way by emphasizing middle- to high-frequency components. The wavelet-based algorithms are based on the multiresolution decomposition of the image to be compressed. For visualization, the LL (low-frequency) band at the highest level is considered as most important, and the other bands, containing edge information, are classified as of minor importance, with the degree of importance decreasing from the top of the pyramid to the

bands at the bottom. As a result, a large number of bits are assigned to encode the LL band, and many coefficients at higher-frequency subbands (finer resolutions) are discarded. However, for pattern recognition, we know that edge information is more important. Accordingly, we suggest encoding the LH and HL bands and discarding the LL and HH bands. In this way, a much higher compression ratio can be obtained while simultaneously enhancing the discrimination performance. Wavelet denoising [7] can be applied to the wavelet coefficients to remove possible noise in the encoding process.

C.4 Simulation Results

Computer simulations were performed to test the performance of the proposed algorithms. The input image used is an infrared image as shown in Fig. C.2, and the truck located at the left is used as the reference image. We compare the performance of the standard JPEG, the standard EZW, and the proposed modified JPEG and EZW. The input image is compressed using these four compression methods respectively, and then reconstructed and cross-correlated with the reference image. Figure C.3(a), C.3(b), C.3(c), and C.3(d) show, respectively, the reconstructed images from the standard JPEG (or EZW); the modified JPEG with q=1, weight matrix $\sqrt{i^2 + j^2}$, an all-ones mask; the modified JPEG with q=1, the zigzag weight matrix, and the example mask; and the modified EZW algorithms. It is seen that the reconstructed image from the standard JPEG or EZW is almost identical to the original image, while in the reconstructed images from the modified algorithms, the middle and high frequency components are emphasized. Figure C.3(b), C.3(c), and C.3(d) will not be acceptable for visualization purposes, but they are good for pattern recognition. This can be seen from the normalized correlation results as shown in Figure C.4. Figure C.4(a) is the normalized correlation output with

the original image (without compression) and the reference. Figure C.4(b) shows the normalized correlation obtained with the decompressed image (using the standard JPEG or EZW compression) as the input image. Figure C.4(a) and C.4(b) are almost the same. The correlation profiles are broad, and there are a lot of sidelobes.



00-00

Figure C2. The input and reference images



Figure C.3. Reconstructed image from (a) the standard JPEG or EZW compression, (b) the modified JPEG compression with the weight matrix $\sqrt{i^2 + j^2}$ and all-ones mask, (c) the modified JPEG compression with the zigzag weight matrix and the example mask, and (d) the modified EZW compression



Figure C.4. Normalized correlation outputs of the reference and the input image processed (a) without compression, (b) with JPEG or EZW compression, (c) with the modified JPEG, and (d) with the modified EZW compression.

However, with the proposed modified JPEG and EZW compression, we can obtain a much sharper correlation peak, as shown in Figure C.4(c) and C.4(d), respectively. For the compression ratio, the standard JPEG baseline and the modified JPEG with the zigzag weight matrix are 30:1 and 73.5:1, respectively whereas the standard EZW and the modified EZW have compression ratios of 533:1 and 624:1, respectively.



Figure C.5: Compression ratio versus quality factor

The compression ratio of the modified JPEG algorithm depends on the quality factor, the weight matrix, and the shape of the mask. Figure C.5 illustrates the relationship between *q*-factor and compression ratio when the zigzag weight matrix and the example mask shown in C.3 are used. Figures C.6 and C.7 show the reconstructed images and correlation distributions using the modified JPEG when the quality factor is 12 (compression ratio 117:1) and 35 (compression ratio 820:1), respectively. It is seen that we can still get sharp correlation peaks at high compression ratio. The quality of the decompressed image is degraded significantly as the compression ratio increases. The edges of the two other smaller vehicles (car and tank) are almost lost in Figure C.7 due to the extremely high compression ratio. Low contrast and small spatial support are other

reasons for losing edges. No compression method can be guaranteed to work for all objects with a high compression ratio.



Figure C.6: Reconstructed image and correlation output when q=12.



Figure C.7: Reconstructed image and correlation output when q=35.

Next, white Gaussian noise was added to the input image as shown in Figure C.8 (a) to test the noise performance of the proposed methods. The noise has zero mean and a standard deviation of 20. The normalized correlation distributions of the reference and the noisy input processed with the modified JPEG (with q=12) and the modified EZW, are obtained as shown in Figure C.8 (b) and C.8 (c), respectively. It is seen that they can tolerate additive noise. When the noise becomes heavy, the performance of the modified JPEG degrades faster than that of the modified EZW. From the results obtained, the modified EZW has the best performance, considering the compression ratio, correlation output, and noise tolerance. This is partially because the modified EZW has better

denoising and edge enhancement capabilities than the modified JPEG. In the modified JPEG, denoising and edge enhancement are simply performed by discarding low and extremely high frequencies. This approach should be improved in the future. The modified JPEG, however, has a lower computationally complexity.



Figure C.8: (a) the noisy input image; (b) normalized correlation outputs of the reference and (a) with the modified JPEG; (c) normalized correlation outputs of the reference and (a) with the modified EZW.

APPENDIX D

Symmetric Indexing

In the examples in Sections IV-A and IV-B we used a symmetric indexing of the data of which we assess the Fisher information: for TDOA estimation we symmetrically index the DFT coefficients in (3.26) and for FDOA estimation we symmetrically index the samples in (3.33). Why do we choose these seemingly arbitrary indexing schemes? Fundamentally, the choice of which sample or DFT coefficient we index as zero is an arbitrary one that arises in the mathematical analysis; thus, such an arbitrary naming should have no effect on the amount of Fisher information is available in from the data. However, for some problems the choice of the data indexing can influence the form of the Fisher information. For example, in the FDOA problem, if instead of the indexing used in (3.33) we changed the indexing of the signal samples to

$$x_{1}[n] = s[n]e^{j(v_{0}+v/2)} + w_{1}[n] \quad n = n_{o} - N/2, \dots, n_{o} - 1, n_{o}, n_{o} + 1, \dots, n_{o} + N/2$$

$$x_{2}[n] = s[n]e^{j(v_{0}-v/2)} + w_{2}[n] \quad n = n_{o} - N/2, \dots, n_{o} - 1, n_{o}, n_{o} + 1, \dots, n_{o} + N/2$$

(D.1)

then the total Fisher information can be made arbitrarily large by increasing the value of $|n_o|$. This motivates choosing the smallest possible value of $|n_o|$, or $n_o = 0$. In other words, the indexing should be symmetric in the TDOA and FDOA problems. In other problems similar attention should be paid to the impact of the indexing on the Fisher information.

APENDIX E

E.1. Correlated TDOA/FDOA Estimates

In this section, we will shows that it is indeed possible to occur that the TDOA/FDOA estimates are correlated. Let the received C-T signal be $r(t) = s(t + \tau)e^{jvt} + w(t)$ so that after sampling we have

$$r[n] = \underbrace{s(nT+\tau)e^{jwnT}}_{\to \mathbf{s}} + w(nT)$$
(E.1)

where the arrow notation means that samples of the indicated signal go into a signal vector called **s**. Then the partial derivatives of this signal vector w.r.t. the TDOA/FDOA parameters are given by

$$\frac{\partial \mathbf{s}}{\partial \tau} = \begin{bmatrix} e^{j v_n T} & \frac{\partial s(nT+\tau)}{\partial \tau} \\ \vdots \end{bmatrix}$$
(E.2)

and

$$\frac{\partial \mathbf{s}}{\partial v} = \begin{bmatrix} \vdots \\ jnTe^{jvnT}s(nT+\tau) \\ \vdots \end{bmatrix}$$
(E.3)

The diagonal elements of the FIM are given by

$$J_{11} = \frac{2}{\sigma_w^2} \left(\frac{\partial \mathbf{s}}{\partial \tau} \right)^H \left(\frac{\partial \mathbf{s}}{\partial \tau} \right) = \frac{2}{\sigma_w^2} \sum_n \left| \left(\frac{\partial s(t)}{\partial t} \Big|_{t=nT} \right) \right|^2$$
(E.4)

$$J_{22} = \frac{2}{\sigma_w^2} \left(\frac{\partial \mathbf{s}}{\partial \nu}\right)^H \left(\frac{\partial \mathbf{s}}{\partial \nu}\right) = \frac{2}{\sigma_w^2} \sum_n (nT)^2 |s(nT)|^2$$
(E.5)

The off-diagonal elements of the FIM are given by

$$J_{12} = \frac{2}{\sigma_w^2} \operatorname{Re}\left\{ \left(\frac{\partial \mathbf{s}}{\partial \nu} \right)^H \left(\frac{\partial \mathbf{s}}{\partial \tau} \right) \right\} = \frac{2}{\sigma_w^2} \operatorname{Re}\left\{ \sum_n -jnTs^*(nT+\tau) \frac{\partial s(nT+\tau)}{\partial \tau} \right\} \quad (E.6)$$

with $J_{21} = J_{12}$. Let the signal be a general complex equivalent lowpass signal; then it has the general form of

$$s(nT+\tau) = A(nT+\tau)e^{j\phi(nT+\tau)}$$
(E.7)

so that we get

$$\frac{\partial}{\partial \tau} s(nT+\tau) = e^{j\phi(nT+\tau)} \frac{\partial}{\partial \tau} A(nT+\tau) + A(nT+\tau) \frac{\partial}{\partial \tau} e^{j\phi(nT+\tau)}$$
$$= e^{j\phi(nT+\tau)} \frac{\partial}{\partial \tau} A(nT+\tau) + js(nT+\tau) \frac{\partial}{\partial \tau} \phi(nT+\tau)$$
(E.8)

Now using (E.8) in (E.6) gives

$$J_{12} = -\frac{2}{\sigma_w^2} \sum_n \operatorname{Re}\left\{ jnTs^*(nT+\tau) \left[e^{j\phi(nT+\tau)} \frac{\partial}{\partial \tau} A(nT+\tau) + js(nT+\tau) \frac{\partial}{\partial \tau} \phi(nT+\tau) \right] \right\}$$
$$= -\frac{2}{\sigma_w^2} \sum_n \operatorname{Re}\left\{ jnTs^*(nT+\tau) \left[e^{j\phi(nT+\tau)} \frac{\partial}{\partial \tau} A(nT+\tau) \right] \right\} + \frac{2}{\sigma_w^2} \sum_n nT |s(nT+\tau)|^2 \frac{\partial}{\partial \tau} \phi(nT+\tau)$$
(E.9)

Now clearly this doesn't go to zero in general, although there may be some special cases [86].

E.2. FIM for TDOA/FDOA of a Chirp Signal

Let the signal be a complex linear chirp signal given by

$$s(t) = e^{j\phi} e^{j(\alpha/2)t^2}$$
(E.10)

so that we get

$$\dot{s}(nT) = j\alpha nT e^{j\phi} e^{j(\alpha/2)(nT)^2} = j\alpha nT s(nT)$$
(E.11)

Now using (E.11) in (E.4)-(E6) gives

$$J_{\tau\tau} = 2\sum_{n} \left| \left(\dot{s}(nT) \right) \right|^{2} = 2\sum_{n} \left| j\alpha nTe^{j\phi} e^{j(\alpha/2)(nT)^{2}} \right|^{2} = 2(\alpha T)^{2} \sum_{n} n^{2}$$

$$J_{\nu\nu} = 2\sum_{n} (nT)^{2} \left| s(nT) \right|^{2} = 2T^{2} \sum_{n} n^{2}$$

$$J_{\nu\tau} = 2\operatorname{Re} \left\{ \sum_{n} -jnTs^{*}(nT) (j\alpha nTs(nT)) \right\} = 2\alpha T^{2} \sum_{n} n^{2}$$
(E.12)

or in matrix form

$$\mathbf{J} = \left(2T^2 \sum_{n} n^2\right) \begin{bmatrix} \alpha^2 & \alpha \\ \alpha & 1 \end{bmatrix}.$$
 (E.13)

In the following simulation of testing if we could apply equations (4.30), (4.31) and (4.37) by using PR filter bank and the short-time Fourier transform (STFT) to numerically compute the FIM elements and then compare the result to the analytically derived theoretical result for the linear chirp signal. Because $J_{\nu\nu}$ in (E.12) is a constant with respect to the sweep rate of the chirp signal α , so α , so we can look at the other terms relative to this term: $J_{11}/J_{22} = \alpha^2$ and $J_{12}/J_{22} = \alpha$. Plots of the numerical and theoretical results for the case of using a complex filter bank and STFT are shown in Figure E.1 and Figure E.2 respectively.



Figure E.1: Quality of numerical evaluation of FIM via complex filter bank


Figure E.2 : Quality of numerical evaluation of FIM via STFT.

Appendix F:

Complex PR Filter Banks

A. Introduction

Let $\{g_0[n], g_1[n], h_0[n], h_1[n]\}$ be a set of *N*-order FIR filters that lead to a perfect reconstruction (PR) filter bank arising from a tree cascade structure; the $g_i[n]$ are the analysis filters and the $h_i[n]$ are the synthesis filters. For this analysis we will assume a "full tree" decomposition; a four-channel example of the analysis bank is shown in Figure F.1 and the corresponding synthesis bank is shown in Figure F.2.

In this example we have assumed that the analysis filters $g_0[n]$ and $g_1[n]$ are chosen to be half-band lowpass and half-band highpass, respectively, as is typical and virtually universal in the literature. As a result of this choice, note that the channel frequency responses shown in Figure F.1 are "two sided," in that each channel passes both positive and negative frequencies.

In many applications it is desirable to have so-called one-sided channels. Although there are known methods for the direct design and implementation of such filters (e.g., DFT-based filter banks and polyphase filter banks), tree-based methods for one-sided filter banks seem to be relatively unknown. Because there are times when a tree-based approach is desirable for one-sided filter banks it is worthwhile to address a simple trick that allows conversion of any two-sided tree-based PR filter bank into a one-sided tree-based PR. This then brings to bear on the one-sided problem a huge body of literature devoted to the two-sided problem.

B. Channel Response Analysis for Two-Sided Filter Banks

To aid our later analysis we will analyze this 4-channel filter bank to verify that it provides the kind of channel responses shown in Figure F.1. The easiest way to do this is to use the decimation identity to move the "inside" decimators to the output side, then find the four channel's outputs before decimation, then impart the effect of the 4x decimation. The decimation



Figure F.1: Example of a four-channel, full-tree cascaded analysis filter bank



Figure F.2 : Example of a four-channel, full-tree cascaded synthesis filter bank

identity says that a 2x decimator followed by some H(z) can be replaced by $H(z^2)$ followed by 2x decimation. Furthermore, we know that if transfer function H(z) gives frequency response $H(\theta)$ then transfer function $H(z^2)$ gives frequency response $H(2\theta)$. Thus, the two-stage cascade for the four-channel analysis filter can be replaced by the system in Figure F.3 and then that in Figure F.4.

Figure F.5 shows the four ideal filters needed for the analysis: $G_0(\theta)$ and $G_1(\theta)$ in the first stage and $G_0(2\theta)$ and $G_1(2\theta)$ in the second stage (assuming the decimation moved to the output side). Figure F.6 shows the pairs of cascaded filter responses superimposed on top of each other and Figure F.7 shows the result of multiplying each pair of cascaded filter responses to get the channel responses– note that each channel has a (positive side) passband width of $\pi/4$ rad/second and that the responses do not overlap but seamlessly cover the whole frequency range. Figure F.8 shows the effect of each channel on the input DTFT and gives the output DTFT *before* any decimation – note that each channel's output is a perfect slice of the input with nothing existing outside the channel's ideal range.

Figure F. 9 through Figure F.12 shows the effect of decimation on the output DTFT of each channel – notice that there is no overlap between the replica components that are present due to decimation and therefore there is no aliasing. Finally, Figure F. 13 shows the outputs after decimation – since there is no aliasing the shapes of these look just like those shown for the predecimation channel outputs in Figure F.8, except that the decimation has stretched them to be "full band" and has flipped them left-right for channels 2 and 3.



Figure F.3: Equivalent filter bank with decimators moved



Figure F. 4: Equivalent filter bank with decimators moved and blocks combined



Figure F. 5: Ideal filters for the various stages (w/ decimation moved to end)



Figure F. 6 : Shows pairs of filters that appear in the four different cascades that form the four channels



Figure F.7: Four channels' frequency response (w/ decimation moved to end)



Figure F. 8 : The four channels and how they split up the input DTFT



Figure F.9 : Channel 1 after decimation - shows all the streched spectral replicas that arise from decimation



Figure F. 10 : Channel 2 after decimation - shows all the stretched spectral replicas that arise from decimation



Figure F. 11: Channel 4 after decimation - shows all the stretched spectral replicas that arise from decimation



Figure F. 12: Channel 3 after decimation - shows all the stretched spectral replicas that arise from decimation



Figure F.13: Channel output DTFT's after decimation - only shows standard $-\pi$ to π range

C. Background: Standard Conditions For PR

First some preliminaries on the conditions used to achieve PR in the two-sided filter banks. The order N of the FIR filters must be odd (and therefore the length L = N+1 is even). Cancellation of aliasing is ensured when the synthesis filters are related to the analysis filters according to

$$H_0(z) = G_1(-z)$$

 $H_1(z) = -G_0(-z)$
(F.1)

which in the time domain becomes

$$h_0[n] = (-1)^n g_1[n], \qquad n = 0, 1, 2, ..., N$$

$$h_1[n] = -(-1)^n g_0[n], \qquad n = 0, 1, 2, ..., N$$
 (F.2)

With aliasing cancelled through this relationship between the synthesis filters and the analysis filters, the transfer function from input on the analysis side to the output on the synthesis side becomes

$$T(z) = G_0(z)G_1(-z) - G_0(-z)G_1(z)$$
(F.3)

The following relationship between the two analysis filters is imposed to help enforce the PR condition:

$$G_1(z) = (-z)^{-N} G_0(-z^{-1})$$
(F.4)

which in the time domain is equivalent to

$$g_1[n] = (-1)^n g_0[N-n], \qquad n = 0, 1, 2, \dots, N$$
 (F.5)

which turns (F.2) into

$$h_0[n] = g_0[N - n],$$
 $n = 0, 1, 2, ..., N$
 $h_1[n] = -(-1)^n g_0[n],$ $n = 0, 1, 2, ..., N$ (F.6)

Substitution of (F.4) into (F.3) leads to

$$T(z) = z^{-N} [G_0(z)G_0(z^{-1}) - G_0(-z)G_0(-z^{-1})].$$
 (F.7)

and thus a sufficient condition for perfect reconstruction is

$$|G_0(\theta)|^2 + |G_0(\theta - \pi)|^2 = 1$$
, (F.8)

where $G_0(\theta)$ is the frequency response of the filter; in the time domain this condition is equivalent to

$$\left\langle g_0[n], g_0[n+2k] \right\rangle = \delta[k], \tag{F.9}$$

which means that the filter $g_0[n]$ is orthogonal to its even-integer translates. Thus, a set of filters designed to satisfy (F.5)–(F.9) will lead to a two-sided PR filter bank. The typical approach is to design a half-band lowpass filter $g_0[n]$ that satisfies (F.8) and (F.9) and then determine the other three filters from (F.5) and (F.6). The relationships in (F.5) and (F.6) constitute a strict linkage between the four filters in the filter banks. The equations given in (F.5) and (F.6) show how to get the other three filters once you have $g_0[n]$; similarly, given any one of the four filters it is possible to get the other three using similar relationships; these relationships are summarized in Figure F.14, where the solid line connections correspond to (F.5) and (F.6) and the dashed line connections can be derived from them.

When the filter $g_0[n]$ is chosen to be a half-band lowpass filter, then $g_1[n]$ will become a halfband highpass filter, as imposed by (F.5); and it is that structure that leads to the standard twosided filter bank. Thus, to get a one-sided filter bank it seems that we may simply need to chose $g_0[n]$ in a slightly different way. As shown in the next subsection, if we first choose $g_0[n]$ to be a complex half-band filter that satisfies (F.8) then we can get the desired complex one-sided filter bank through a similar set of relationships as (F.5) and (F.6).



Figure F. 14 : Relationships between the filters for real PR filter banks.

D. One-Sided Filter Banks

We consider the same structure as in Figure F.1, but with a different choice for $g_i[n]$ than was made above. To aid in seeing the structure we will analyze the structure using ideal filters in the same that was done in Subsection C of this appendix. Let $\{\tilde{g}_0[n], \tilde{g}_1[n], \tilde{h}_0[n], \tilde{h}_1[n]\}$ be the filters that give rise to the desired complex one-sided filter bank. Choose $\tilde{g}_0[n]$ to be a positiveband filter (i.e., a half-band filter that covers the positive frequencies); this can be easily done by choosing

$$\widetilde{g}_0[n] = e^{j(\pi/2)n} g_0[n], \qquad (F.10)$$

which is just a modulation up by $\pi/2$ rad/sample of $g_0[n]$, a standard half-band lowpass filter for a two-sided filter bank. Now to get the corresponding negative-band filter, (F.5) states that we should use

$$\widetilde{g}_{1}[n] \stackrel{?}{=} (-1)^{n} \widetilde{g}_{0}[N-n]$$
 (F.11)

where the question mark above the equal sign shows our uncertainty if this result still holds. A quick numerical test of this idea shows that, in fact, (F.11) does NOT give a negative-frequency band filter! Why is that? Why did it work for the real filter case? Well, for the real filter case

the time reversal does not change the magnitude of the frequency response and then the modulation converts this into a highpass filter. However, for a complex filter the time reversal DOES change the magnitude of the frequency response; it changes it into a negative-band filter! So maybe we could just do the time reversal of (F.11) to get the desired negative-band filter; but that would provide us with a theory for the complex case that does NOT naturally degenerate to the known real-valued theory.

To find the answer we have to look at (F.4), which is the truer starting point than (F.5). The inverse-z in (F.4) (which is responsible for the time reversal) causes the filter transfer function zeros to move as show in Figure F.15; for real filters the zeros must appear in conjugate pairs and they then move into conjugate pairs. Half-band lowpass filters will have zeros along the left-had side of the unit circle and the inverse-z mapping will keep them on that side; thus, they remain half-band lowpass filters after time reversal. However, for half-band positive-pass filters a zero on the negative frequency side of the unit circle will get mapped into the positive frequency side (which is not what we want). But... conjugating the filter coefficients and doing a time reversal will leave them on the negative frequency side as shown in Figure F.16. Note that this operation will naturally degenerate to the known real-valued theory.

Thus, we have demonstrated that a time reversal and coefficient conjugation is the mapping needed to get the negative-band filter from the positive band filter; thus, (F.11) should be changed to

$$\widetilde{g}_{1}[n] = (-1)^{n} \widetilde{g}_{0}^{*}[N-n], \qquad (F.12)$$

which now covers the negative frequency band. This is shown in Figure F.17.

If we now walk our way through the equivalent four-channel cascade with the decimators moved to the end we will see what kind of analysis bank this choice of filters will give. The plots in Figure F.17 show the interaction between the two filters in each channel. The plots in Figure F.18 show the frequency response magnitudes for each of the four filter bank channels. Figure F.19 shows the four channels side by side, where the filter responses are shown as trapezoids to enable on to easily see the different channels. From this is it is clear that this arrangement gives us the desired one-sided channel frequency responses. But... it is important to note that the frequency ordering of the channels can be counter-intuitive (e.g., note in Figure F.19 that the response of $|G_1(\theta)| |G_1(2\theta)|$ does not give the most negative frequency channel as one might expect).

The diagram in Figure F.20 shows the modifications made to the relationships in Figure F.14 to handle the complex filter case. Note that the diagonal relationships that ensure aliasing cancellation are the same. Thus, all we have to ensure is that by choosing according to (F.12) still ensures PR when $g_0[n]$ is a standard real-valued PR lowpass filter.



Figure F. 15: Movement of real filter zeros under $z \rightarrow z^{-1}$.



Figure F.16 : Movement of complex filter zeros under coefficient conjugation and $z \rightarrow z^{-1}$; the first move to the gray star location is due to the $z \rightarrow z^{-1}$ and the second move is due to coefficient conjugation.



Figure F.16 : Ideal Filters for the first stage of the one-sided filter bank



Figure F.17: Shows pairs of filters that appear in the four different cascades that form the four channels



Figure F. 18 : Four Channels' frequency response (w/ decimation moved to end)



Figure F.19: One-Sided filter bank's frequency responses (w/ decimation moved to end)



Figure F.20: Relationships between the filters for complex PR filter banks.

One question remains: Do we still get perfect reconstruction from this scheme? The answer is yes, as long as the original two-sided $g_0[n]$ leads to a perfect reconstruction filter bank. That this is true follows easily.

Let $\hat{g}_0[n] = \tilde{g}_0^*[n]$. Then (F.12) becomes

$$\widetilde{G}_{1}(z) = (-z)^{-N} \hat{G}_{0}(-z^{-1})$$
 (F.13)

and total transfer function in (F.3) becomes

$$T(z) = \widetilde{G}_0(z)\widetilde{G}_1(-z) - \widetilde{G}_0(-z)\widetilde{G}_1(z) .$$
(F.14)

which after substituting (F.13) becomes

$$T(z) = z^{-N} \left[\widetilde{G}_0(z) \hat{G}_0(z^{-1}) - \widetilde{G}_0(-z) \hat{G}_0(-z^{-1}) \right].$$
(F.15)

where the part in front of the brackets is a pure delay needed for causality. To get PR we need the frequency response of the part in the brackets to be constant over frequency. To check this we have

$$\widetilde{G}_{0}(z)\widehat{G}_{0}(z^{-1}) - \widetilde{G}_{0}(-z)\widehat{G}_{0}(-z^{-1})$$

$$\downarrow$$

$$\widetilde{G}_{0}(\theta)\widetilde{G}_{0}^{*}(\theta) - \widetilde{G}_{0}(\theta - \pi)\widetilde{G}_{0}^{*}(\theta - \pi)$$
(F.16)

where the conjugated terms in bottom of (F.16) come from

$$\hat{G}_{0}(z^{-1}) = \sum_{n} \widetilde{g}_{0}^{*}[n] z^{+n}$$

$$\rightarrow \sum_{n} \widetilde{g}_{0}^{*}[n] e^{+j\theta n} = \left[\sum_{n} \widetilde{g}_{0}[n] e^{-j\theta n} \right]^{*}, \quad (F.17)$$

$$= \widetilde{G}_{0}^{*}(\theta)$$

Then the bottom of (F.16) becomes

$$\left|\widetilde{G}_{0}(\theta)\right|^{2} + \left|\widetilde{G}_{0}(\theta - \pi)\right|^{2} = 1.$$
(F.18)

That this is true whenever the original filter satisfies (F.8) follows from

$$\left| \widetilde{G}_{0}(\theta) \right|^{2} + \left| \widetilde{G}_{0}(\theta - \pi) \right|^{2} = \left| G_{0}(\theta - \pi/2) \right|^{2} + \left| G_{0}(\theta - \pi - \pi/2) \right|^{2}$$
$$= \left| G_{0}(\xi) \right|^{2} + \left| G_{0}(\xi - \pi) \right|^{2} = 1$$
, (F.19)

where we have used a change of variables.

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